

Windows zero-day exploit used in targeted attacks by FruityArmor APT

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A few days ago, Microsoft published the “critical” [MS16-120 security bulletin](#) with fixes for vulnerabilities in Microsoft Windows, Microsoft Office, Skype for Business, Silverlight and Microsoft Lync.

One of the vulnerabilities – CVE-2016-3393 – was reported to Microsoft by Kaspersky Lab in September 2016.

MS16-120	Win32k Elevation of Privilege Vulnerability	CVE-2016-3270	pgboy, zhong_sf of Qihoo 360 Vulcan Team
MS16-120	Windows Graphics Component RCE Vulnerability	CVE-2016-3393	Anton Ivanov of Kaspersky Lab
MS16-120	True Type Font Parsing Elevation of Privilege Vulnerability	CVE-2016-7182	Mateusz Jurczyk of Google Project Zero
MS16-119	Microsoft Browser Information Disclosure Vulnerability	CVE-2016-3267	Wenxiang Qian of Tencent QQBrowser

Here’s a bit of background on how this zero-day was discovered. A few of months ago, we deployed a new set of technologies in our products to identify and block zero-day attacks. These technologies proved their effectiveness earlier this year, when we discovered two [Adobe Flash zero-day exploits](#) – CVE-2016-1010 and CVE-2016-4171. Two Windows EoP exploits have also been found with the help of this technology. One is CVE-2016-0165. The other is CVE-2016-3393.

Like most zero-day exploits found in the wild today, CVE-2016-3393 is used by an APT group we call **FruityArmor**. FruityArmor is perhaps a bit unusual due to the fact that it leverages an attack platform that is built entirely around PowerShell. The group’s primary malware implant is written in PowerShell and all commands from the operators are also sent in the form of PowerShell scripts.

In this report we describe the vulnerability that was used by this group to elevate privileges on a victim’s machine. Please keep in mind that we will not be publishing all the details about this vulnerability because of the risk that other threat actors may use them in their attacks.

Attack chain description

To achieve remote code execution on a victim’s machine, FruityArmor normally relies on a browser exploit. Since many modern browsers are built around sandboxes, a single exploit is generally not sufficient to allow full access to a targeted machine. Most of the recent attacks we’ve seen that rely on a browser exploit are combined with an EoP exploit, which allows for a reliable sandbox escape.

In the case of FruityArmor, the initial browser exploitation is always followed by an EoP exploit. This comes in the form of a module, which runs directly in memory. The main goal of this module is to unpack a specially crafted TTF font containing the CVE-2016-3393 exploit. After unpacking, the module directly loads the code exploit from memory with the help of AddFontMemResourceEx. After successfully leveraging CVE-2016-3393, a second stage payload is executed with higher privileges to execute PowerShell with a meterpreter-style script that connects to the C&C.

EOP zero-day details

The vulnerability is located in the **cjComputeGLYPHSET_MSFT_GENERAL** function from the Win32k.sys system module. This function parses the cmap table and fills internal structures. The CMAP structure looks like this:

Type	Name
USHORT	format
USHORT	length
USHORT	language
USHORT	segCountX2
USHORT	searchRange
USHORT	entrySelector
USHORT	rangeShift
USHORT	endCount[segCount]
USHORT	reservedPad
USHORT	startCount[segCount]
SHORT	idDelta[segCount]
USHORT	idRangeOffset[segCount]
USHORT	glyphIdArray[]

The most interesting parts of this structure are two arrays – endCount and startCount. The exploit contains the next cmap table with segments:

```

Length:      48
Version:     0
segCount:    28 (X2 = 56)
searchRange: 32
entrySelector: 4
rangeShift: 24
Seg  1 : St = 0000, En = 1388, D =      0, RO =      1, gId# = -28
Seg  2 : St = 1388, En = 1770, D =      0, RO =      1, gId# = -27
Seg  3 : St = 1770, En = 1B58, D =      0, RO =      1, gId# = -26
Seg  4 : St = 1B58, En = 1F40, D =      0, RO =      1, gId# = -25
Seg  5 : St = 1F40, En = 2328, D =      0, RO =      1, gId# = -24
Seg  6 : St = 2328, En = 238C, D =      0, RO =      1, gId# = -23
Seg  7 : St = 238C, En = 23F0, D =      0, RO =      1, gId# = -22
Seg  8 : St = 23F0, En = 2454, D =      0, RO =      1, gId# = -21
Seg  9 : St = 2454, En = 24B8, D =      0, RO =      1, gId# = -20
Seg 10 : St = 24B8, En = 251C, D =      0, RO =      1, gId# = -19
Seg 11 : St = 251C, En = 2580, D =      0, RO =      1, gId# = -18
Seg 12 : St = 2580, En = 25E4, D =      0, RO =      1, gId# = -17
Seg 13 : St = 25E4, En = 2648, D =      0, RO =      1, gId# = -16
Seg 14 : St = 2648, En = 2710, D =      0, RO =      1, gId# = -15
Seg 15 : St = 2710, En = 3A98, D =      0, RO =      1, gId# = -14
Seg 16 : St = 3A98, En = 4E20, D =      0, RO =      1, gId# = -13
Seg 17 : St = 4E20, En = 61A8, D =      0, RO =      1, gId# = -12
Seg 18 : St = 61A8, En = 7530, D =      0, RO =      1, gId# = -11
Seg 19 : St = 7530, En = 88B8, D =      0, RO =      1, gId# = -10
Seg 20 : St = 88B8, En = 9C40, D =      0, RO =      1, gId# = -9
Seg 21 : St = 9C40, En = AFC8, D =      0, RO =      1, gId# = -8
Seg 22 : St = AFC8, En = C350, D =      0, RO =      1, gId# = -7
Seg 23 : St = C350, En = D6D8, D =      0, RO =      1, gId# = -6
Seg 24 : St = D6D8, En = EA60, D =      0, RO =      1, gId# = -5
Seg 25 : St = EA60, En = EE48, D =      0, RO =      1, gId# = -4
Seg 26 : St = EE48, En = F230, D =      0, RO =      1, gId# = -3
Seg 27 : St = F231, En = FFFE, D =      0, RO =      1, gId# = -2
Seg 28 : St = FFFF,      = FFFF, D =      0, RO =      1, gId# = -1

