

Tackling Anti-Analysis Techniques of GuLoader and RedLine Stealer

By Mark Lim, Zong-Yu Wu

Published: 2024-01-05 · Archived: 2026-04-05 13:51:36 UTC

Executive Summary

Malware, like many complex software systems, relies on the concept of software configuration. Configurations establish guidelines for malware behavior and they are a common feature among the various malware families we examine. The configuration data embedded within malware can offer invaluable insights into the intentions of cybercriminals. However, due to its significance, malware authors deliberately make configuration data challenging to parse statically from the file.

Over the past few years, we have developed a system to extract internal malware configurations. We will share code from our extractors for multiple malware families with the research community. These extractors, written in Python, are designed to scan and extract configuration data from memory dumps associated with specific malware samples.

We will also introduce selected configuration protection techniques employed by two malware families: GuLoader and RedLine Stealer. For those interested in more details, please look into the [whitepaper](#), [slides](#) or [video](#) we [presented](#) at Virus Bulletin 2023 in London.

Palo Alto Networks customers are better protected from these threats through our [Next-Generation Firewall](#) with [cloud-delivered security services](#) including [WildFire](#). If you think you might have been compromised or have an urgent matter, get in touch with the [Unit 42 Incident Response team](#).

Technical Analysis of GuLoader

The GuLoader authors went to great lengths to obfuscate their C2 configuration. Figure 1 provides a timeline illustrating the evolution of GuLoader obfuscation techniques.

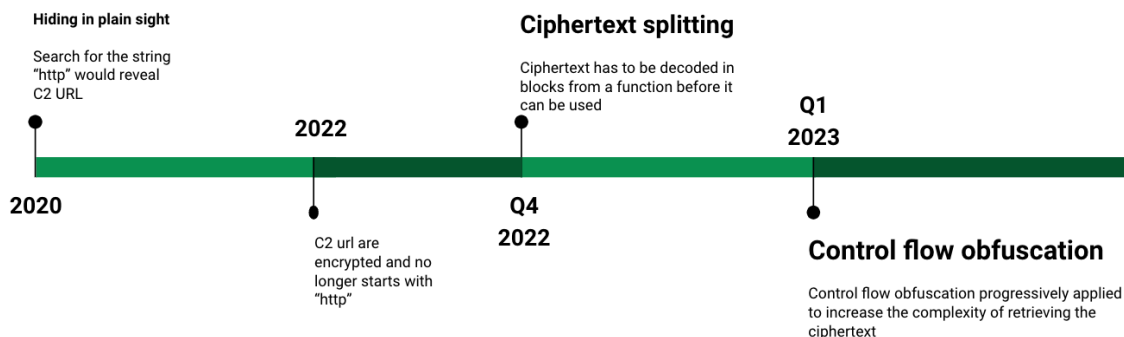


Figure 1. Timeline showing the evolution of obfuscation techniques used by Guloader.

This evolution has defeated our [previous approach](#) to extracting GuLoader malware configuration. The GuLoader authors' newer techniques include [ciphertext splitting](#) and [control flow obfuscation](#).

Ciphertext Splitting

We have labeled GuLoader's previous method of storing encrypted configuration data (ciphertext) in the top section of Figure 2 as the "old method." In this old method, the ciphertext was stored as a continuous sequence of bytes.

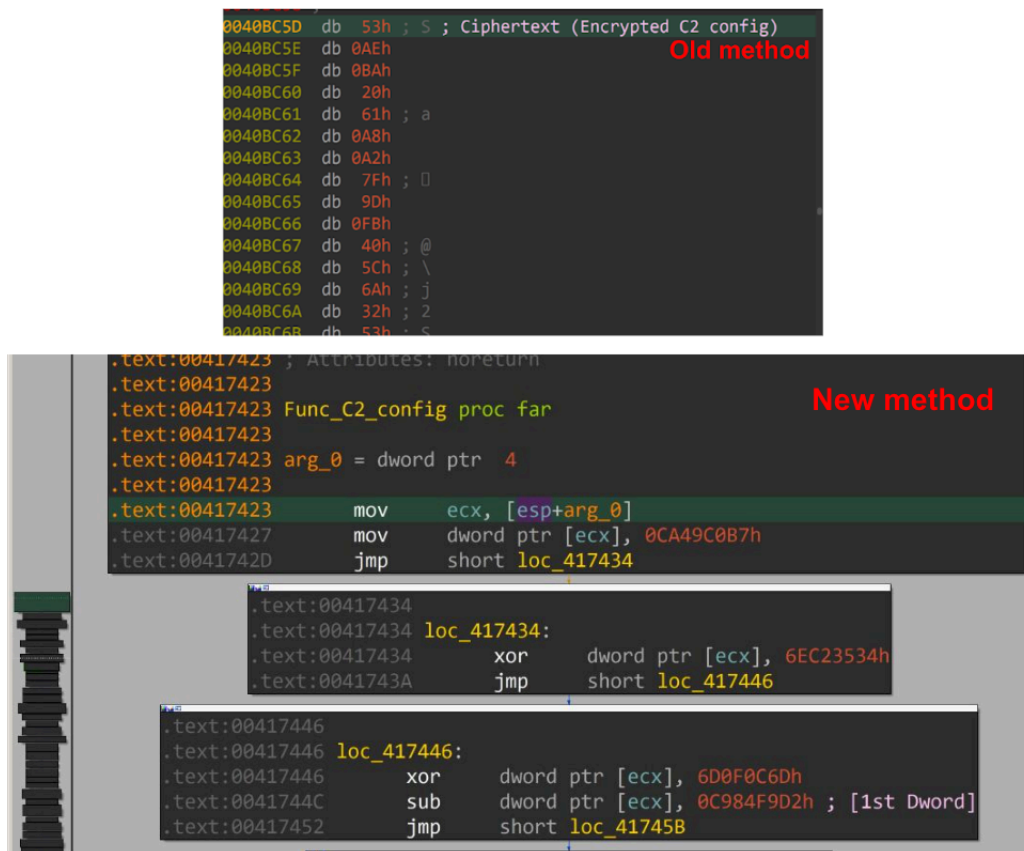


Figure 2. Comparing old and new methods of storing ciphertext.

In the lower section of Figure 2 above, we have labeled the new approach as GuLoader's "new method," where the ciphertext is computed from a function. In this function, the ciphertext is first divided into a 4-byte [DWORD](#). Each DWORD is individually encrypted using randomized mathematical operations.

For example, to retrieve the first DWORD of the ciphertext from GuLoader's new method, we must perform the mathematical operations illustrated below in Figure 3.

```

1  B = [0xCA9C0B7, 0x6EC23534, 0x6D0F0C6D, 0xC984F9D2]
2  b1 = (B[0] ^ B[1]) & 0xFFFFFFFF
3  c1 = (b1 ^ B[2]) & 0xFFFFFFFF
4  result = (c1 - B[3]) & 0xFFFFFFFF

```

Figure 3. An example of computing a DWORD of the ciphertext from Guloder’s new method of storing ciphertext.

To acquire the complete ciphertext from this new method, we perform a series of operations similar to the method shown in Figure 3 above for each individual DWORD. Subsequently, we concatenate these DWORD values together, resulting in the complete ciphertext.

Control Flow Obfuscation

In early 2023, we encountered a [GuLoader sample](#) that originally had zero VirusTotal (VT) detections. Using Hex-Rays IDA Pro to disassemble and analyze this malware sample, we found instructions that attempted to prevent further analysis. These anti-analysis instructions were designed to cause EXCEPTION_BREAKPOINT, EXCEPTION_ACCESS_VIOLATION and EXCEPTION_SINGLE_STEP violations.