```

To compute how much memory to allocate to internal structures, the function executes this code:

```

USHORT cnt=0;
for (int i=0; i<cnt;i++)
    cnt += (Seg[i].end - Seg[i].start)+1

```

After computing this number, the function allocates memory for structures in the following way:

```

text:BF9EAB40 loc_BF9EAB40:                ; CODE XREF: cjComputeGLYPHSET_MSFT_GENERAL(x,x,x,x)+757j
text:BF9EAB40      movzx  eax, dx
text:BF9EAB43      push  64667454h      ; Tag
text:BF9EAB48      shl   eax, 3
text:BF9EAB4B      push  eax            ; int
text:BF9EAB4C      push  ebx            ; char
text:BF9EAB4D      call  EngAllocMem@12 ; EngAllocMem(x,x,x)

```

The problem is that if we compute the entire table, we will achieve an integer overflow and the `cnt` variable will contain an incorrect value.

In kernel, we see the following picture:

<pre>kd> bc 1 kd> g Breakpoint 2 hit win32k!cjComputeGLYPHSET_MSFT_GENERAL+0xb7: 9556ab40 0fb7c2 movzx eax,dx kd> r edx edx=00010018 kd> p win32k!cjComputeGLYPHSET_MSFT_GENERAL+0xba: 9556ab43 6854746664 push 64667454h kd> r eax eax=00000018 kd> p win32k!cjComputeGLYPHSET_MSFT_GENERAL+0xbf: 9556ab48 c1e003 shl eax,3 kd> r eax eax=00000018 = Cnt kd> p win32k!cjComputeGLYPHSET_MSFT_GENERAL+0xc2: 9556ab4b 50 push eax kd> r eax eax=000000c0 Cnt*sizeof(internalStruct)</pre>	<pre>9556ab3e 33db xor ebx,ebx 9556ab40 0fb7c2 movzx eax,dx 9556ab43 6854746664 push 64667454h 9556ab48 c1e003 shl eax,3 9556ab4b 50 push eax 9556ab4c 53 push ebx 9556ab4d e8318ee9ff call win32k!EngAllocMem (954 9556ab52 8945ec mov dword ptr [ebp-14h],eax 9556ab55 3bc3 cmp eax,ebx 9556ab57 750c jne win32k!cjComputeGLYPHSE 9556ab59 8b4510 mov eax,dword ptr [ebp+10h] 9556ab5c 8918 mov dword ptr [eax],ebx 9556ab5e 33c0 xor eax,eax 9556ab60 e935010000 jmp win32k!cjComputeGLYPHSE 9556ab65 3bfb cmp edi,ebx 9556ab67 8b55f0 mov edx,dword ptr [ebp-10h] 9556ab6a 8955f0 mov dword ptr [ebp-10h],edx 9556ab6d 895dfc mov dword ptr [ebp-4],ebx 9556ab70 895d08 mov dword ptr [ebp+8],ebx 9556ab73 0f8e04010000 jle win32k!cjComputeGLYPHSE 9556ab79 8b4df4 mov ecx,dword ptr [ebp-0Ch] 9556ab7c 33c0 xor eax,eax 9556ab7e 2bca sub ecx,edx 9556ab80 894de4 mov dword ptr [ebp-1Ch],ecx 9556ab83 eb03 jmp win32k!cjComputeGLYPHSE</pre>
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The code allocates memory only for 0x18 InternalStruct but then there is a loop for all the segments range (this value was extracted from the file directly):

```

v38 = EngAllocMem(v12, 8 * (unsigned __int16)v14, 'dftT');
if ( v38 )
{
    v44 = 0;
    v46 = 0;
    if ( v13 > 0 )
    {
        v18 = 0;
        v19 = v40 - (_DWORD)v39;
        for ( j = v40 - (_DWORD)v39; ; v19 = j )
        {
            LOBYTE(v20) = *(unsigned __int16 *)((char *)v39 + v19) >> 8;
            HIBYTE(v20) = *(unsigned __int16 *)((char *)v39 + v19);
            v21 = v20;
            HIBYTE(v22) = *v39;
            v34 = v20;
            LOBYTE(v22) = *v39 >> 8;
            v23 = v22;
            v33 = v22;
            if ( v22 >= v34 && v21 != -1 )
            {
                LOWORD(v24) = v21;
                v42 = v21;
                if ( v21 <= v23 )
                {
                    v25 = (unsigned __int16 *)(v35 + 2 * v18);
                    v41 = v21;
                    do
                    {
                        v26 = v38 + 8 * v44;
                        *(_DWORD *)v20 = 0,