Figure 4 illustrates how GuLoader implemented all these instructions for anti-analysis.

```

02B993E3 Func_C2_ciphertext_construct:
02B993E3 mov     eax, [esp+4]
02B993E7 int     3 ; EXCEPTION_BREAKPOINT triggered!
02B993E7 ;
02B993E8 db     0B4h,0A1h,'1',13h,0D8h,0FFh,'+',0A5h,0E3h,30h,26h,'t'
02B993F6 ;
02B993F6 push   eax
02B993F7 mov     eax, 433D4EDBh
02B993FC xor     eax, 9436930Fh
02B99401 xor     eax, 0A78DF586h
02B99406 sub     eax, 392289D1h
02B99408 sub     eax, 37639D81h ; eax=0x100
02B99410 push   ebx
02B99411 pushf
02B99412 mov     ebx, esp
02B99414 or      [ebx], eax ; enable Trap flag
02B99416 popf
02B99417 test    edi, ecx ; EXCEPTION_SINGLE_STEP triggered!
02B99417 ;
02B99419 aW db     'w',6,0B7h,0B6h,0CFh,14h,'\\',0ACh,0A6h,0B8h,'% ',0
02B99425 ;
02B99425 cmp     ebx, ecx
02B99427 pop     ebx
02B99428 cmp     eax, edx
02B9942A pop     eax
02B9942B mov     dword ptr [eax], 3D6E3A33h
02B99431 push   edi
02B99432 mov     edi, 82A8E946h
02B99437 sub     edi, 9707FB40h
02B9943D sub     edi, 0EBA0EE06h
02B99443 mov     [edi], edi
02B99445 mov     esp, offset unk_78C8DF
02B9944A pop     edi
02B9944B xor     dword ptr [eax], 8344C565h
02B99451 xor     dword ptr [eax], 74FD1098h
02B99457 sub     dword ptr [eax], 0CAD7EF8Ch
02B9945D push   edi
02B9945E mov     edi, 0AA2AD49Dh
02B99463 add     edi, 676C9CAAh
02B99469 add     edi, 0EE688EB9h ; edi = 0x0
02B9946F mov     [edi], ecx ; EXCEPTION_ACCESS_VIOLATION triggered!
02B9946F ;
    
```

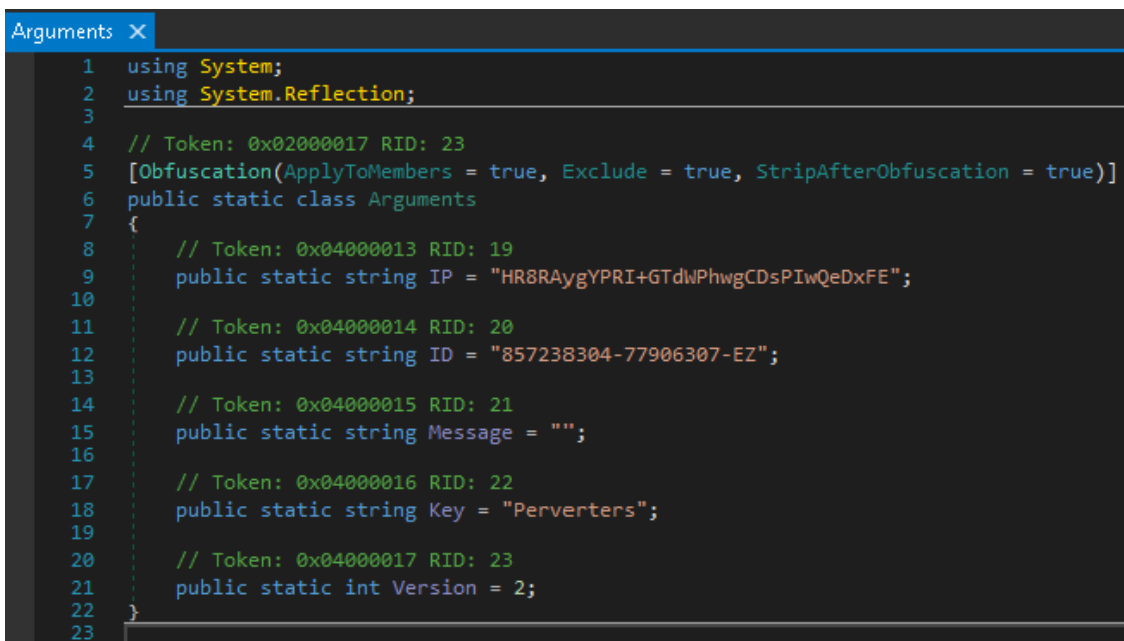
Figure 4. Disassembler analysis of the Guloder sample revealing the anti-analysis instructions.

The anti-analysis instructions noted in Figure 4 above rendered [our previous solution](#) of writing an IDA processor module extension ineffective. Due to the variable nature of the length of Intel x86 CPU instructions, we could not detect the huge combination of instructions that triggered EXCEPTION_ACCESS_VIOLATION and EXCEPTION_SINGLE_STEP exceptions.

Since our previous solution was no longer effective, we had to manually analyze the code to find these anti-analysis instructions and bypass them to extract the configuration. We explained in detail how we extracted the configuration in [our whitepaper for Virus Bulletin](#).

Technical Analysis of RedLine Stealer

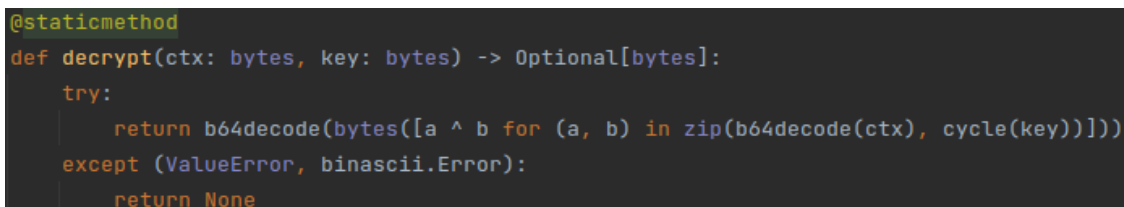
The SHA256 hash for the RedLine Stealer sample used in this analysis is a4cf69f849e9ea0ab4eba1cdc1ef2a973591bc7bb55901fdbceb412fb1147ef9. Using an MSIL decompiler called [dnSpy](#), we quickly identified the configuration data as shown below in Figure 5.



```
Arguments X
1 using System;
2 using System.Reflection;
3
4 // Token: 0x02000017 RID: 23
5 [Obfuscation(ApplyToMembers = true, Exclude = true, StripAfterObfuscation = true)]
6 public static class Arguments
7 {
8     // Token: 0x04000013 RID: 19
9     public static string IP = "HR8RAYgYPRI+GTdWPhwgCDsPIwQeDxFE";
10
11     // Token: 0x04000014 RID: 20
12     public static string ID = "857238304-77906307-EZ";
13
14     // Token: 0x04000015 RID: 21
15     public static string Message = "";
16
17     // Token: 0x04000016 RID: 22
18     public static string Key = "Perversers";
19
20     // Token: 0x04000017 RID: 23
21     public static int Version = 2;
22 }
23
```

Figure 5. Screenshot taken from dnSpy that contains the RedLine Stealer sample’s encrypted configuration block.

We implemented a decryption routine in Python as shown below in Figure 6. We invite readers to manually grab example ciphertexts and keys to test whether [the script from Figure 6](#) decrypts correctly.



```
@staticmethod
def decrypt(ctx: bytes, key: bytes) -> Optional[bytes]:
    try:
        return b64decode(bytes([a ^ b for (a, b) in zip(b64decode(ctx), cycle(key))]))
    except (ValueError, binascii.Error):
        return None
```

Figure 6. The decryption routine written in Python.

Next, we located the configuration (shown in Figures 7 and 8) and prepared the decrypt function in Python. However, before decrypting the data, we had to manually grab the ciphertext and key from the decompiled result

generated by dnSpy.

When writing C code, we directly access system memory, so we sometimes call the executables compiled from C code as native executables. However, in [.NET MSIL](#), everything is managed. A pointer leads to the character array stored somewhere in the binary in native C code, but all we see in the compiled MSIL are the tokens.

When accessing these tokens, the runtime library ([CLR](#)) parses where the ciphertext is actually stored, which is one less thing for an analyst to worry about. For example, in Figure 7 comments generated by dnSpy show that the string IP is a token number 0x04000013.

```

1 using System;
2 using System.Reflection;
3
4 // Token: 0x02000017 RID: 23
5 [Obfuscation(ApplyToMembers = true, Exclude = true, StripAfterObfuscation = true)]
6 public static class Arguments
7 {
8     // Token: 0x04000013 RID: 19
9     public static string IP = "HR8RAYgYPRI+GTdWPhwgCDsPIwQeDxFE";
10
11     // Token: 0x04000014 RID: 20
12     public static string ID = "857238304-77906307-EZ";
13
14     // Token: 0x04000015 RID: 21
15     public static string Message = "";
16
17     // Token: 0x04000016 RID: 22
18     public static string Key = "Perverters";
19
20     // Token: 0x04000017 RID: 23
21     public static int Version = 2;
22 }
    
```

Figure 7. DnSpy labeling the token number of the decompiled strings.

Next, we open the RedLine Stealer sample in IDA Pro and navigate to the same function. Figure 8 shows that the [ldstr](#) commands push object reference for the metadata strings located at seg000:29F1, seg000:29FB, seg000:2A05 and seg000:2A0F. The object references are enclosed in black boxes in Figure 8.