```

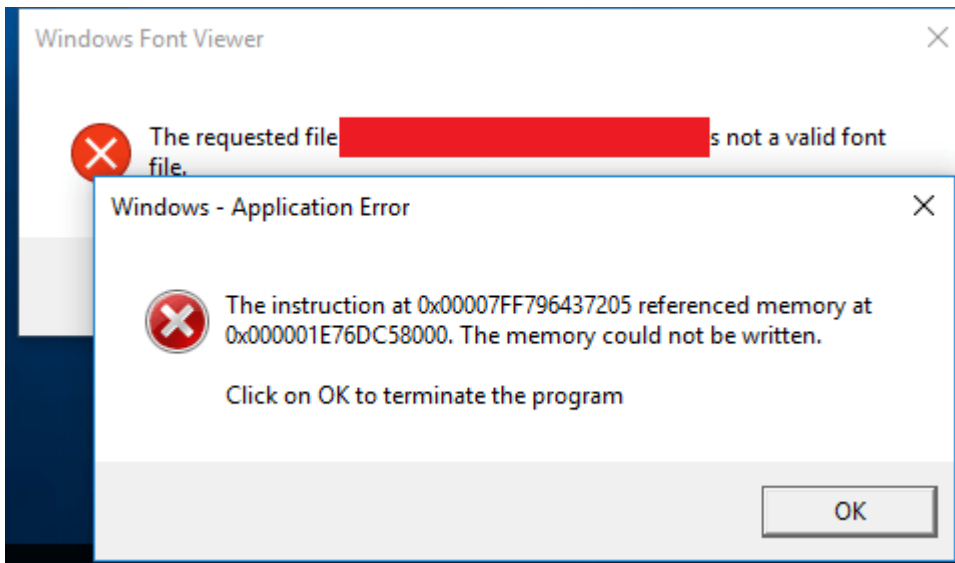
Using the cmap table, the v44 variable (index) could be controlled and, as a result, we get memory corruption. To achieve it, the attacker can do the following:

1. 1 Make an integer overflow in win32k!cjComputeGLYPHSET_MSFT_GENERAL
2. 2 Make a specific segment ranges in font file to access interesting memory.

What about Windows 10? As most of you know, the font processing in Windows 10 is performed in a special user mode process with restricted privileges. This is a very good solution but the code has the same bug in the TTF processing.

```
.text:00000001400565B9      jmp     qa, r100
.text:00000001400565BB      iz     short loc_1400565C5
.text:00000001400565BE      sub    ax, cx
.text:00000001400565C1      inc    ax
.text:00000001400565C5      add    r8u, ax
.text:00000001400565C5      loc_1400565C5:                                ; CODE XREF:
.text:00000001400565C5      ; cjComputeGl
.text:00000001400565C8      inc    dx
.text:00000001400565CC      add    rbx, 2
.text:00000001400565CF      movzx eax, dx
.text:00000001400565D3      add    r9, 2
.text:00000001400565D5      cmp    eax, ebp
.text:00000001400565D7      jl     short loc_1400565A1
.text:00000001400565D7      loc_1400565D7:                                ; CODE XREF:
.text:00000001400565D7      movzx edx, r8u
.text:00000001400565DB      xor    ecx, ecx ; uFlags
.text:00000001400565DD      shl    rdx, 3 ; duBytes
.text:00000001400565E1      call  cs:imp_Globa101loc
.text:00000001400565E7      mov    r01, rax
```

As a result, if you load/open this font exploit in Windows 10, you will see the crash of fontdrvhost.exe:



Kaspersky Lab detects this exploit as:

- HEUR:Exploit.Win32.Generic
- PDM:Exploit.Win32.Generic

We would like to thank Microsoft for their swift response in closing this security hole.

* More information about the FruityArmor APT group is available to customers of Kaspersky Intelligence Services. Contact: intelreports@kaspersky.com

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Source: <https://securelist.com/windows-zero-day-exploit-used-in-targeted-attacks-by-fruityarmor-apt/76396/>