These metadata strings are set by instructions located at seg000:29F6, seg000:2A00, seg000:2A0A and seg000:2A14 respectively. The [stsfld](#) instructions replace the value of a static field with a value from the evaluation stack. The values on the evaluation stack for each field are enclosed in red boxes.

```

seg000:25A0      .method private static hidebysig specialname rtspecialname void .cctor()
seg000:25A0      {
seg000:25A0      }
seg000:25A0      .maxstack 8
seg000:25A0      nop
seg000:25A1      72 FF 08 00 70 ldstr  aCieqviapnmm4nf // "CieqViAPNnM4NFmUKwkKACozNR88LSQdKgYwTg"...
seg000:25A6      80 03 00 00 04 stsfld string Arguments::IP
seg000:25A8      72 A3 03 00 70 ldstr  asc_A3A2 // ""
seg000:25B0      80 04 00 00 04 stsfld string Arguments::ID
seg000:25B5      72 A3 03 00 70 ldstr  asc_A3A2 // ""
seg000:25BA      80 05 00 00 04 stsfld string Arguments::Message
seg000:25BF      72 51 09 00 70 ldstr  aGuggles // "Guggles"
seg000:25C4      80 06 00 00 04 stsfld string Arguments::Key
seg000:25C9      18 ldc.i4.2
seg000:25CA      80 07 00 00 04 stsfld int32 Arguments::Version
seg000:25CF      2A ret
seg000:25CF      }
seg000:25CF      }
    
```

Figure 8. IDA Pro disassembler view of the configuration setup function.

The IP field from Figure 7 is not enough to statically extract the configuration. The source of the string that was pushed onto the stack for the IP field has not yet been identified. The operand type of the instruction `ldstr` shown in Figure 8 is, [according to Microsoft](#), a string token, and string tokens are stored in the `#US` (User-Stream) table.

To find the string token, we used an open-source library called [dnfile](#), which is like a .NET version of [PEfile](#). Dnfile allows us to easily access the [#US tokens](#) by just giving the [.NET runtime identifier](#) (RID). Dnfile also provides the interface to access the user streams and [a lot more](#).

The Python implementation shown in Figure 9 is an example of how we accessed user streams by offset. We passed the user string into the decryption routine shown in Figure 9 once we got the user stream by the token. This should return the decrypted configuration.

```
7 def get_us_stream_by_offset(dn: dnfile.dnPE, offset: int) -> Optional[bytes]:
8     us: dnfile.stream.UserStringHeap = dn.net.metadata.streams.get(b"#US", None)
9     if not us:
10         return None
11
12     if us.sizeof() == 0:
13         return None
14
15     ret: Optional[Tuple[bytes, int]] = us.get_with_size(offset)
16     if ret is None:
17         return None
18     buf, _ = ret
19     try:
20         s = dnfile.stream.UserString(buf[:-1])
21         return s.value.encode()
22     except UnicodeDecodeError:
23         return None
```

Figure 9. An implementation that uses dnfile to get the resource by a given .NET MSIL token.

Conclusion

By delving into the methods used for GuLoader and RedLine Stealer, we shed light on the process of locating and extracting C2 configurations from various malware families.

Leveraging our insights gained from analyzing these malware configurations, we can enhance our ability to detect, analyze and develop effective countermeasures against malicious software. Through continuous collaboration and knowledge sharing, we can collectively stay ahead of cybercriminals to help safeguard our digital systems and networks.

Palo Alto Networks customers are better protected from the threats discussed in this article through the following products:

- [Next-Generation Firewall](#) with [cloud-delivered security services](#) including [WildFire](#) detect the files mentioned within this report as malicious.

If you think you might have been compromised or have an urgent matter, get in touch with the [Unit 42 Incident Response team](#) or call:

- North America Toll-Free: 866.486.4842 (866.4.UNIT42)

- EMEA: +31.20.299.3130
- APAC: +65.6983.8730
- Japan: +81.50.1790.0200

Palo Alto Networks has shared these findings with our fellow Cyber Threat Alliance (CTA) members. CTA members use this intelligence to rapidly deploy protections to their customers and to systematically disrupt malicious cyber actors. Learn more about the [Cyber Threat Alliance](#).

Indicators of Compromise

SHA256 Hash of the GuLoader Sample Analyzed in This Article

- 32ea41ff050f09d0b92967588a131e0a170cb46baf7ee58d03277d09336f89d9

SHA256 Hash of the RedLine Stealer Sample Analyzed in This Article

- a4cf69f849e9ea0ab4eba1cdc1ef2a973591bc7bb55901fdbceb412fb1147ef9

Source: <https://unit42.paloaltonetworks.com/malware-configuration-extraction-techniques-guloader-redline-stealer/>